

Optimization for lattices, packings, and coverings

Lecture 3

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G.F. Voronoi (1868–1908)

2. Voronoi's lattice reduction theory

Parameter space of lattice & Voronoi's second reduction theory

$$\tilde{S}_+^n = \text{cone}_{\mathbb{Q}}\{xx^T : x \in \mathbb{Z}^n\} \cup S_{++}^n \subset S_+^n \text{ rational closure of } S_{++}^n$$

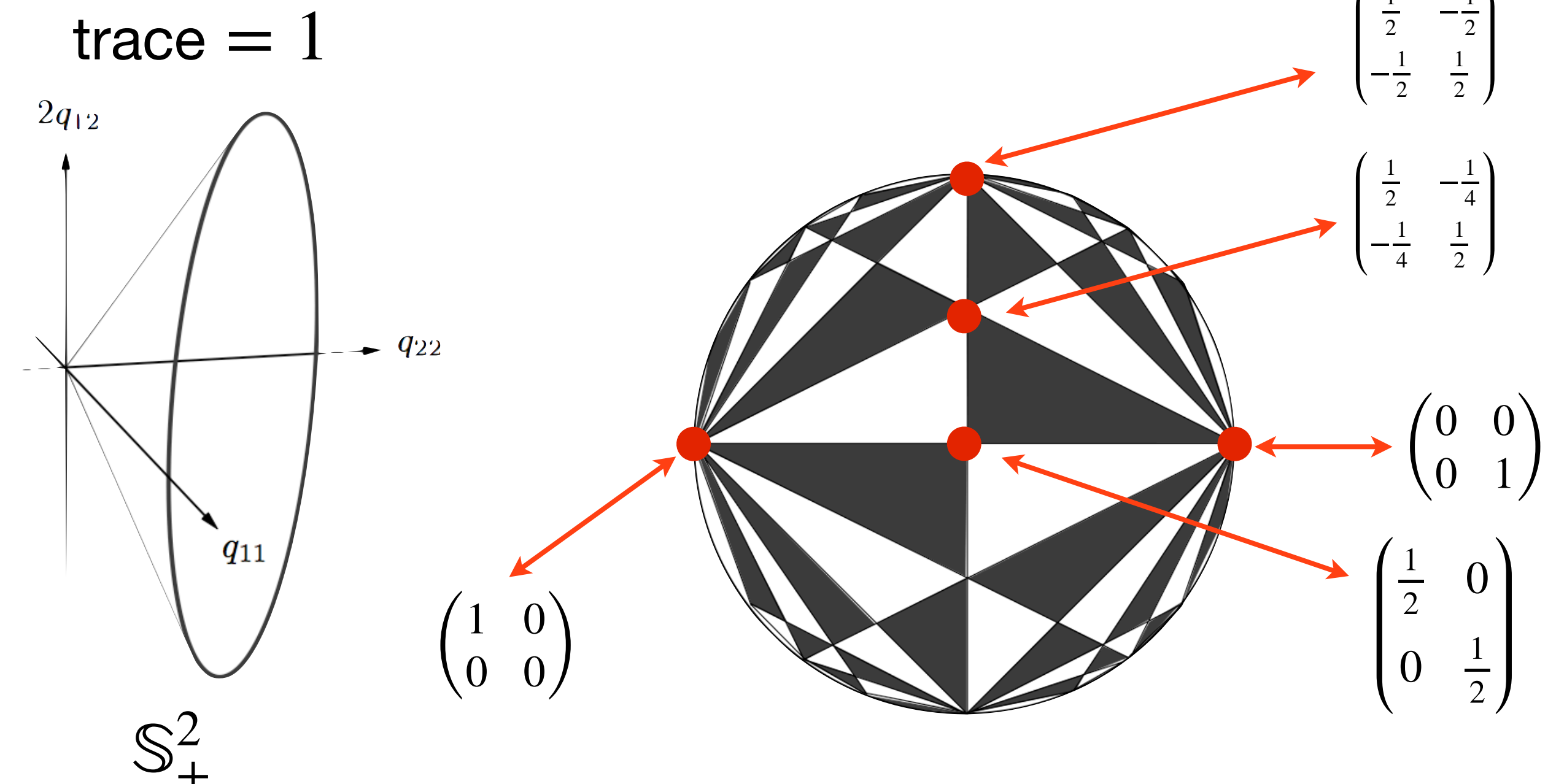
$$GL_n(\mathbb{Z}) \text{ acts on } \tilde{S}_+^n \text{ by } (g, Q) \mapsto g^T Q g$$

reduction theory of lattices = find „nice“ fundamental domain for $\tilde{S}_+^n / GL_n(\mathbb{Z})$

properties

infinite polyhedral face-to-face tiling

all triangular polyhedra $GL_2(\mathbb{Z})$ -equivalent





Construction of Delaunay polyhedra



B.N. Delaunay
= Б.Н. Делоне
= B.N. Delone (1890–1980)

