

Mathematics of Data: From Theory to Computation

Prof. Volkan Cevher
volkan.cevher@epfl.ch

Lecture 12: Primal-dual optimization II

Laboratory for Information and Inference Systems (LIONS)
École Polytechnique Fédérale de Lausanne (EPFL)

EE-556 (Fall 2020)



License Information for Mathematics of Data Slides

- ▶ This work is released under a [Creative Commons License](#) with the following terms:
- ▶ **Attribution**
 - ▶ The licensor permits others to copy, distribute, display, and perform the work. In return, licensees must give the original authors credit.
- ▶ **Non-Commercial**
 - ▶ The licensor permits others to copy, distribute, display, and perform the work. In return, licensees may not use the work for commercial purposes – unless they get the licensor's permission.
- ▶ **Share Alike**
 - ▶ The licensor permits others to distribute derivative works only under a license identical to the one that governs the licensor's work.
- ▶ [Full Text of the License](#)

Outline

- ▶ This class:
 1. Algorithms for solving min-max optimization
- ▶ Next class
 1. Additional scalable optimization methods for constrained minimization

A roadmap to algorithms for convex-concave minimax optimization

Recall: A restricted minimax formulation

Let us consider

$$\min_{\mathbf{x} \in \mathcal{X}} \max_{\mathbf{y} \in \mathcal{Y}} \Phi(\mathbf{x}, \mathbf{y}), \quad (1)$$

where $\Phi(\mathbf{x}, \mathbf{y})$ is convex in \mathbf{x} and concave in \mathbf{y} .

o In the sequel, we consider the following cases

1. $\mathcal{X} \subset \mathbb{R}^p$ and $\mathcal{Y} \subset \mathbb{R}^n$; and $\Phi(\mathbf{x}, \mathbf{y})$ is smooth, or bilinear, or strongly convex/strongly concave
 - ▶ Algorithms: Proximal-Point [18], Extra-gradient [8, 13, 12], OGDA [13, 12]
2. $\mathcal{X} \subset \mathbb{R}^p$ and $\mathcal{Y} \subset \mathbb{R}^n$ with tractable “mirror maps”; and $\Phi(\mathbf{x}, \mathbf{y})$ is smooth and continuously differentiable
 - ▶ Algorithm: Mirror-Prox [14]
3. $\mathcal{X} = \mathbb{R}^p$ and $\mathcal{Y} = \mathbb{R}^n$; and $\Phi(\mathbf{x}, \mathbf{y}) = h(\mathbf{x}) + f(\mathbf{x}) + \langle \mathbf{A}\mathbf{x}, \mathbf{y} \rangle - g^*(\mathbf{y})$
 - ▶ Algorithms: Chambolle-Pock [4], Condat-Vu [5, 21], PD3O [23]

Smooth unconstrained minimax optimization

Details of the restricted minimax formulation

$$\min_{\mathbf{x} \in \mathbb{R}^d} \max_{\mathbf{y} \in \mathbb{R}^n} \Phi(\mathbf{x}, \mathbf{y}).$$

We assume that

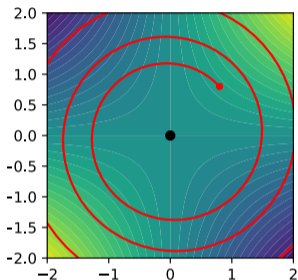
- ▶ $\Phi(\cdot, \mathbf{y})$ is convex for all $\mathbf{y} \in \mathbb{R}^n$,
- ▶ $\Phi(\mathbf{x}, \cdot)$ is concave for all $\mathbf{x} \in \mathbb{R}^d$,
- ▶ $\Phi(\mathbf{x}, \mathbf{y})$ is continuously differentiable in \mathbf{x} and \mathbf{y} ,
- ▶ Φ is smooth in the following sense.

$$\|\mathbf{V}(\mathbf{z}_1) - \mathbf{V}(\mathbf{z}_2)\| := \left\| \begin{bmatrix} \nabla_{\mathbf{x}} \Phi(\mathbf{x}_1, \mathbf{y}_1) \\ -\nabla_{\mathbf{y}} \Phi(\mathbf{x}_1, \mathbf{y}_1) \end{bmatrix} - \begin{bmatrix} \nabla_{\mathbf{x}} \Phi(\mathbf{x}_2, \mathbf{y}_2) \\ -\nabla_{\mathbf{y}} \Phi(\mathbf{x}_2, \mathbf{y}_2) \end{bmatrix} \right\| \leq L \left\| \begin{bmatrix} \mathbf{x}_1 - \mathbf{x}_2 \\ \mathbf{y}_1 - \mathbf{y}_2 \end{bmatrix} \right\|, \text{ where } \mathbf{z} = \begin{pmatrix} \mathbf{x} \\ \mathbf{y} \end{pmatrix} \quad (2)$$

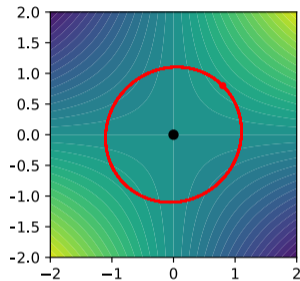
- Remarks:**
- GDA (i.e., $\mathbf{z}^{k+1} = \mathbf{z}^k - \tau \mathbf{V}(\mathbf{z}^k)$) **diverges** even for the simple bilinear objective (Lecture 11).
 - Roughly speaking, minimax is harder than just optimization (Lecture 11).

A running, bilinear example: $\min_{x \in \mathbb{R}} \max_{y \in \mathbb{R}} xy$

○ GDA



○ Alternating GDA



GDA

1. Choose $\mathbf{x}^0, \mathbf{y}^0$ and τ .
2. For $k = 0, 1, \dots$, perform:
 $\mathbf{x}^{k+1} := \mathbf{x}^k - \tau \nabla_{\mathbf{x}} \Phi(\mathbf{x}^k, \mathbf{y}^k)$.
 $\mathbf{y}^{k+1} := \mathbf{y}^k + \tau \nabla_{\mathbf{y}} \Phi(\mathbf{x}^k, \mathbf{y}^k)$.

AltGDA

1. Choose $\mathbf{x}^0, \mathbf{y}^0$ and α_k .
2. For $k = 0, 1, \dots$, perform:
 $\mathbf{x}^{k+1} := \mathbf{x}^k - \tau \nabla_{\mathbf{x}} \Phi(\mathbf{x}^k, \mathbf{y}^k)$.
 $\mathbf{y}^{k+1} := \mathbf{y}^k + \tau \nabla_{\mathbf{y}} \Phi(\mathbf{x}^{k+1}, \mathbf{y}^k)$.

A preview of algorithms to be covered

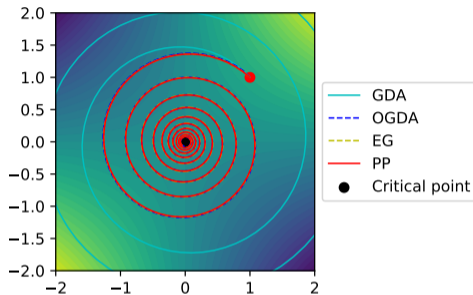


Figure: Trajectory of different algorithms for a simple bilinear game $\min_x \max_y xy$.

- Convergent algorithms in the sequel
 - ▶ Proximal point method (PPM)
 - ▶ Extra-gradient (EG)
 - ▶ Optimistic Gradient Descent Ascent (OGDA)

A preview of algorithms to be covered

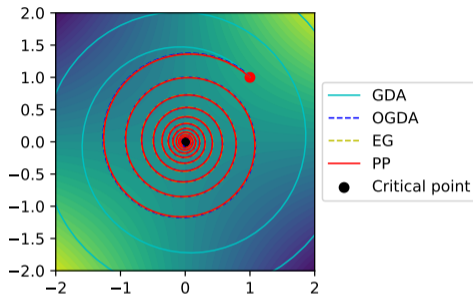


Figure: Trajectory of different algorithms for a simple bilinear game $\min_x \max_y xy$.

