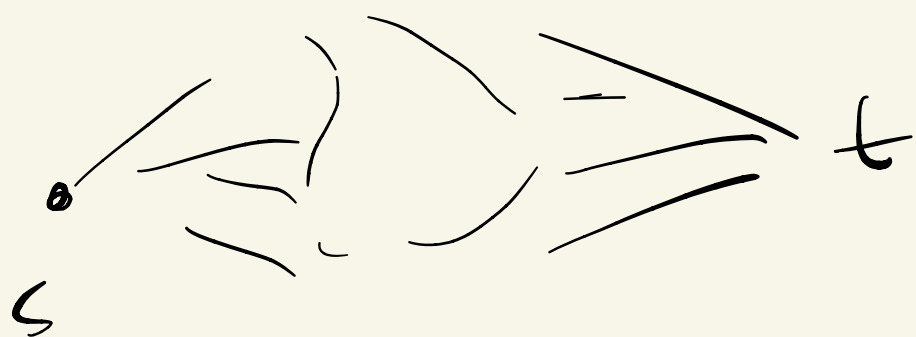


Course overview

ideas from convex geometry and
convex optimization changed the
way we think about algorithms

max from problem



5 lectures \rightarrow basic theorems convex set
separating hyperplane
theorems
etc

6-17 lectures \rightarrow optimization algorithms and some applications min cut
Caratheodory theorem
matroid intersection

18-24 lectures \rightarrow convex geometry

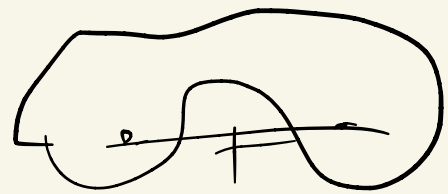
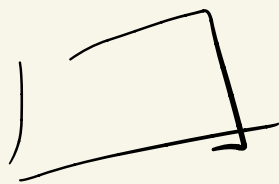
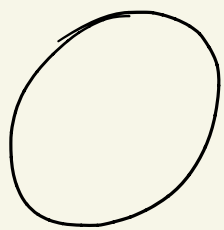
Convex sets

$C \subseteq \mathbb{R}^n$ is convex iff $\forall x, y \in C, 0 \leq \theta \leq 1$

$\theta x + (1-\theta)y$ | a set is convex if ^{for} any two points

that lie in the set also the line between

them is contained in the set.



practical purposes enough to check $\frac{x+y}{2}$

convex combination

$$\theta_1 x_1 + \dots + \theta_k x_k \quad x_1, x_2, \dots, x_k \quad \text{if } \theta_i \geq 0$$

$$\text{and } \theta_1 + \dots + \theta_k = 1$$

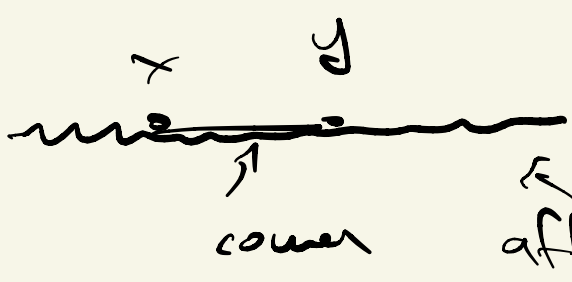
convex hull

set of convex combinations

$$C \rightarrow \text{conv } C = \{ \theta_1 x_1 + \dots + \theta_k x_k \mid x_i \in C, \theta_i \geq 0, 1 \leq i \leq k, \theta_1 + \dots + \theta_k = 1 \}$$

Affine set

A set $C \subseteq \mathbb{R}^n$ is affine if for any $x, y \in C$ and $\theta \in \mathbb{R}$ we have $\theta x + (1-\theta)y$