Empty sphere construction

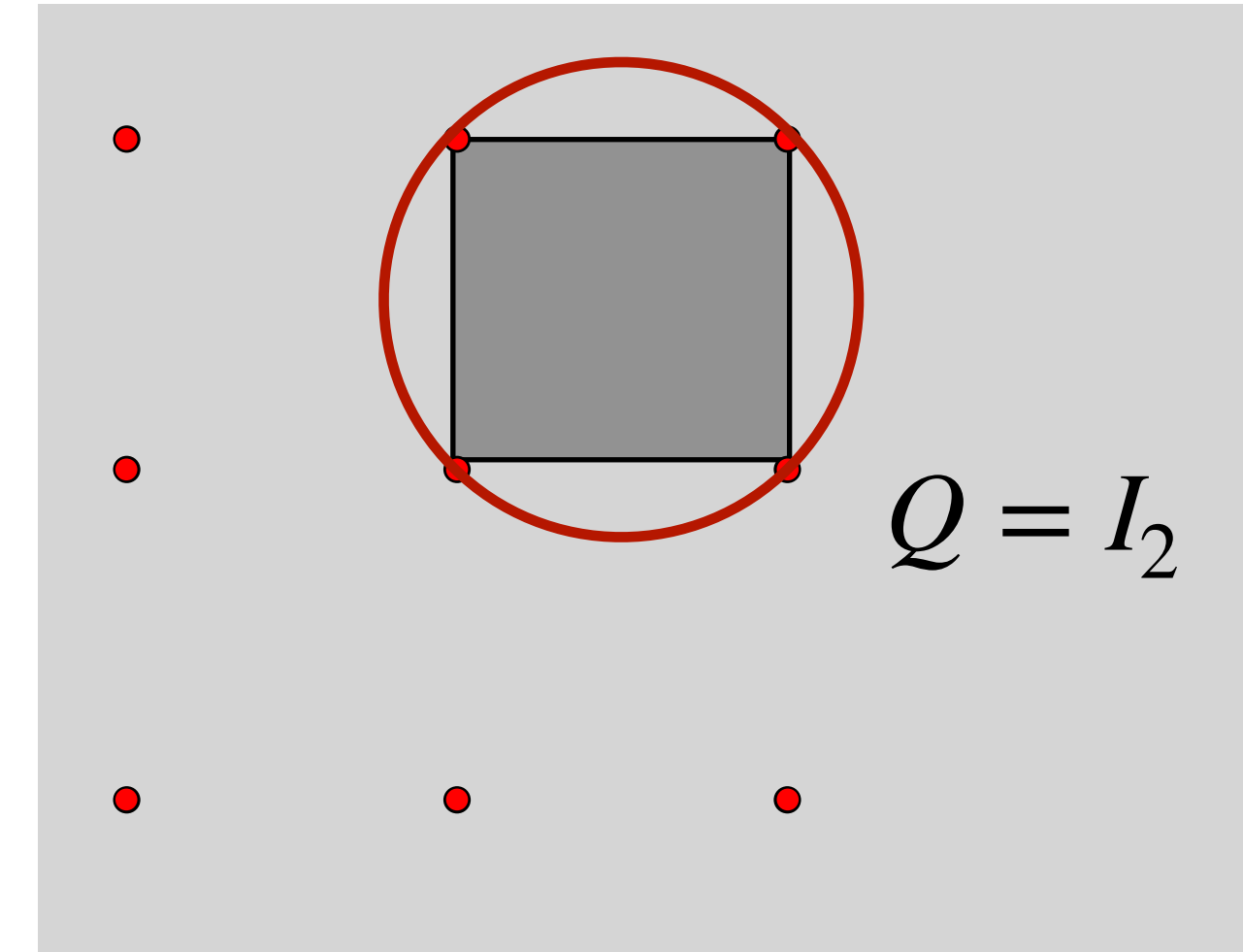
Delaunay polyhedra of $Q \in \tilde{S}_+^n$

$$P = \text{conv}\{v_1, v_2, \dots\}, v_i \in \mathbb{Z}^n, \text{ where}$$

there is a center $c \in \mathbb{R}^n$ and a radius $r > 0$ so that

$$Q[v_i - c] = r^2 \text{ and } Q[w - c] > r^2 \text{ for all } w \in \mathbb{Z}^n \setminus \{v_1, v_2, \dots\}$$

with $Q[x] = x^T Q x$

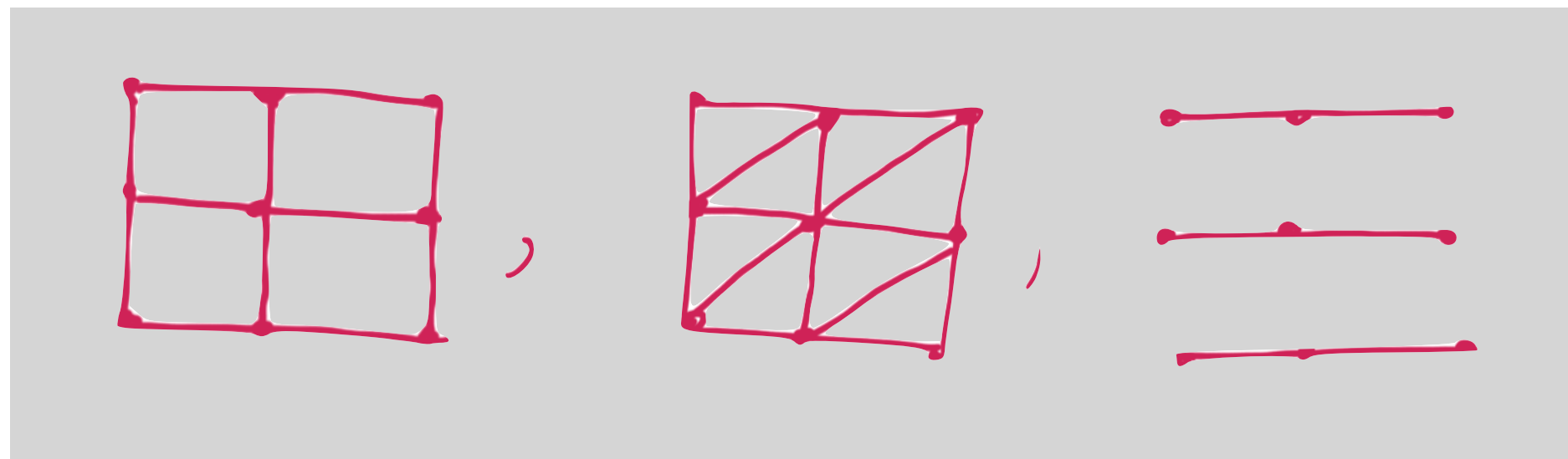




Delaunay subdivisions & secondary cones

Delaunay subdivision of $Q \in \tilde{S}_+^n$

$\text{Del}(Q) = \{P : P \text{ is a Delaunay polyhedron of } Q\}$ is a polyhedral complex