- Convergent algorithms in the sequel
 - ▶ Proximal point method (PPM)
 - ▶ Extra-gradient (EG)
 - ▶ Optimistic Gradient Descent Ascent (OGDA)
- EG and OGDA are approximations of the PPM [12]

Proximal point method (PPM)

- Consider following smooth unconstrained optimization problem:

$$\min_{\mathbf{x} \in \mathbb{R}^p} f(\mathbf{x})$$

Proximal point method for convex minimization.

For a step-size $\tau > 0$, PPM can be written as follows

$$\mathbf{x}^{k+1} = \arg \min_{\mathbf{x} \in \mathbb{R}^p} \left\{ f(\mathbf{x}) + \frac{1}{2\tau} \|\mathbf{x} - \mathbf{x}^k\|^2 \right\} := \text{prox}_{\tau f}(\mathbf{x}^k) \quad (3)$$

- Observations:**
- The optimality condition of (3) reveals a simpler PPM recursion for smooth f :

$$\mathbf{x}^{k+1} = \mathbf{x}^k - \tau \nabla f(\mathbf{x}^{k+1}).$$

- PPM is an **implicit**, non-practical algorithm since we need the point \mathbf{x}^{k+1} for its update.
- Each step of PPM can be as hard as solving the original problem.
- Convergence properties are well understood due to Rockafellar [18].

PPM and minimax optimization

PPM applied to the minimax template: $\min_{\mathbf{x} \in \mathbb{R}^d} \max_{\mathbf{y} \in \mathbb{R}^n} \Phi(\mathbf{x}, \mathbf{y})$

Define $\mathbf{z} = [\mathbf{x}, \mathbf{y}]^\top$ and $\mathbf{V}(\mathbf{z}) = [\nabla_{\mathbf{x}} \Phi(\mathbf{x}, \mathbf{y}), -\nabla_{\mathbf{y}} \Phi(\mathbf{x}, \mathbf{y})]^\top$. PPM iterations with a step-size $\tau > 0$ is given by

$$\mathbf{z}^{k+1} = \mathbf{z}^k - \tau \mathbf{V}(\mathbf{z}^{k+1}).$$

Derivation: o For $\tau > 0$, $(\mathbf{x}^{k+1}, \mathbf{y}^{k+1})$ is the unique solution to the saddle point problem,

$$\min_{\mathbf{x} \in \mathbb{R}^d} \max_{\mathbf{y} \in \mathbb{R}^n} \Phi(\mathbf{x}, \mathbf{y}) + \frac{1}{2\tau} \|\mathbf{x} - \mathbf{x}^k\|^2 - \frac{1}{2\tau} \|\mathbf{y} - \mathbf{y}^k\|^2 \quad (4)$$

o Writing the optimality condition of the update in (4)

$$\mathbf{x}^{k+1} = \mathbf{x}^k - \tau \nabla_{\mathbf{x}} \Phi(\mathbf{x}^{k+1}, \mathbf{y}^{k+1}), \quad \mathbf{y}^{k+1} = \mathbf{y}^k + \tau \nabla_{\mathbf{y}} \Phi(\mathbf{x}^{k+1}, \mathbf{y}^{k+1}) \quad (5)$$

Observation: o **PPM is an implicit algorithm.**

o For the bilinear problem, PPM is implementable!

PPM guarantees for minimax optimization

Theorem (Convergence of PPM [18])

Suppose $(\mathbf{x}^k, \mathbf{y}^k)$ be the iterates generated by PPM (i.e., (5)), then for the averaged iterates, it holds that

$$\left| \Phi \left(\frac{1}{K} \sum_{k=1}^K \mathbf{x}^k, \frac{1}{K} \sum_{k=1}^K \mathbf{y}^k \right) - \Phi(\mathbf{x}^*, \mathbf{y}^*) \right| \leq \frac{\|\mathbf{x}^0 - \mathbf{x}^*\|^2 + \|\mathbf{y}^0 - \mathbf{y}^*\|^2}{\tau K}.$$

Theorem (Linear convergence [18])

Suppose $(\mathbf{x}^k, \mathbf{y}^k)$ be the iterates generated by (5), $\Phi(\cdot, \cdot)$ is μ_x -strongly convex in \mathbf{x} and μ_y -strongly concave in \mathbf{y} . Let $\mu = \max\{\mu_x, \mu_y\}$. Then, for any $\tau > 0$, $(\mathbf{x}^k, \mathbf{y}^k)$ satisfies the following

$$r^{k+1} \leq \frac{1}{1 + \mu\tau} r^k,$$

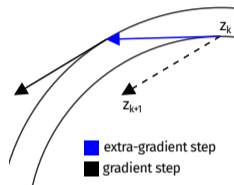
where $r^k = \|\mathbf{x}^k - \mathbf{x}^*\|^2 + \|\mathbf{y}^k - \mathbf{y}^*\|^2$.

- Remark:**
- Still need an implementable and convergent algorithm beyond the stylized bilinear case.
 - Note what happens when $\tau \rightarrow \infty$.

Extra-gradient algorithm (EG) [8]

EG method for saddle point problems

1. Choose $\mathbf{x}^0, \mathbf{y}^0$ and τ .
2. For $k = 0, 1, \dots$, perform:
 $\tilde{\mathbf{x}}^k := \mathbf{x}^k - \tau \nabla_{\mathbf{x}} \Phi(\mathbf{x}^k, \mathbf{y}^k),$
 $\tilde{\mathbf{y}}^k := \mathbf{y}^k + \tau \nabla_{\mathbf{y}} \Phi(\mathbf{x}^k, \mathbf{y}^k).$
 $\mathbf{x}^{k+1} := \mathbf{x}^k - \tau \nabla_{\mathbf{x}} \Phi(\tilde{\mathbf{x}}^k, \tilde{\mathbf{y}}^k).$
 $\mathbf{y}^{k+1} := \mathbf{y}^k + \tau \nabla_{\mathbf{y}} \Phi(\tilde{\mathbf{x}}^k, \tilde{\mathbf{y}}^k).$



- o Idea: Predict the gradient at the next point

$$\mathbf{z}^{k+1} = \mathbf{z}^k - \tau \mathbf{V} \left(\underbrace{\mathbf{z}^k - \tau \mathbf{V}(\mathbf{z}^k)}_{\text{prediction of } \mathbf{z}^{k+1}} \right)$$

(EG)

- Remark:**
- o 1-extra-gradient computation per iteration

Extra-gradient algorithm: Convergence

Theorem (General case [12])

Let $0 < \tau \leq \frac{1}{L}$. It holds that

- ▶ Iterates $(\mathbf{x}^k, \mathbf{y}^k)$ remains bounded in a convex compact set.
- ▶ Primal-dual gap reduces: $\text{Gap} \left(\frac{1}{K} \sum_{k=1}^K \mathbf{x}^k, \frac{1}{K} \sum_{k=1}^K \mathbf{y}^k \right) \leq \mathcal{O} \left(\frac{1}{K} \right)$.