Difference with convex sets = 

affine combination
affine hulls

$$\Rightarrow \theta_1 + \dots + \theta_k = 1$$
$$\theta_i \in \mathbb{R}$$

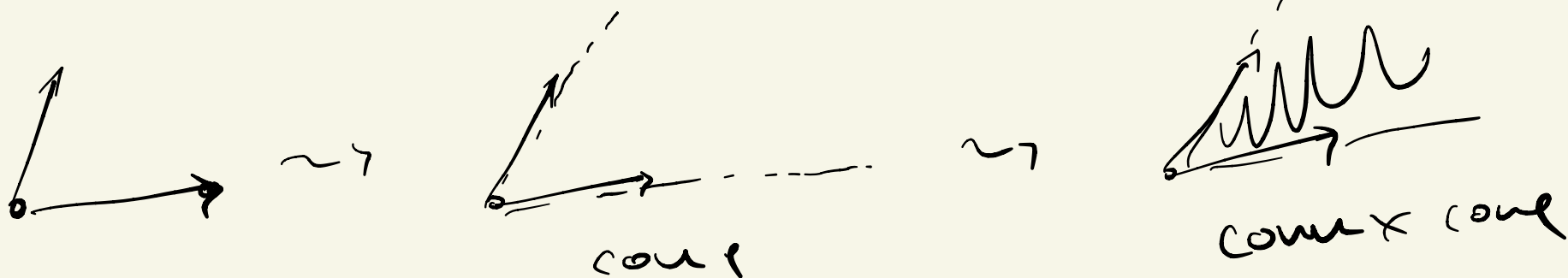


Cone

$C \subseteq \mathbb{R}^n$ is a cone if $\forall x \in C, \theta > 0, \theta x \in C$

// // convex cone if for every

$$x_1, x_2 \in C, \theta_1, \theta_2 > 0, \theta_1 x_1 + \theta_2 x_2 \in C$$



Polyhedra

hyperplane $\{x \mid a^T x = b\}$

half space $\{x \mid a^T x \leq b\}$

polyhedron = intersection of finite number of

banded polyhedron = polytope

LP \rightarrow optimize a linear function over a polyhedron

Positive definite case

$S_n = \{X \in \mathbb{R}^{n \times n} \mid X^T = X\} \leftarrow$ set of real symmetric matrices,

let $U \in S_n$ then the spectral theorem for symmetric matrices tells us that

$$U = V \Lambda V^T = \begin{bmatrix} | & & | \\ v_1 & \dots & v_n \\ | & & | \end{bmatrix} \begin{bmatrix} \lambda_1 & & \\ & \ddots & \\ & & \lambda_n \end{bmatrix} \begin{bmatrix} - & & - \\ v_1 & & \\ & \ddots & \\ - & & v_n \end{bmatrix}$$

$\underbrace{\hspace{10em}}_{\text{real}} \quad \underbrace{\hspace{10em}}_{\text{eigenvalues orthonormal}}$

positive definite ($\mu > 0$) (if $\lambda_i > 0$)

equivalent characterization

① $\lambda_i > 0$

② $x^T \mu x > 0 \ \forall x \in \mathbb{R}^n$

③ $\mu = V^T \cdot V$ for some $V \in \mathbb{R}^{n \times n}$

(positive definite $\Leftrightarrow \mu > 0$)

$S_+^n = \{x \in S^n \mid x > 0\} \leftarrow$ is a convex cone
 $x^T(\theta_1 A + \theta_2 B)x > 0$

$S_{++}^n = \{x \in S^n \mid x > 0\}$

Semidefinite programming

optimize a linear function over the

positive semidefinite cone (half space)
 \nearrow intersecting

Ellipsoids (important)

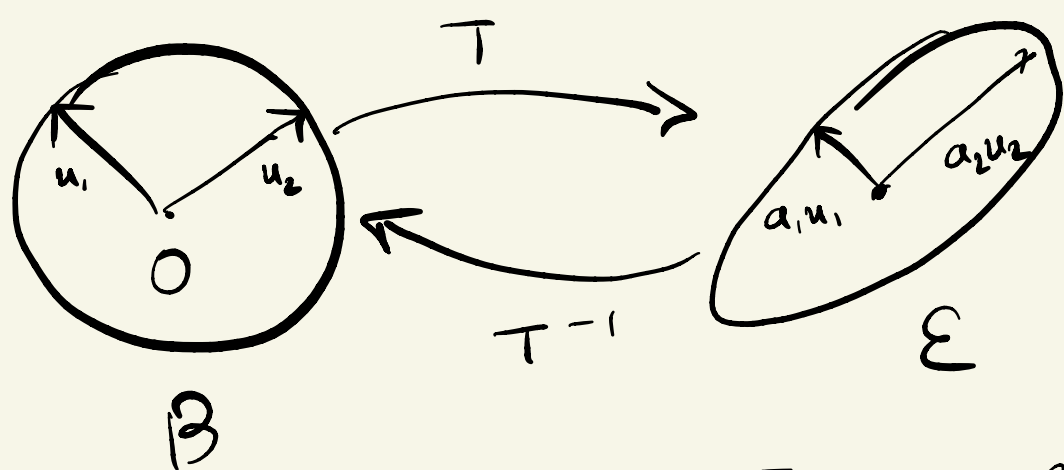
Informally: $\text{Ball} \xrightarrow[\text{+ transformation}]{\text{affine}} \text{ellipsoid}$