Secondary cone of Delaunay subdivision

$$\Delta(\text{Del}(Q)) = \{Q' \in \tilde{S}_+^n : \text{Del}(Q') = \text{Del}(Q)\}$$

is (open) polyhedral cone in \tilde{S}_+^n ; facets of $\Delta(\text{Del}(Q))$ are determined by facets in $\text{Del}(Q)$



Some facts about secondary cones

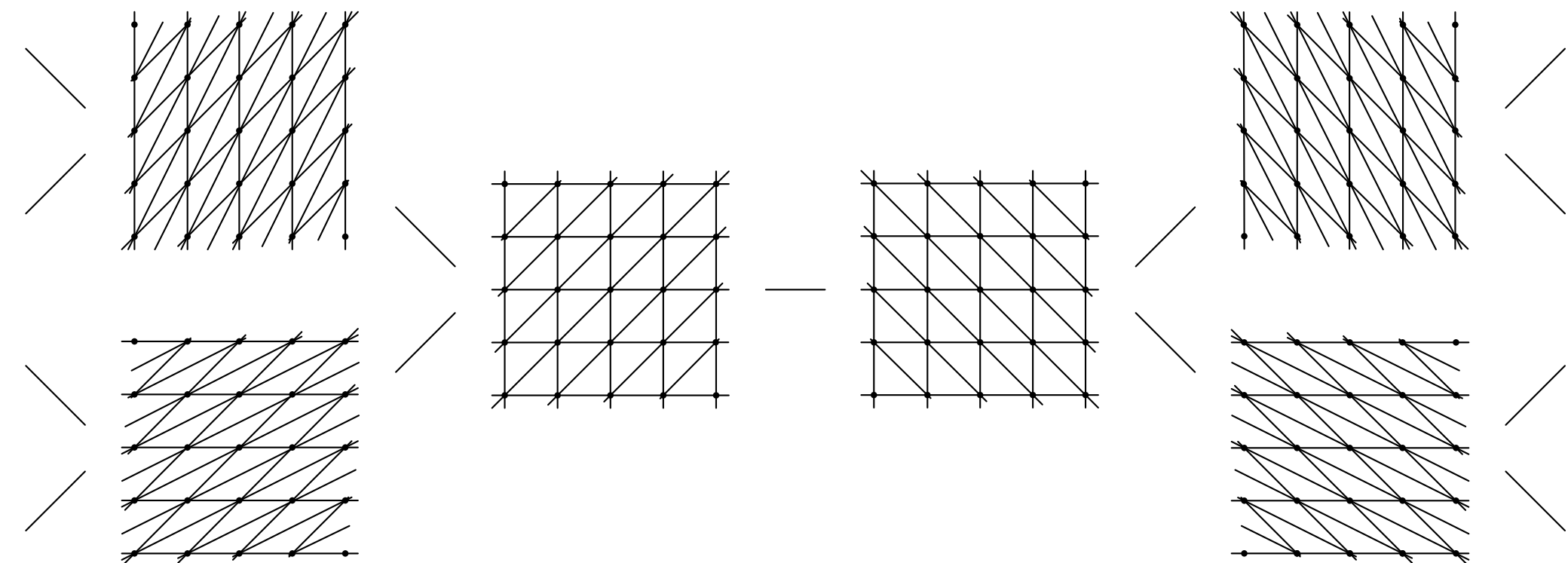
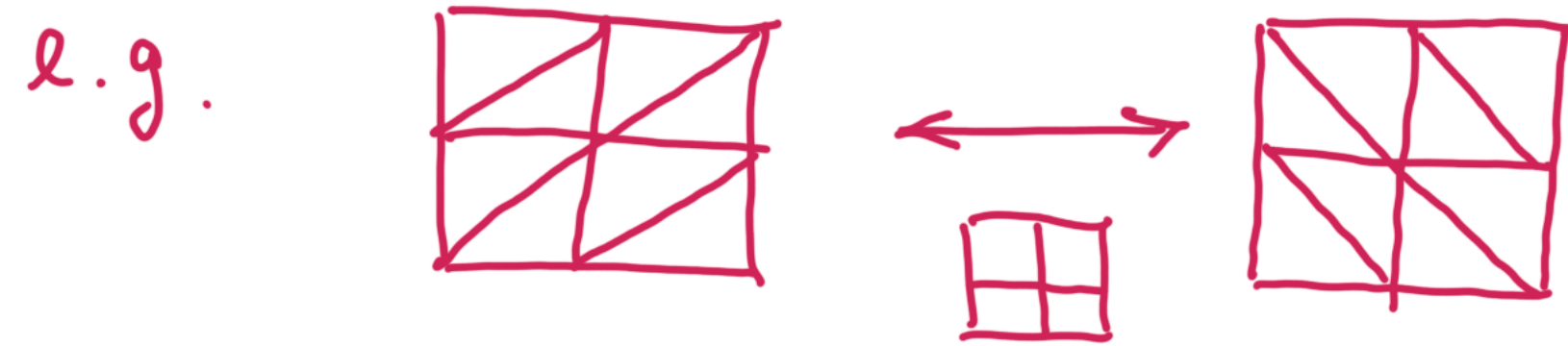
$\Delta(\text{Del}(Q))$ is full-dimensional $\iff \text{Del}(Q)$ is a triangulation