Theorem (Linear convergence [13])

Suppose $(\mathbf{x}^k, \mathbf{y}^k)$ be the iterates generated by Extra-gradient algorithm, $\Phi(\cdot, \cdot)$ is μ_x -strongly convex in \mathbf{x} and μ_y -strongly concave in \mathbf{y} . Let $\mu = \max\{\mu_x, \mu_y\}$. Then, for $\tau = \frac{1}{4L}$, $(\mathbf{x}^k, \mathbf{y}^k)$ satisfies,

$$r^{k+1} \leq \left(1 - \frac{1}{c\kappa}\right)^k r^0,$$

where $r^k = \|\mathbf{x}^k - \mathbf{x}^*\|^2 + \|\mathbf{y}^k - \mathbf{y}^*\|^2$, $\kappa = \frac{L}{\mu}$ is the condition number of the problem, and c is a constant which is independent of the problem parameters.

Optimistic gradient descent algorithm (OGDA) [17]

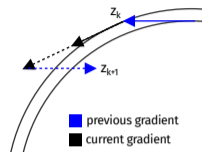
OGDA for saddle point problems

1. Choose $\mathbf{x}^0, \mathbf{y}^0, \mathbf{x}^1, \mathbf{y}^1$ and τ .

2. For $k = 1, \dots$, perform:

$$\mathbf{x}^{k+1} := \mathbf{x}^k - 2\tau \nabla_{\mathbf{x}} \Phi(\mathbf{x}^k, \mathbf{y}^k) + \tau \nabla_{\mathbf{x}} \Phi(\mathbf{x}^{k-1}, \mathbf{y}^{k-1}).$$

$$\mathbf{y}^{k+1} := \mathbf{y}^k + 2\tau \nabla_{\mathbf{y}} \Phi(\mathbf{x}^k, \mathbf{y}^k) - \tau \nabla_{\mathbf{y}} \Phi(\mathbf{x}^{k-1}, \mathbf{y}^{k-1}).$$



- o Main difference from the GDA: Add a “momentum” or “reflection” term to the updates

$$\mathbf{z}^{k+1} = \mathbf{z}^k - \tau \left[\mathbf{V}(\mathbf{z}^k) + \underbrace{(\mathbf{V}(\mathbf{z}^k) - \mathbf{V}(\mathbf{z}^{k-1}))}_{\text{momentum}} \right]. \quad (\text{OGDA})$$

- o Known as Popov’s method [16], it is also a special case of the Forward-Reflected-Backward method [11].
- o It has ties to the Reflected-Forward-Backward Splitting (RFBS) method [3]:

$$\mathbf{z}^{k+1} = \mathbf{z}^k - \tau \mathbf{V}(2\mathbf{z}^k - \mathbf{z}^{k-1}). \quad (\text{RFBS})$$

Remark:

- o Advanced material at the end: OGDA is an approximation of PPM for bilinear problems.

OGDA: Convergence

Theorem (General case [12])

Let $0 < \tau \leq \frac{1}{2L}$, $\mathbf{x}^1 = \mathbf{x}^0, \mathbf{y}^1 = \mathbf{y}^0$. It holds that

- ▶ Iterates $(\mathbf{x}^k, \mathbf{y}^k)$ remains bounded in a convex compact set.
- ▶ Primal-dual gap reduces: $\text{Gap} \left(\frac{1}{K} \sum_{k=1}^K \mathbf{x}^k, \frac{1}{K} \sum_{k=1}^K \mathbf{y}^k \right) \leq \mathcal{O} \left(\frac{1}{K} \right)$.

Theorem (Linear convergence [13])

Suppose $(\mathbf{x}^k, \mathbf{y}^k)$ be the iterates generated by OGDA, $\Phi(\cdot, \cdot)$ is μ_x -strongly convex in \mathbf{x} and μ_y -strongly concave in \mathbf{y} . Let $\mu = \max\{\mu_x, \mu_y\}$. Then, for $\tau = \frac{1}{4L}$, $(\mathbf{x}^k, \mathbf{y}^k)$ satisfies,

$$r^{k+1} \leq \left(1 - \frac{1}{c\kappa} \right)^k r^0,$$

where $r^k = \|\mathbf{x}^k - \mathbf{x}^*\|^2 + \|\mathbf{y}^k - \mathbf{y}^*\|^2$, $\kappa = \frac{L}{\mu}$ is the condition number of the problem, and c is a constant which is independent of the problem parameters.

A generalization of EG: The Mirror-Prox Algorithm

Definition: Bregman distance

Let $\omega : \mathcal{X} \times \mathcal{Y} \rightarrow \mathbb{R}$ be a distance generating function where ω is 1-strongly convex w.r.t. some norm $\|\cdot\|$ on the underlying space and is continuously differentiable. The Bregman distance induced by $\omega(\cdot)$ is given by

$$D_\omega(\mathbf{z}, \mathbf{z}') = \omega(\mathbf{z}) + \omega(\mathbf{z}') - \nabla\omega(\mathbf{z}')^\top (\mathbf{z} - \mathbf{z}').$$

Mirror-Prox algorithm

1. Choose $\mathbf{x}^0, \mathbf{y}^0$ and τ .

2. For $k = 0, 1, \dots$, perform:

$$\tilde{\mathbf{z}}^k = \arg \min_{\mathbf{z} \in \mathcal{X} \times \mathcal{Y}} (D_\omega(\mathbf{z}, \mathbf{z}^k) - \langle \tau \mathbf{V}(\mathbf{z}^k), \mathbf{z} \rangle).$$

$$\mathbf{z}^{k+1} = \arg \min_{\mathbf{z} \in \mathcal{X} \times \mathcal{Y}} (D_\omega(\mathbf{z}, \tilde{\mathbf{z}}^k) - \langle \tau \mathbf{V}(\tilde{\mathbf{z}}^k), \mathbf{z} \rangle).$$

Theorem (Mirror-Prox convergence)

Denote by $\Omega := \max_{\mathbf{z} \in \mathcal{X} \times \mathcal{Y}} D_\omega(\mathbf{z}, \mathbf{z}')$. The mirror-prox algorithm with $\tau \leq \frac{1}{L}$,

$$\text{Gap} \left(\frac{1}{K} \sum_{k=1}^K \mathbf{x}^k, \frac{1}{K} \sum_{k=1}^K \mathbf{y}^k \right) \leq \mathcal{O} \left(\frac{\Omega}{K} \right).$$

Comparison of convergence rates for **smooth** convex-concave minimax

Method	Assumption on $\Phi(\cdot, \cdot)$	Convergence rate	Reference	Note
PP	convex-concave	$\mathcal{O}(\epsilon^{-1})$	[18]	
PP	strongly convex- strongly concave	$\mathcal{O}(\kappa \log(\epsilon^{-1}))$	[18]	Implicit algorithm
PP	Bilinear	$\mathcal{O}(\kappa \log(\epsilon^{-1}))$	[18]	
EG	convex-concave	$\mathcal{O}(\epsilon^{-1})$	[12]	
EG	strongly convex- strongly concave	$\mathcal{O}(\kappa \log(\epsilon^{-1}))$	[13, 12]	1 extra-gradient evaluation per iteration
EG	Bilinear	$\mathcal{O}(\kappa \log(\epsilon^{-1}))$	[13, 12]	
OGDA	convex-concave	$\mathcal{O}(\epsilon^{-1})$	[12]	
OGDA	strongly convex- strongly concave	$\mathcal{O}(\kappa \log(\epsilon^{-1}))$	[13, 12]	no obvious downside
OGDA	Bilinear	$\mathcal{O}(\kappa \log(\epsilon^{-1}))$	[13, 12]	

Primal-dual methods for composite minimization: minimax reformulation

- Quest: Looking for algorithms such that $(\mathbf{x}^k, \mathbf{y}^k) \rightarrow (\mathbf{x}^*, \mathbf{y}^*)$ (with rates?)