$$B(x_c, r) = \{x \mid \|x - x_c\|_2 \leq r\} \quad \begin{array}{l} x_c \text{ center} \\ r \text{ radius} \end{array}$$

$$\|u\|_2 = \sqrt{u^T u}$$

$u_1, \dots, u_n \leftarrow$ set of orthonormal vectors (semi-axes)

$a_1, a_2, \dots, a_n \geq 0$ scaling factors



\leftarrow I stretched more along the u_2 direction

$$x \mapsto Tx$$

$$T = \begin{bmatrix} | & & | \\ u_1 & \dots & u_n \\ | & & | \end{bmatrix} \begin{bmatrix} a_1 & & \\ & \ddots & \\ & & a_n \end{bmatrix} \begin{bmatrix} -u_1 - \\ \vdots \\ -u_n - \end{bmatrix}$$

$$a_i \geq 0$$

$$T = U D U^T$$

$$T^{-1} = (U D U^T)^{-1} = (U^T)^{-1} D^{-1} U^T = U D^{-1} U^T$$

$$y \in E \Leftrightarrow T^{-1} y \in B \Rightarrow \|U D^{-1} U^T y\|_2 \leq 1 \Rightarrow$$

$$\Rightarrow \|U D^{-1} U^T y\|_2 \leq 1 \Rightarrow y^T \underbrace{U D^{-1} U^T}_T U D^{-1} U^T y \leq 1 \Rightarrow$$

$$\Rightarrow y^T U D^{-2} U^T y \leq 1$$

$$y^T U \bar{D}^2 U^T y \leq 1 \Rightarrow$$

$$\Rightarrow \mathcal{E} := \{y \mid \sum_{i=1}^n \frac{(y^T u_i)^2}{\alpha_i^2} \leq 1\}$$

another representation

\rightarrow the bigger α_i the more stretch I can get in this direction

$$P \in S_{++}^n \quad (P = U \bar{D}^2 U^T)$$

$$\mathcal{E} = \{y \mid y^T P y \leq 1\}$$

Note that $P \in S_{++}^n \Rightarrow P = V D' V^T$

let $P'^{1/2} = V (D')^{1/2} V^T$ then $y^T P y \leq 1 \Rightarrow$

$$\Rightarrow \|P'^{1/2} y\|_2 \leq 1$$

$$\mathcal{E} = \{y \mid \|P'^{1/2} y\| \leq 1\}$$

\rightarrow inverse representation
you go back to the ball

Why ellipsoids?

simple enough objects to have closed

formulas for quadrants

line

→ optimizing a linear function over an ellipsoid

→ computing the volume

$$\left(\prod_{i=1}^n a_i \right), \text{vol}(\mathcal{B}_2^n)$$

//

$$\det(T)$$

Jouli's theorem

rich enough to approximate reasonably

"well" any convex body

Norm and norm balls

$$\|\cdot\| : \mathbb{R}^n \rightarrow \mathbb{R}_+$$

$$(1) \|x\| = 0 \text{ iff } x = 0$$

$$(2) \|t \cdot x\| = |t| \cdot \|x\| \quad \forall t \in \mathbb{R}, x \in \mathbb{R}^n$$