$\Delta(\text{Del}(Q)) \subseteq \overline{\Delta(\text{Del}(Q'))} \iff \text{Del}(Q')$ is a refinement of $\text{Del}(Q)$



$\text{Del}(Q), \text{Del}(Q')$ triangulations, $\dim(\overline{\Delta(\text{Del}(Q))} \cap \overline{\Delta(\text{Del}(Q'))}) = \binom{n+1}{2} - 1$

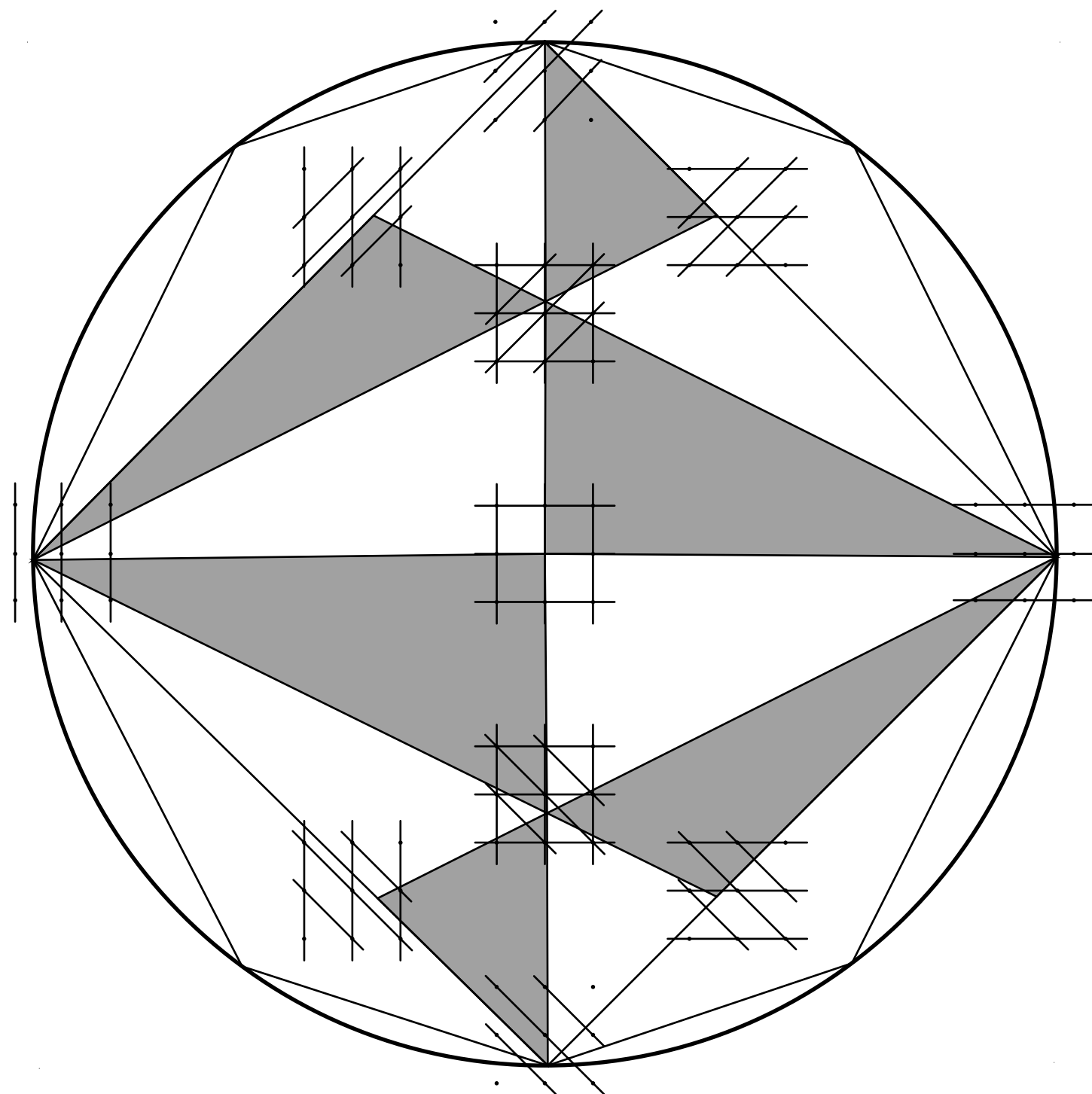
$\iff \text{Del}(Q), \text{Del}(Q')$ differ by a *bistellar flip*



Main theorem of Voronoi's second reduction theory

There are only finitely many n -dimensional Delaunay triangulations up to $GL_n(\mathbb{Z})$ -equivalence.

Furthermore: $\tilde{\mathcal{S}}_+^n = \bigcup_{\mathcal{D} \text{ Delaunay triangulation}} \overline{\Delta(\mathcal{D})}$



Known classification

n	# triang.	
2	1	} Voronoi (1908)
3	1	
4	3	
5	221+1	Baranovski &
6	many...	Ryshkov (1976)

3. Application to lattice sphere covering

Lattice sphere covering problem

lattice sphere covering problem

given: dimension $n \in \mathbb{N}$

find: lattice $L \subseteq \mathbb{R}^n$ s.t. covering density $\Theta(L)$ is minimized