Another restricted minimax template

$$\min_{\mathbf{x} \in \mathcal{X}} h(\mathbf{x}) + f(\mathbf{x}) + g(\mathbf{A}\mathbf{x}) = \min_{\mathbf{x} \in \mathcal{X}} \max_{\mathbf{y} \in \mathcal{Y}} \Phi(\mathbf{x}, \mathbf{y}) := h(\mathbf{x}) + f(\mathbf{x}) + \langle \mathbf{A}\mathbf{x}, \mathbf{y} \rangle - g^*(\mathbf{y}).$$

We assume that

- ▶ $f(\mathbf{x}) : \mathcal{X} \rightarrow \mathbb{R}$ is proper, convex and lower-semicontinuous (l.s.c.),
- ▶ $h(\mathbf{x}) : \mathcal{X} \rightarrow \mathbb{R}$ is proper, convex, l.s.c. and differentiable with a $\frac{1}{\beta}$ -Lipschitz continuous gradient,
- ▶ $g^*(\mathbf{y}) : \mathcal{Y} \rightarrow \mathbb{R}$ is proper, convex and l.s.c.
- ▶ $\mathcal{X} \subseteq \mathbb{R}^p$ and $\mathcal{Y} \subseteq \mathbb{R}^n$,
- ▶ $\mathbf{A} : \mathcal{X} \rightarrow \mathcal{Y}$ is a bounded linear operator,
- ▶ Problem has at least one solution $(\mathbf{x}^*, \mathbf{y}^*) \in \mathcal{X} \times \mathcal{Y}$

Primal-dual hybrid gradient method (PDHG, aka Chambolle-Pock)

$$\min_{\mathbf{x} \in \mathcal{X}} h(\mathbf{x}) + f(\mathbf{x}) + g(\mathbf{A}\mathbf{x}) = \min_{\mathbf{x} \in \mathcal{X}} \max_{\mathbf{y} \in \mathcal{Y}} \Phi(\mathbf{x}, \mathbf{y}) := h(\mathbf{x}) + f(\mathbf{x}) + \langle \mathbf{A}\mathbf{x}, \mathbf{y} \rangle - g^*(\mathbf{y})$$

PDHG [4], ($h(\mathbf{x}) = 0$)

1. Choose $\tilde{\mathbf{x}}^0, \mathbf{x}^0, \mathbf{y}^0$ and $\tau, \sigma > 0$.

2. For $k = 0, 1, \dots$, perform:

$$\mathbf{y}^{k+1} = \text{prox}_{\sigma g^*}(\mathbf{y}^k + \sigma \mathbf{A}\tilde{\mathbf{x}}^k).$$

$$\mathbf{x}^{k+1} = \text{prox}_{\tau f}(\mathbf{x}^k - \tau \mathbf{A}^T \mathbf{y}^{k+1}).$$

$$\tilde{\mathbf{x}}^{k+1} = 2\mathbf{x}^{k+1} - \mathbf{x}^k.$$

Theorem ([4])

Let $L = \|\mathbf{A}\|$, and choose τ and σ such that we have $\tau\sigma L^2 < 1$. Then, it holds that

- ▶ Iterates $(\mathbf{x}^k, \mathbf{y}^k)$ remains bounded in a convex compact set.
- ▶ Primal-dual gap satisfies $\text{Gap} \left(\frac{1}{K} \sum_{k=1}^K \mathbf{x}^k, \frac{1}{K} \sum_{k=1}^K \mathbf{y}^k \right) \leq \mathcal{O} \left(\frac{1}{K} \right)$.
- ▶ $(\mathbf{x}^k, \mathbf{y}^k)$ converges to saddle point $(\mathbf{x}^*, \mathbf{y}^*)$.
- ▶ If f and g are smooth, the rate improves to $\mathcal{O}(1/K^2)$.
- ▶ If f and g are also strongly convex, the convergence is linear.

Stochastic PDHG

$$\min_{\mathbf{x} \in \mathcal{X}} f(\mathbf{x}) + \sum_{i=1}^n g_i(\mathbf{A}_i \mathbf{x}) = \min_{\mathbf{x} \in \mathcal{X}} \max_{\mathbf{y} \in \mathcal{Y}} \underbrace{\Phi(\mathbf{x}, \mathbf{y})}_{=0} := h(\mathbf{x}) + f(\mathbf{x}) + \sum_{i=1}^n \langle \mathbf{A}_i \mathbf{x}, \mathbf{y}_i \rangle - \sum_{i=1}^n g_i^*(\mathbf{y}_i) \quad (6)$$

Algorithm 1 Stochastic Primal-Dual Hybrid Gradient

Input: Pick step sizes σ_i, τ and $\mathbf{x}^0 \in \mathcal{X}, \mathbf{y}^0 = \mathbf{y}^1 = \bar{\mathbf{y}}^1 \in \mathcal{Y}$. Given $\mathbf{P} = \text{diag}(p_1, \dots, p_n)$.

for $k = 1, 2, \dots$ **do**

$$\mathbf{x}^k = \text{prox}_{\tau f}(\mathbf{x}^{k-1} - \tau \sum_i \mathbf{A}_i^\top \bar{\mathbf{y}}_i^k)$$

Draw $j_k \in \{1, \dots, n\}$ such that $\mathbb{P}(j_k = j) = p_j$.

$$\mathbf{y}_{j_k}^{k+1} = \text{prox}_{\sigma_{j_k} g_{j_k}^*}(\mathbf{y}_{j_k}^k + \sigma_{j_k} \mathbf{A}_{j_k} \mathbf{x}^k)$$

$$\mathbf{y}_j^{k+1} = \mathbf{y}_j^k, \forall j \neq j_k$$

$$\bar{\mathbf{y}}_i^{k+1} = \mathbf{y}_i^{k+1} + \mathbf{P}^{-1}(\mathbf{y}_i^{k+1} - \mathbf{y}_i^k), \forall i,$$

end for

Remarks:

- Note: $p_i^{-1} \tau \sigma_i \|\mathbf{A}_i\|^2 < 1$.
- Only one dual vector is updated at each iteration.
- Especially effective when \mathbf{A}_i is row-vector.

SPDHG: Convergence [1]

Theorem (Almost sure convergence)

Almost surely, there exists $(\mathbf{x}^*, \mathbf{y}^*) \in \mathcal{Z}^*$, such that the iterates of SPDHG satisfy $\mathbf{x}^k \rightarrow \mathbf{x}^*$ and $\mathbf{y}^k \rightarrow \mathbf{y}^*$.

Theorem (Sublinear convergence)

Define the ergodic sequences $\mathbf{x}_{av}^K = \sum_{k=1}^K \mathbf{x}^k$ and $\mathbf{y}_{av}^{K+1} = \sum_{k=1}^K \mathbf{y}^{k+1}$, and define the gap function

$$\text{Gap}(\mathbf{x}_{av}^K, \mathbf{y}_{av}^{K+1}) = \sup_{\mathbf{x}, \mathbf{y}} f(\mathbf{x}_{av}^K) + \langle A\mathbf{x}_{av}^K, \mathbf{y} \rangle - g^*(\mathbf{y}) - f(\mathbf{x}) - \langle A\mathbf{x}, \mathbf{y}_{av}^{K+1} \rangle + g^*(\mathbf{y}_{av}^{K+1}).$$

The following result holds for the expected primal-dual gap, which is expectation of a supremum

$$\mathbb{E} \left[\text{Gap}(\mathbf{x}_{av}^K, \mathbf{y}_{av}^{K+1}) \right] = \mathcal{O} \left(\frac{1}{K} \right). \quad (7)$$

Primal-dual algorithms for minimax: The zoo!

$$\min_{\mathbf{x} \in \mathcal{X}} h(\mathbf{x}) + f(\mathbf{x}) + g(\mathbf{Ax}) = \min_{\mathbf{x} \in \mathcal{X}} \max_{\mathbf{y} \in \mathcal{Y}} \Phi(\mathbf{x}, \mathbf{y}) := h(\mathbf{x}) + f(\mathbf{x}) + \langle \mathbf{Ax}, \mathbf{y} \rangle - g^*(\mathbf{y})$$