$$(3) \forall x, y \in \mathbb{R}^n \quad \|x+y\| \leq \|x\| + \|y\|$$

Common norms

$$l_\infty \quad \|x\|_\infty = \max \{ |x_1|, \dots, |x_n| \}$$

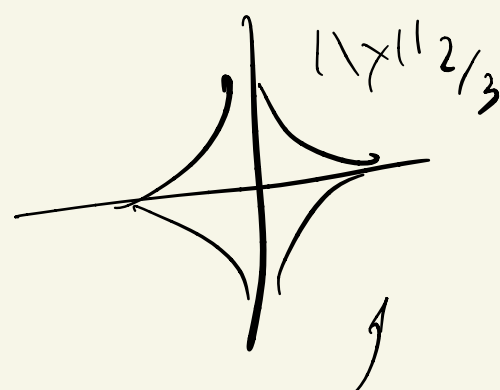
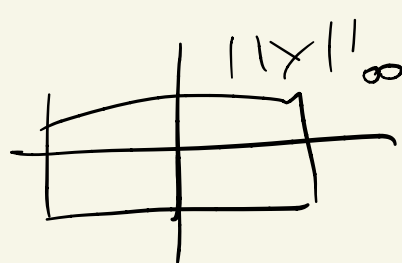
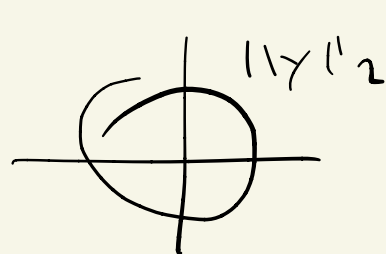
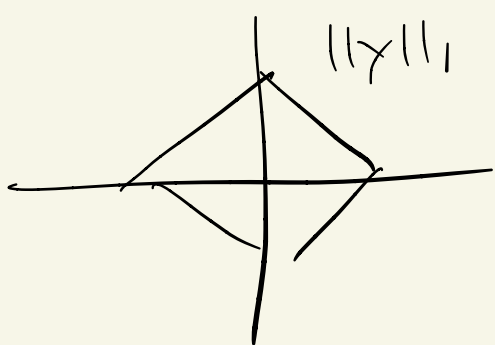
$$l_p \quad \|x\|_p = \left(\sum_{i=1}^n |x_i|^p \right)^{1/p} \quad p \geq 1$$

quadratic norm

$$\|x\|_M = \sqrt{x^T M x} = \|M^{1/2} x\|_2, \quad M \in S_{++}^n$$

given a norm we can define a

unit ball as $B = \{x \mid \|x\| \leq 1\}$



Euclidean balls are 2 dimensional

just some special cases of norm balls

unit ball
for l_p
with $p < 1$ (not norm)

Convex functions

$f: \mathbb{R}^n \rightarrow \mathbb{R}$ is convex if $\text{dom } f$ is convex
and $\forall x, y \in \text{dom } f, \theta \in [0, 1]$

$$f(\theta x + (1-\theta)y) \leq \theta f(x) + (1-\theta)f(y)$$

strictly convex $\leadsto \leq \leadsto <$

concave $\leadsto \geq \leadsto >$

epigraph connection between convex functions and convex sets

$\text{epi } f = \{(x, t) : x \in \text{dom } f, t \geq f(x)\}$

lemma

f is convex iff $\text{epi } f$ is a convex set

f is convex $\Rightarrow \text{epi } f$ convex set

$(x_1, t_1), \dots, (x_n, t_n) \in \text{epi } f, \lambda_1, \dots, \lambda_n \in [0, 1], \sum \lambda_i = 1$

$(x, t) = (\sum x_i \lambda_i, \sum \lambda_i t_i) \in \text{epi } f?$

$$f(\sum x_i \lambda_i) \leq \sum \lambda_i f(x_i) \leq \sum \lambda_i t_i \quad \checkmark$$

epi f is convex set $\Rightarrow f$ is convex

$$(x, f(x)) \in \text{epi } f$$

$$(y, f(y)) \in \text{epi } f$$

want to prove

$$f(\theta x + (1-\theta)y) \leq \theta f(x) + (1-\theta)f(y)$$

\Rightarrow

$$(\theta x + (1-\theta)y, \underbrace{\theta f(x) + (1-\theta)f(y)}_{\geq f(\theta x + (1-\theta)y)}) \in \text{epi } f$$