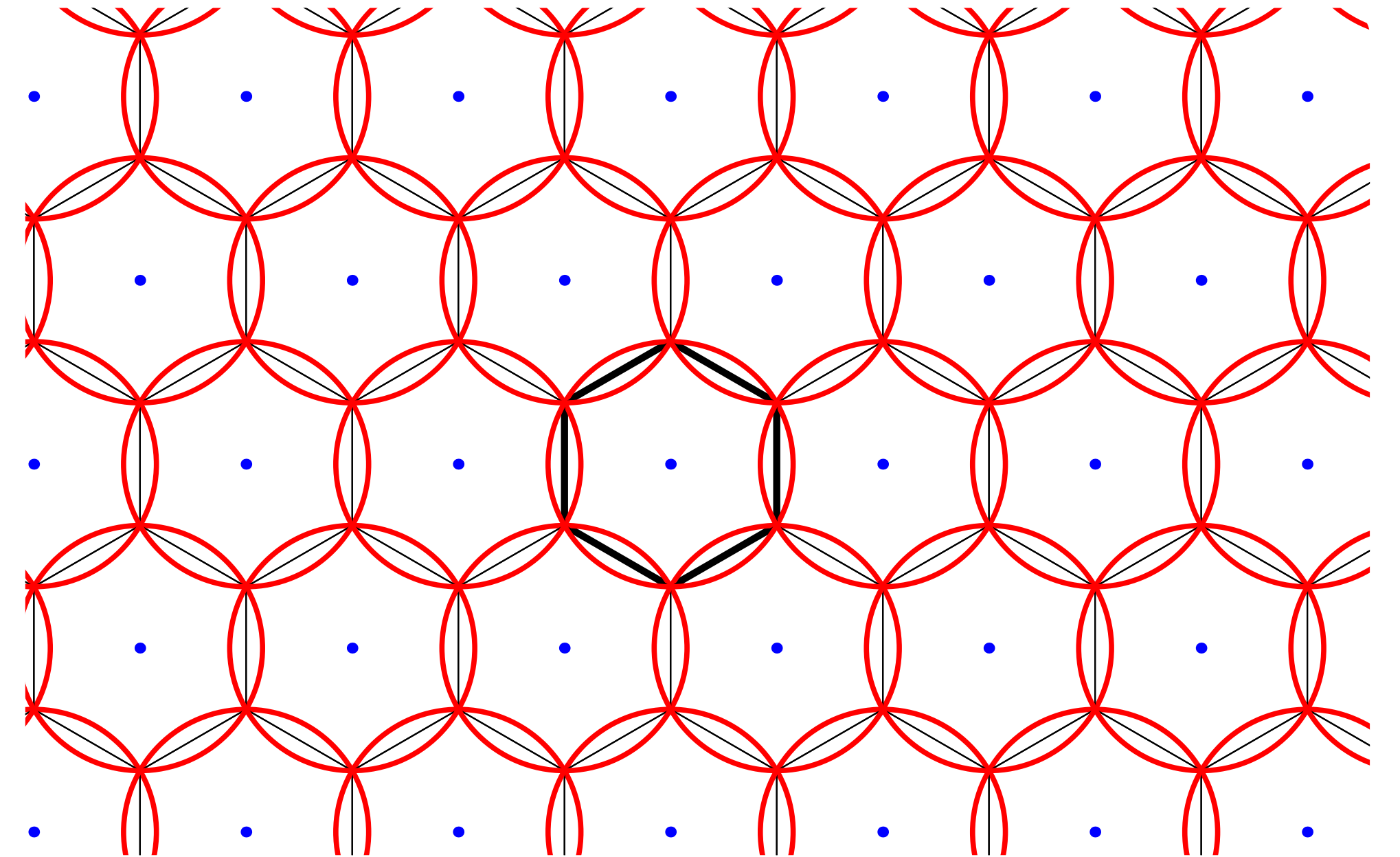
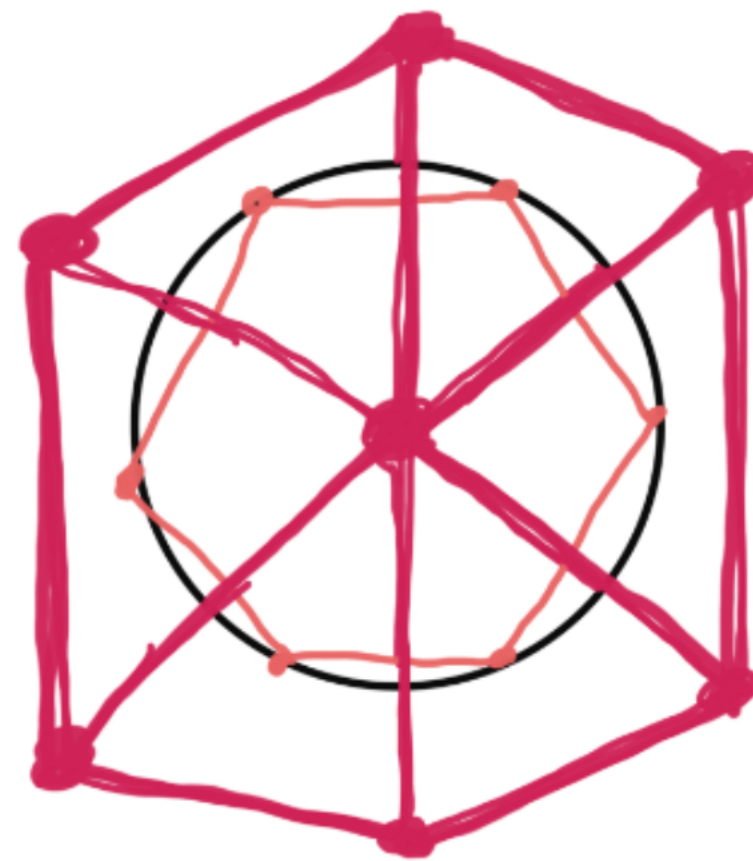
$$\Theta(L) = \frac{\mu(L)^n}{\det L} \cdot \text{vol } B_n$$