3 operator splitting [6], ($\mathbf{A} = \mathbb{I}$)

1. Choose $\hat{\mathbf{x}}^0, \mathbf{x}^0, \mathbf{y}^0$ and $\tau > 0$.

2. For $k = 0, 1, \dots$, perform:

$$\mathbf{x}^{k+1} = \text{prox}_{\tau f}(\tilde{\mathbf{x}}^k).$$

$$\mathbf{y}^{k+1} = \frac{1}{\tau} (\mathbb{I} + \text{prox}_{\tau^{-1}g}) (2\mathbf{x}^{k+1} - \tilde{\mathbf{x}}^k - \tau \nabla h(\mathbf{x}^{k+1})).$$

$$\tilde{\mathbf{x}}^{k+1} = \mathbf{x}^{k+1} - \tau \nabla h(\mathbf{x}^{k+1}) - \tau \mathbf{y}^{k+1}.$$

Primal-dual algorithms for minimax: The zoo!

$$\min_{\mathbf{x} \in \mathcal{X}} h(\mathbf{x}) + f(\mathbf{x}) + g(\mathbf{Ax}) = \min_{\mathbf{x} \in \mathcal{X}} \max_{\mathbf{y} \in \mathcal{Y}} \Phi(\mathbf{x}, \mathbf{y}) := h(\mathbf{x}) + f(\mathbf{x}) + \langle \mathbf{Ax}, \mathbf{y} \rangle - g^*(\mathbf{y})$$

3 operator splitting [6], ($\mathbf{A} = \mathbb{I}$)

1. Choose $\hat{\mathbf{x}}^0, \mathbf{x}^0, \mathbf{y}^0$ and $\tau > 0$.

2. For $k = 0, 1, \dots$, perform:

$$\mathbf{x}^{k+1} = \text{prox}_{\tau f}(\tilde{\mathbf{x}}^k).$$

$$\mathbf{y}^{k+1} = \frac{1}{\tau} (\mathbb{I} + \text{prox}_{\tau^{-1}g}) (2\mathbf{x}^{k+1} - \tilde{\mathbf{x}}^k - \tau \nabla h(\mathbf{x}^{k+1})).$$

$$\tilde{\mathbf{x}}^{k+1} = \mathbf{x}^{k+1} - \tau \nabla h(\mathbf{x}^{k+1}) - \tau \mathbf{y}^{k+1}.$$

o There is a stochastic variant [24].

Primal-dual algorithms for minimax: The zoo!

$$\min_{\mathbf{x} \in \mathcal{X}} h(\mathbf{x}) + f(\mathbf{x}) + g(\mathbf{Ax}) = \min_{\mathbf{x} \in \mathcal{X}} \max_{\mathbf{y} \in \mathcal{Y}} \Phi(\mathbf{x}, \mathbf{y}) := h(\mathbf{x}) + f(\mathbf{x}) + \langle \mathbf{Ax}, \mathbf{y} \rangle - g^*(\mathbf{y})$$

Condat-Vu [5, 21]

1. Choose $\hat{\mathbf{x}}^0, \mathbf{x}^0, \mathbf{y}^0$ and $\tau, \sigma > 0$.

2. For $k = 0, 1, \dots$, perform:

$$\mathbf{y}^{k+1} = \text{prox}_{\sigma g^*}(\mathbf{y}^k + \sigma \mathbf{A} \tilde{\mathbf{x}}^k).$$

$$\mathbf{x}^{k+1} = \text{prox}_{\tau f}(\mathbf{x}^k - \tau \nabla h(\mathbf{x}^k) - \tau \mathbf{A}^T \mathbf{y}^{k+1}).$$

$$\tilde{\mathbf{x}}^{k+1} = 2\mathbf{x}^{k+1} - \mathbf{x}^k.$$

Primal-dual algorithms for minimax: The zoo!

$$\min_{\mathbf{x} \in \mathcal{X}} h(\mathbf{x}) + f(\mathbf{x}) + g(\mathbf{Ax}) = \min_{\mathbf{x} \in \mathcal{X}} \max_{\mathbf{y} \in \mathcal{Y}} \Phi(\mathbf{x}, \mathbf{y}) := h(\mathbf{x}) + f(\mathbf{x}) + \langle \mathbf{Ax}, \mathbf{y} \rangle - g^*(\mathbf{y})$$

PD3O splitting [23]

1. Choose $\hat{\mathbf{x}}^0, \mathbf{x}^0, \mathbf{y}^0$ and $\tau, \sigma > 0$.

2. For $k = 0, 1, \dots$, perform:

$$\mathbf{y}^{k+1} = \text{prox}_{\sigma g^*}(\mathbf{y}^k + \sigma \mathbf{A} \tilde{\mathbf{x}}^k).$$

$$\mathbf{x}^{k+1} = \text{prox}_{\tau f}(\mathbf{x}^k - \tau \nabla h(\mathbf{x}^k) - \tau \mathbf{A}^T \mathbf{y}^{k+1}).$$

$$\tilde{\mathbf{x}}^{k+1} = 2\mathbf{x}^{k+1} - \mathbf{x}^k + \tau \nabla h(\mathbf{x}^k) - \tau \nabla h(\mathbf{x}^{k+1}).$$

Between convex-concave and nonconvex-nonconcave

Nonconvex-concave problems

$$\min_{\mathbf{x} \in \mathcal{X}} \max_{\mathbf{y} \in \mathcal{Y}} \Phi(\mathbf{x}, \mathbf{y})$$

- $\Phi(\mathbf{x}, \mathbf{y})$ is nonconvex in \mathbf{x} , concave in \mathbf{y} , smooth in \mathbf{x} and \mathbf{y} .

Recall

Define $f(\mathbf{x}) = \max_{\mathbf{y} \in \mathcal{Y}} \Phi(\mathbf{x}, \mathbf{y})$.

- Gradient descent applied to nonconvex f requires $\mathcal{O}(\epsilon^{-2})$ iterations to give an ϵ -stationary point.
- (Sub)gradient of f can be computed using Danskin's theorem:

$$\nabla_{\mathbf{x}} \Phi(\cdot, \mathbf{y}^*(\cdot)) \in \partial f(\cdot), \text{ where } \mathbf{y}^*(\cdot) \in \arg \max_{\mathbf{y} \in \mathcal{Y}} \Phi(\cdot, \mathbf{y}),$$

which is **tractable since Φ is concave in \mathbf{y} [9]**.

- Remark:**
- “Conceptually” much easier than nonconvex-nonconcave case.

A summary of results for **nonconvex-concave** setting

- o A summary of gradient complexities to reach ϵ —first order stationary point in terms of **gradient mapping**.

Method	Assumption on $\Phi(\cdot, \cdot)$	Convergence rate	Reference
GDA	nonconvex-concave	$\tilde{\mathcal{O}}(\epsilon^{-6})$	[9]
GDA	nonconvex- strongly concave	$\mathcal{O}(\epsilon^{-2})$	[9]
GDmax	nonconvex-concave	$\tilde{\mathcal{O}}(\epsilon^{-6})$	[7]
GDmax	nonconvex- strongly concave	$\mathcal{O}(\epsilon^{-2})$	[7]
HiBSA, AGP	nonconvex-concave	$\tilde{\mathcal{O}}(\epsilon^{-4})$	[10], [22]
HiBSA, AGP	nonconvex- strongly concave	$\mathcal{O}(\epsilon^{-2})$	[10], [22]

The elephant in the room: Nonsmooth, nonconvex optimization

$$\min_{\mathbf{x} \in \mathbb{R}^p} f(\mathbf{x})$$

- Finding a stationary point of nonsmooth nonconvex minimization problems are hard [25]
 - ▶ A traditional ϵ -stationarity can not be obtained in finite time
- Even the relax notions are hard [19]
- Really puzzling how deep learning approaches with ReLu etc. work...

How about purely primal approaches?

$$\min_{\mathbf{x} \in \mathbb{R}^p} \left\{ f(\mathbf{x}) : \mathbf{Ax} = \mathbf{b} \right\}$$

Penalty methods

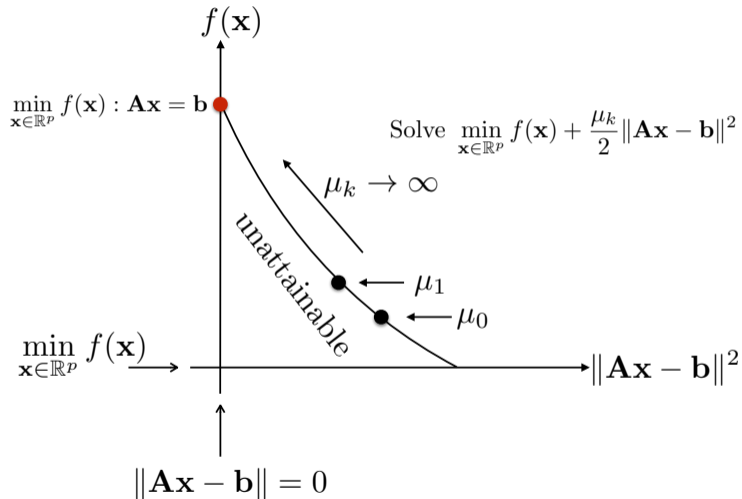
- Convert constrained problem (**difficult**) to unconstrained (**easy**).
- Penalized function with penalty parameter $\mu > 0$:

$$F_\mu(\mathbf{x}) := \left\{ f(\mathbf{x}) + \frac{\mu}{2} \|\mathbf{Ax} - \mathbf{b}\|^2 \right\} \xleftrightarrow{\mu \rightarrow \infty} \min_{\mathbf{x} \in \mathbb{R}^p} \left\{ f(\mathbf{x}) : \mathbf{Ax} = \mathbf{b} \right\}.$$

○ Observations:

- ▶ Minimize a weighted combination of $f(\mathbf{x})$ and $\|\mathbf{Ax} - \mathbf{b}\|^2$ at the same time.
- ▶ μ determines the weight of $\|\mathbf{Ax} - \mathbf{b}\|^2$.
- ▶ As $\mu \rightarrow \infty$, we enforce $\mathbf{Ax} = \mathbf{b}$.
- ▶ Other functions than the **quadratic** $\frac{1}{2} \|\cdot\|^2$ are also possible e.g., exact nonsmooth penalty functions:
 - ▶ $\mu \|\mathbf{Ax} - \mathbf{b}\|_2$ or $\mu \|\mathbf{Ax} - \mathbf{b}\|_1$
 - ▶ They work with finite μ , but they are difficult to solve [15, Section 17.2], [2]

Quadratic penalty: Intuition



Quadratic penalty: Conceptual algorithm

Quadratic penalty method (QP):

1. Choose $\mathbf{x}_0 \in \mathbb{R}^p$ and $\mu_0 > 0$.
2. For $k = 0, 1, \dots$, perform:
 - 2.a. $\mathbf{x}_k := \arg \min_{\mathbf{x} \in \mathbb{R}^p} \left\{ f(\mathbf{x}) + \frac{\mu_k}{2} \|\mathbf{Ax} - \mathbf{b}\|^2 \right\}$.
 - 2.b. Update $\mu_{k+1} > \mu_k$.

Theorem [15, Theorem 17.1]

Assume that f is smooth and $\mu_k \rightarrow \infty$. Then, every limit point $\bar{\mathbf{x}}$ of the sequence $\{\mathbf{x}_k\}$ is a solution of the constrained problem

$$\mathbf{x}^* \in \arg \min_{\mathbf{x} \in \mathbb{R}^p} \left\{ f(\mathbf{x}) : \mathbf{Ax} = \mathbf{b} \right\}.$$

Limitation

- The minimization problems of step 2.a. of the algorithm become ill-conditioned as $\mu_k \rightarrow \infty$.
- Common improvements:
 - ▶ Solve the subproblem inexactly, *i.e.*, up to ϵ accuracy.
 - ▶ **Linearization** to simplify subproblems (**up next**).

Quadratic penalty: Linearization

Generalized quadratic penalty method:

1. Choose $\mathbf{x}_0 \in \mathbb{R}^p$, $\mu_0 > 0$ and positive semidefinite matrix \mathbf{Q}_k .

2. For $k = 0, 1, \dots$, perform:

2.a. $\mathbf{x}_k := \arg \min_{\mathbf{x} \in \mathbb{R}^p} \left\{ f(\mathbf{x}) + \frac{\mu_k}{2} \|\mathbf{A}\mathbf{x} - \mathbf{b}\|^2 + \frac{1}{2} \|\mathbf{x} - \mathbf{x}_{k-1}\|_{\mathbf{Q}_k}^2 \right\}$.

2.b. Update $\mu_{k+1} > \mu_k$.

Ideas

- Minimize a **majorizer** of $F_\mu(\mathbf{x})$, parametrized by \mathbf{Q}_k in step 2.a..
- $\mathbf{Q}_k = \mathbf{0}$ gives the standard QP; $\mathbf{Q}_k = \mathbf{I}$ gives strongly convex subproblems.
- $\mathbf{Q}_k = \alpha_k \mathbf{I} - \mu_k \mathbf{A}^\top \mathbf{A}$, with $\alpha_k \geq \mu_k \|\mathbf{A}\|^2$ gives

$$\mathbf{x}_k = \text{prox}_{\frac{1}{\alpha_k} f} \left(\mathbf{x}_{k-1} - \frac{\mu_k}{\alpha_k} \mathbf{A}^\top (\mathbf{A}\mathbf{x}_{k-1} - \mathbf{b}) \right) \quad \text{Only one proximal operator!}$$

and picking $\alpha_k = \mu_k \|\mathbf{A}\|^2$ gives

$$\mathbf{x}_k = \text{prox}_{\frac{1}{\mu_k \|\mathbf{A}\|^2} f} \left(\mathbf{x}_{k-1} - \frac{1}{\|\mathbf{A}\|^2} \mathbf{A}^\top (\mathbf{A}\mathbf{x}_{k-1} - \mathbf{b}) \right).$$

Per-iteration time: The key role of the prox-operator

Recall: Prox-operator

$$\text{prox}_f(\mathbf{x}) := \arg \min_{\mathbf{z} \in \mathbb{R}^p} \left\{ f(\mathbf{z}) + \frac{1}{2} \|\mathbf{z} - \mathbf{x}\|^2 \right\}.$$

Key properties:

- ▶ **single valued & non-expansive** since f is a proper convex function.
- ▶ **distributes** when the primal problem has **decomposable** structure:

$$f(\mathbf{x}) := \sum_{i=1}^m f_i(\mathbf{x}_i), \quad \text{and} \quad \mathcal{X} := \mathcal{X}_1 \times \cdots \times \mathcal{X}_m.$$

where $m \geq 1$ is the **number of components**.

- ▶ **often efficient & has closed form expression**. For instance, if $f(\mathbf{z}) = \|\mathbf{z}\|_1$, then the prox-operator performs coordinate-wise soft-thresholding by 1.