\uparrow unit ball

covering radius

$$= \max_{x \in \mathbb{R}^n} \min_{v \in L} \|x - v\|$$

= max. circumradius of Delaunay polytope



$$\Theta(L) = 1.20 \dots$$

MAXDET optimization

→ see Vandenberghe, Boyd, Wu (1998)

→ here: conic formulation

$$\mathcal{D}^n = \{ (X, s) \in S_+^n \times \mathbb{R}_+ : (\det X)^{1/n} \geq s \} \text{ proper convex cone}$$

computation of its dual cone by

• Minkowski's determinant inequality $\forall X, Y \in S_+^n : \det(X+Y)^{1/n} \geq (\det X)^{1/n} + (\det Y)^{1/n}$

• AM-GM inequality $\forall X \in S_+^n : \text{Tr}(X) - n(\det X)^{1/n} \geq 0$

gives

$$(\mathcal{D}^n)^* = \{ (Y, t) \in S_+^n \times \mathbb{R} : (\det Y)^{1/n} \geq -\frac{t}{n} \}$$

→ can optimize over \mathcal{D}^n and $(\mathcal{D}^n)^*$ in polynomial time

(under usual conditions for the ellipsoid method)

Basic idea

Fix Delaunay triangulation \mathcal{D} , then the restricted lattice covering problem becomes a tractable MAXDET problem.

$$\max (\det Q)^{1/n}$$

$$\text{s.t. } Q \in \overline{\Delta(\mathcal{D})}$$

(linear cond.)

$$\mu_2(Q) \leq 1$$

for every Delaunay
simplex $L = \text{conv}\{0, v_1, \dots, v_n\}$



Gadget

Lemma $L = \text{conv} \{0, v_1, \dots, v_n\}$ simplex. Then

$\mu_2(Q) \leq 1$ iff the following matrix is in S_+^{n+1} :

$$\begin{bmatrix} 4 & v_1^T Q v_1 & v_2^T Q v_2 & \dots & v_n^T Q v_n \\ v_1^T Q v_1 & & & & \\ \vdots & & & & \\ v_n^T Q v_n & & & & \end{bmatrix}$$

$v_i^T Q v_j$

\leadsto Cayley - Menger determinant

Algorithmic approach to lattice covering

Schürmann, Vallentin (2006):

C++ implementation of Voronoi's second reduction theory (scc = second cone cruiser)

finding solutions for dimensions $n \leq 5$,

reproving results of Baranovski, Ryshkov (1976)

finding new, conjecturally optimal lattices for $n \geq 6$

2020: still only conjectures

The Λ_{24} sphere covering problem

The Leech lattice gives clearly the least dense sphere covering in dimension 24.

Currently, we don't have a good method (like LP bounds) to come anywhere close to a proof.

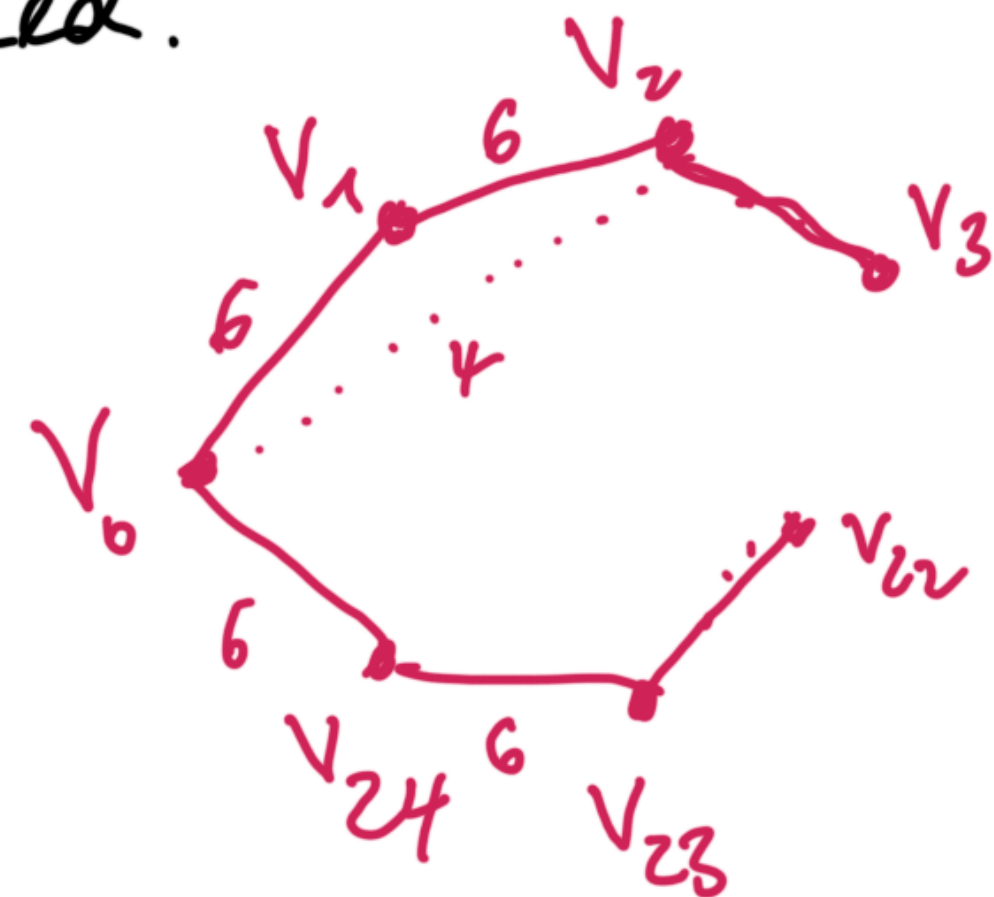
Λ_{24} is locally optimal for covering

Delannoy polytopes of Λ_{24} were classified by
Conway, Parker, Sloane (1982)
and Borcherds, Conway, Queen (1984)