Quadratic penalty: Linearized methods

Linearized QP method (LQP)

1. Choose $\mathbf{x}_0 \in \mathbb{R}^p$, $\sigma_0 = 1$, $\mu_0 > 0$.

2. For $k = 0, 1, \dots$:

2.a. $\mathbf{x}_{k+1} := \text{prox}_{\frac{1}{\mu_k \|\mathbf{A}\|^2} f} \left(\mathbf{x}_k - \frac{1}{\|\mathbf{A}\|^2} \mathbf{A}^\top (\mathbf{A} \mathbf{x}_k - \mathbf{b}) \right)$.

2.b. Update σ_{k+1} s.t. $\frac{(1 - \sigma_{k+1})^2}{\sigma_{k+1}} = \frac{1}{\sigma_k}$.

2.c. Update $\mu_{k+1} = \sqrt{\sigma_{k+1}}$.

Accelerated linearized QP method (ALQP)

1. Choose $\mathbf{x}_0, \mathbf{y}_0 \in \mathbb{R}^p$, $\tau_0 = 1$, $\mu_0 > 0$.

2. For $k = 0, 1, \dots$:

2.a. $\mathbf{x}_{k+1} := \text{prox}_{\frac{1}{\mu_k \|\mathbf{A}\|^2} f} \left(\mathbf{y}_k - \frac{1}{\|\mathbf{A}\|^2} \mathbf{A}^\top (\mathbf{A} \mathbf{y}_k - \mathbf{b}) \right)$.

2.b. $\mathbf{y}_{k+1} := \mathbf{x}_{k+1} + \frac{\tau_{k+1}(1 - \tau_k)}{\tau_k} (\mathbf{x}_{k+1} - \mathbf{x}_k)$.

2.c. Update $\mu_{k+1} = \mu_k (1 + \tau_{k+1})$.

2.d. Update $\tau_{k+1} \in (0, 1)$ as the unique positive root of $\tau^3 + \tau^2 + \tau_k^2 \tau - \tau_k^2 = 0$.

Theorem (Convergence [20])

o **LQP:**

$$|f(\mathbf{x}_k) - f(\mathbf{x}^*)| \leq \mathcal{O}(\mu_0 k^{-1/2} + \mu_0^{-1} k^{-1/2})$$

$$\|\mathbf{A} \mathbf{x}_k - \mathbf{b}\| \leq \mathcal{O}(\mu_0^{-1} k^{-1/2})$$

o **ALQP:**

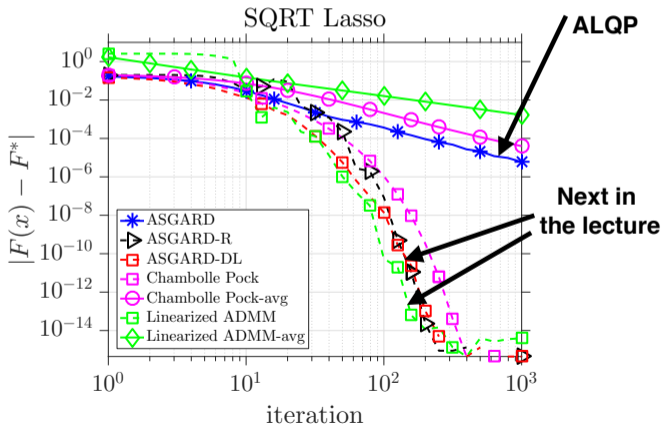
$$|f(\mathbf{x}_k) - f(\mathbf{x}^*)| \leq \mathcal{O}(\mu_0 k^{-1} + \mu_0^{-1} k^{-1})$$

$$\|\mathbf{A} \mathbf{x}_k - \mathbf{b}\| \leq \mathcal{O}(\mu_0^{-1} k^{-1})$$

In practice: **poor (worst case) performance**

- A nonsmooth problem: SQR T Lasso

$$\min_{\mathbf{x} \in \mathbb{R}^p} \|\mathbf{A}\mathbf{x} - \mathbf{b}\|_2 + \lambda \|\mathbf{x}\|_1.$$



Wrap up!

- Recitation continues with Homework #2 on Friday...

*OGDA as an approximation of PPM

Claim: OGDA is an approximation of PPM.

○ Consider the bilinear case $\Phi(\mathbf{x}, \mathbf{y}) = \langle \mathbf{x}, \mathbf{B}\mathbf{y} \rangle$, where $\mathbf{B} \in \mathbb{R}^{p \times p}$ is a square full rank matrix. The point $(\mathbf{x}^*, \mathbf{y}^*) = (\mathbf{0}, \mathbf{0})$ is a unique saddle point.

○ OGDA updates are

$$\mathbf{x}^{k+1} = \mathbf{x}^k - 2\tau\mathbf{B}\mathbf{y}^k + \tau\mathbf{B}\mathbf{y}^{k-1}, \quad \mathbf{y}^{k+1} = \mathbf{y}^k + 2\tau\mathbf{B}^\top\mathbf{x}^k - \tau\mathbf{B}^\top\mathbf{x}^{k-1}$$

○ From (5), PP update on the variable \mathbf{x} is

$$\mathbf{x}^{k+1} = \mathbf{x}^k - \tau\mathbf{B}\mathbf{y}^{k+1} = \mathbf{x}^k - \tau\mathbf{B}(\mathbf{y}^k + \tau\mathbf{B}^\top\mathbf{x}^{k+1}),$$

where we used $\mathbf{y}^{k+1} = \mathbf{y}^k + \tau\mathbf{B}^\top\mathbf{x}^{k+1}$. So, PP method update on the variable \mathbf{x} can be rewritten as

$$\mathbf{x}^{k+1} = (\mathbb{I} + \tau^2\mathbf{B}\mathbf{B}^\top)^{-1}(\mathbf{x}^k - \tau\mathbf{B}\mathbf{y}^k)$$

○ Use the fact that $(\mathbb{I} - \tau^2\mathbf{B}\mathbf{B}^\top)$ is an approximation $(\mathbb{I} + \tau^2\mathbf{B}\mathbf{B}^\top)^{-1}$ with an error $o(\tau^2)$.

$$(\mathbb{I} + \tau^2\mathbf{B}\mathbf{B}^\top)^{-1} = (\mathbb{I} - \tau^2\mathbf{B}\mathbf{B}^\top + o(\tau^2)) \tag{8}$$

*OGDA as an approximation of PPM

- Using (8), rewrite the update on \mathbf{x} for PPM as

$$\mathbf{x}^{k+1} = \left(\mathbb{I} - \tau^2 \mathbf{B} \mathbf{B}^\top + o(\tau^2) \right) (\mathbf{x}^k - \tau \mathbf{B} \mathbf{y}^k)$$

- Adding and subtracting $\mathbf{B} \mathbf{y}^k$ to the right hand side, using the PP updates and reorganizing the terms

$$\begin{aligned} \mathbf{x}^{k+1} &= \mathbf{x}^k - \tau \mathbf{B} \mathbf{y}^k - \tau \mathbf{B} \left(\tau \mathbf{B}^\top \mathbf{x}^k - \tau^2 \mathbf{B}^\top \mathbf{B} \mathbf{y}^k \right) + o(\tau^2) \\ &= \mathbf{x}^k - 2\tau \mathbf{B} \mathbf{y}^k - \tau \mathbf{B} \left(\tau \mathbf{B}^\top \mathbf{x}^k - (\mathbb{I} + \tau^2 \mathbf{B}^\top \mathbf{B}) \mathbf{y}^k \right) + o(\tau^2) \\ &= \mathbf{x}^k - 2\tau \mathbf{B} \mathbf{y}^k - \tau \mathbf{B} \left(\tau \mathbf{B}^\top \mathbf{x}^k - \mathbf{y}^{k-1} - \tau \mathbf{B}^\top \mathbf{x}^{k-1} \right) + o(\tau^2) \\ &= \mathbf{x}^k - 2\tau \mathbf{B} \mathbf{y}^k - \tau \mathbf{B} \mathbf{y}^{k-1} + o(\tau^2) \end{aligned}$$

- The last equation is OGDA update for variable \mathbf{x} plus an additional error of $o(\tau^2)$. Similarly for variable \mathbf{y} .