Important class of Delannoy simplices
where covering radius is attained.

$$A_{25} = \text{Conv} \{V_0, V_1, \dots, V_{24}\}$$

center of circumsphere = barycenter



Proof by MAXDET

- first proof (but more complicated) by Schürmann, Valentin (2005)
- here: new proof using MAXDET framework (with support from Krupp, Rolfe and Zimmermann)

Relaxation

$$\frac{1}{2} \leq \max_{Q \in S_+^{24}} (\det Q)^{1/n}$$

↑
because Λ_{24} is feasible solution

$$\mu_{TA_{25}}(Q) \leq 1 \quad \forall T \in C_0$$

$$= \max s$$

$$(Q, s) \in \mathcal{D}^n$$

$$V_T \in S_+^{n+1}, T \in C_0$$

$$\langle V_T, E_{00} \rangle = 4$$

$$\langle V_T, E_{0i} \rangle = \langle Q, V_{T,i} V_{T,i}^T \rangle \quad i = 1, \dots, n$$

$$\langle V_T, E_{ij} \rangle = \langle Q, \frac{1}{2} (V_{T,i} V_{T,j}^T + V_{T,j} V_{T,i}^T) \rangle \quad \begin{matrix} 1 \leq i < j \leq n \\ T \in C_0 \end{matrix}$$

Dualization

$$= \max s$$

$$(Q, s) \in \mathcal{Q}^n$$

$$V_T \in S_+^{m+1}, T \in C_0$$

$$\lambda_{T,0} \langle V_T, E_{00} \rangle = 4$$

$$\lambda_{T,i} \langle V_T, E_{0i} \rangle = \langle Q, V_{T,i} V_{T,i}^T \rangle \quad i=1, \dots, n$$

$$\lambda_{T,ij} \langle V_T, E_{ij} \rangle = \langle Q, \frac{1}{2} (V_{T,i} V_{T,j}^T + V_{T,j} V_{T,i}^T) \rangle \quad \substack{1 \leq i < j \leq n \\ T \in C_0}$$

$$\leq \min 4 \sum_T \lambda_{T,0}$$

$$\lambda_{T,0} E_{00} - \sum_i \lambda_{T,i} E_{0i} - \sum_{ij} \lambda_{T,ij} E_{ij} \geq 0 \quad \forall T$$

$$\left(\underbrace{\sum \lambda_{T,i} V_{T,i} V_{T,i}^T + \sum \lambda_{T,ij} \frac{1}{2} (V_{T,i} V_{T,j}^T + V_{T,j} V_{T,i}^T)}_{Q(\lambda)}, 0 \right) - (0, 1) \in (\mathcal{Q}^n)^*$$

$$Q(\lambda) \rightarrow (\det Q(\lambda))^{1/n} \geq \frac{1}{n}$$

Ansatz

$$\leq \min_T 4 \sum_T \lambda_{T,0}$$

$$\lambda_{T,0} E_{00} - \sum_i \lambda_{T,i} E_{0i} - \sum_{i,j} \lambda_{T,i,j} E_{ij} \geq 0 \quad \forall T$$

$$\left(\underbrace{\sum \lambda_{T,i} v_{T,i} v_{T,i}^T + \sum \lambda_{T,i,j} \frac{1}{2} (v_{T,i} v_{T,j}^T + v_{T,j} v_{T,i}^T)}_{Q(\lambda)}, 0 \right) - (0,1) \in (\mathbb{D}^n)^*$$

$$Q(\lambda) \rightarrow (\det Q(\lambda))^{1/n} \geq \frac{1}{n}$$

Ansatz: Set $\alpha = \lambda_{T,0} \quad \forall T$

$$2\beta = \lambda_{T,i} \quad \forall T,i$$

$$\gamma = \lambda_{T,i,i} \quad \forall T,i$$

$$2\gamma = \lambda_{T,i,j} \quad \forall T, i,j, i \neq j$$

Then

$$\lambda_{T,0} E_{00} - \sum_i \lambda_{T,i} E_{0i} - \sum_{i,j} \lambda_{T,i,j} E_{ij} = \begin{bmatrix} \alpha & -\beta & \dots & -\beta \\ -\beta & \boxed{-\gamma} & & \\ \vdots & & \boxed{-\gamma} & \\ -\beta & & & \boxed{-\gamma} \end{bmatrix} \geq 0$$



First condition

$$\begin{bmatrix} \alpha & -\beta & \dots & -\beta \\ -\beta & \boxed{-\gamma} & & \\ \vdots & & & \\ -\beta & & & \end{bmatrix} \succ 0$$

by
 \Leftrightarrow
Scher
complement

$\alpha > 0$ and

$$-\gamma \gamma - (-\beta e) \frac{1}{\alpha} (-\beta e)^T \geq 0$$

$$\Leftrightarrow -\gamma \gamma - \frac{\beta^2}{\alpha} \gamma \geq 0$$

$$\Leftrightarrow -\gamma - \frac{\beta^2}{\alpha} \geq 0$$

Second condition

$$\left(\underbrace{\sum \lambda_{T,i} v_{T,i} v_{T,i}^T + \sum \lambda_{T,i,j} \frac{1}{2} (v_{T,i} v_{T,j}^T + v_{T,j} v_{T,i}^T)}_{Q(\lambda)}, 0 \right) - (0, 1) \in (\mathbb{D}^n)^*$$

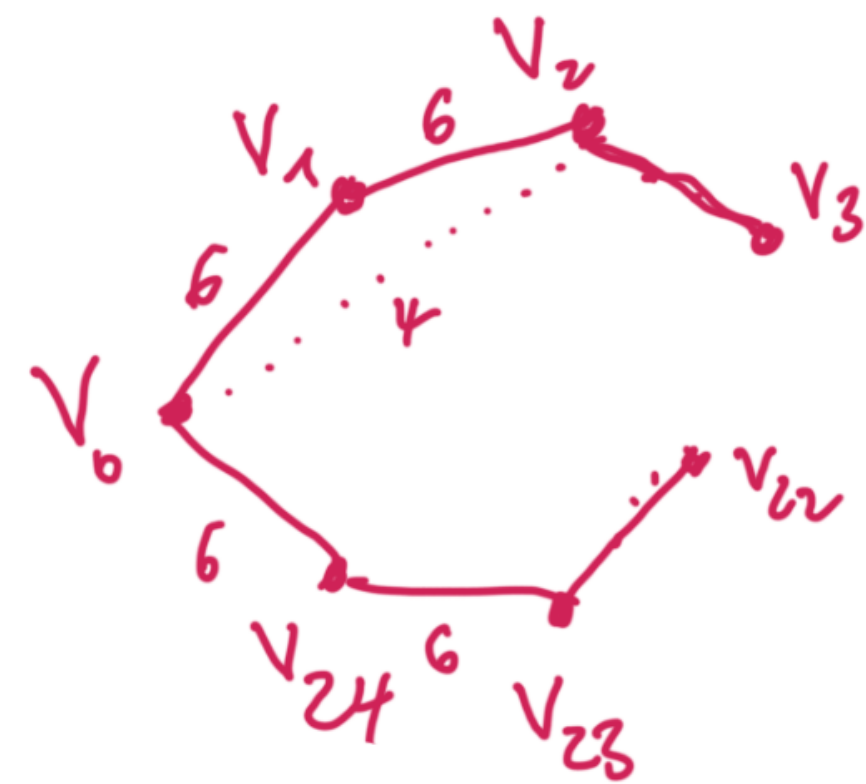
$$Q(\lambda) \rightarrow (\det Q(\lambda))^{1/n} \geq \frac{1}{n}$$

1st term

$$\sum_{T,i} \lambda_{T,i} v_{T,i} v_{T,i}^T = 2\beta \sum_T \sum_i T v_i (T v_i)^T$$

$$= 2\beta \left(\frac{22 \cdot |C_{00}|}{|S_4|} \sum_{e \in S_4} e e^T + \frac{2 |C_{00}|}{|S_6|} \sum_{e \in S_6} e e^T \right)$$

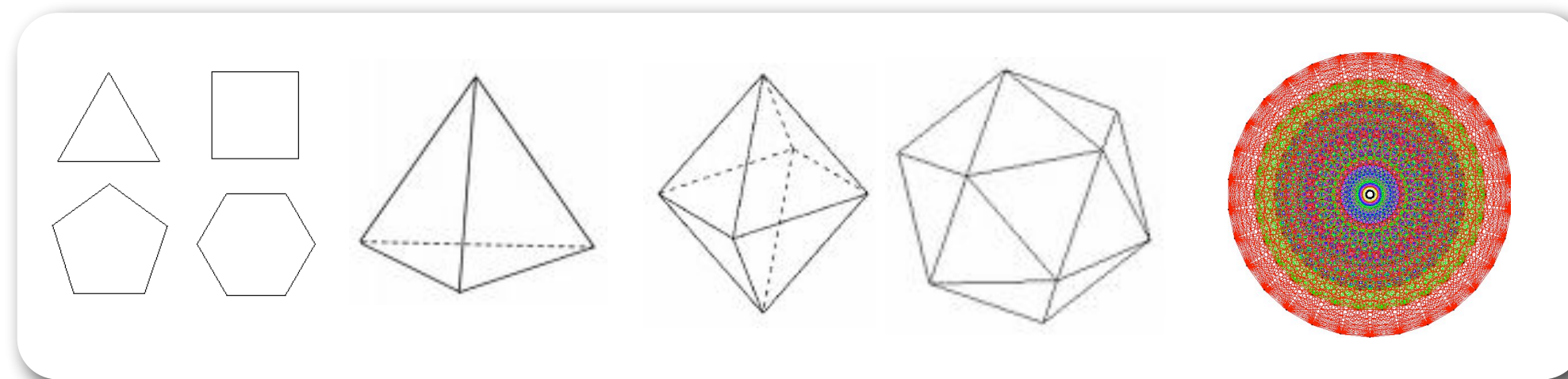
$$S_k = \{v \in \Lambda_{24} : v \cdot v = k\}$$



Spherical t -designs

$X \subseteq S^{n-1}$ spherical t -design

$$\int_{S^{n-1}} f(x) d\omega(x) = \frac{1}{|X|} \sum_{x \in X} f(x) \quad \text{for all polynomials of degree } \leq t.$$



configuration	strength
n -gon	$n - 1$
simplex	2
cross polytope	3
icosahedron	5
240	7
196560	11

Useful for optimality conditions: X spherical 2-design $\iff \sum_{x \in X} xx^T = \frac{|X|}{n} I$

Putting it together

know: $\frac{1}{T_k} S_k$ is spherical t_1 -design

$$\text{So } \sum_{e \in S_k} ee^T = \frac{k |S_k|}{24} I$$

First term of $Q(\lambda)$ equals

$$\frac{50}{6} |C_0| \beta I$$

Similarly, second term of $Q(x)$ equals

$$\frac{1250}{24} |C_0| \gamma I.$$

Together

$$\alpha = \frac{1}{|C_0|} \cdot \frac{1}{8}$$

$$\beta = \frac{1}{|C_0|} \cdot \frac{1}{100}$$

$$\gamma = \frac{1}{|C_0|} \cdot \frac{1}{1250}$$

gives optimal solution of dual problem.
This shows that Λ_{24} is locally optimal
for lattice sphere coverings. \square