*OGDA as an approximation of PPM

- Using (8), rewrite the update on \mathbf{x} for PPM as

$$\mathbf{x}^{k+1} = (\mathbb{I} - \tau^2 \mathbf{B} \mathbf{B}^\top + o(\tau^2)) (\mathbf{x}^k - \tau \mathbf{B} \mathbf{y}^k)$$

- Adding and subtracting $\mathbf{B} \mathbf{y}^k$ to the right hand side, using the PP updates and reorganizing the terms

$$\begin{aligned} \mathbf{x}^{k+1} &= \mathbf{x}^k - \tau \mathbf{B} \mathbf{y}^k - \tau \mathbf{B} (\tau \mathbf{B}^\top \mathbf{x}^k - \tau^2 \mathbf{B}^\top \mathbf{B} \mathbf{y}^k) + o(\tau^2) \\ &= \mathbf{x}^k - 2\tau \mathbf{B} \mathbf{y}^k - \tau \mathbf{B} (\tau \mathbf{B}^\top \mathbf{x}^k - (\mathbb{I} + \tau^2 \mathbf{B}^\top \mathbf{B}) \mathbf{y}^k) + o(\tau^2) \\ &= \mathbf{x}^k - 2\tau \mathbf{B} \mathbf{y}^k - \tau \mathbf{B} (\tau \mathbf{B}^\top \mathbf{x}^k - \mathbf{y}^{k-1} - \tau \mathbf{B}^\top \mathbf{x}^{k-1}) + o(\tau^2) \\ &= \mathbf{x}^k - 2\tau \mathbf{B} \mathbf{y}^k - \tau \mathbf{B} \mathbf{y}^{k-1} + o(\tau^2) \end{aligned}$$

- The last equation is OGDA update for variable \mathbf{x} plus an additional error of $o(\tau^2)$. Similarly for variable \mathbf{y} .

Proposition

Given a point $(\mathbf{x}^k, \mathbf{y}^k)$, let $(\hat{\mathbf{x}}^{k+1}, \hat{\mathbf{y}}^{k+1})$ be the point generated by performing a PP update on $(\mathbf{x}^k, \mathbf{y}^k)$, and let $(\mathbf{x}^{k+1}, \mathbf{y}^{k+1})$ be the point generated by performing an OGDA update on $(\mathbf{x}^k, \mathbf{y}^k)$. For $\eta > 0$

$$\|\mathbf{x}^{k+1} - \hat{\mathbf{x}}^{k+1}\| \leq o(\tau^2), \quad \|\mathbf{y}^{k+1} - \hat{\mathbf{y}}^{k+1}\| \leq o(\tau^2).$$

*Tools for the algorithms: resolvent operator and prox-mapping

- o We need to solve problems of type (9) at each iteration.

$$\boxed{\mathbf{x}^+ = \arg \min_{\mathbf{x}} \left\{ f(\mathbf{x}) + \frac{\|\mathbf{x} - \mathbf{y}\|^2}{2\tau} \right\}} := \text{prox}_{\tau f}(\mathbf{y}) \quad (9)$$

- o Writing the optimality condition gives

$$0 \in \partial f(\mathbf{x}^+) + \frac{1}{\tau}(\mathbf{x}^+ - \mathbf{y}) \quad \Rightarrow \quad \mathbf{x}^+ = \underbrace{(\mathbb{I} + \tau \partial f)^{-1}}_{\text{resolvent operator}}(\mathbf{y}), \quad (10)$$

where ∂f is the subgradient of convex function f and \mathbb{I} is the identity operator.

- o We assume resolvent operator defined through (10) is either
 - ▶ have a closed form solution, or
 - ▶ can be efficiently solved.

*Tools for the algorithms: Moreau's identity

- Similarly, for the dual parameter update, we need the following proximal operator of g^* .

$$\mathbf{y}^+ = \text{prox}_{\sigma g^*}(\mathbf{x})$$

- A fundamental equality for the prox operator: Moreau's identity

$$\mathbf{x} = \text{prox}_g(\mathbf{x}) + \text{prox}_{g^*}(\mathbf{x}) \quad (\text{Moreau's Identity})$$

- It is easy to compute $\text{prox}_{\sigma g^*}(\mathbf{x})$ by using the proximal mapping of function g as

$$\text{prox}_{\sigma g^*}(\mathbf{x}) = \mathbf{x} - \sigma \text{prox}_{\sigma^{-1}g}\left(\frac{\mathbf{x}}{\sigma}\right) \quad (\text{Extended Moreau's Identity})$$

*Extended Moreau's identity

$$\text{prox}_{\sigma g^*}(\mathbf{x}) = \mathbf{x} - \sigma \text{prox}_{\sigma^{-1}g}\left(\frac{\mathbf{x}}{\sigma}\right)$$

Proof: Extended Moreau's identity

First prove that Moreau's identity holds: $\mathbf{x} = \text{prox}_g(\mathbf{x}) + \text{prox}_{g^*}(\mathbf{x})$

$$\mathbf{y} = \text{prox}_g(\mathbf{x}) \iff \mathbf{x} - \mathbf{y} \in \partial g(\mathbf{y}) \quad (\text{Optimality of prox})$$

$$\iff \mathbf{y} \in \partial g^*(\mathbf{x} - \mathbf{y}) \quad (\text{Conjugate subgradient theorem})$$

$$\iff \mathbf{x} - \mathbf{y} = \text{prox}_{g^*}(\mathbf{x})$$

$$\iff \mathbf{x} = \text{prox}_g(\mathbf{x}) + \text{prox}_{g^*}(\mathbf{x}) \quad (\mathbf{y} = \text{prox}_g(\mathbf{x}))$$

Now applying Moreau's identity to function σg

$$\mathbf{x} = \text{prox}_{\sigma g}(\mathbf{x}) + \text{prox}_{(\sigma g)^*}(\mathbf{x})$$

$$= \text{prox}_{\sigma g}(\mathbf{x}) + \sigma \text{prox}_{\sigma^{-1}g^*}\left(\frac{\mathbf{x}}{\sigma}\right) \quad ((\sigma g)^*(\mathbf{y}) = \sigma g^*\left(\frac{\mathbf{y}}{\sigma}\right))$$

*Primal-dual with **random extrapolation** and coordinate descent: PURE-CD

Input: $\mathbf{x}_0 \in \mathbb{R}^n$, $\mathbf{y}_0 \in \mathbb{R}^m$

Parameters: $\theta = \text{diag}(\theta_1, \dots, \theta_m)$ is chosen as $\theta_j = \frac{\pi_j}{\underline{p}}$, where $\pi_j = \sum_{i \in I(j)} p_i$, and $\underline{p} = \min_i p_i$, and

$$\tau_i < \frac{2p_i - \underline{p}}{\beta_i p_i + \underline{p}^{-1} p_i \sum_{j=1}^m \pi_j \sigma_j A_{j,i}^2} 1.$$

for $k \in \mathbb{N}$ **do**

$$\bar{\mathbf{y}}_{k+1} = \text{prox}_{\sigma g^*}(\mathbf{y}_k + \sigma \mathbf{A} \mathbf{x}_k)$$

$$\bar{\mathbf{x}}_{k+1} = \text{prox}_{\tau f}(\mathbf{x}_k - \tau \nabla h(\mathbf{x}_k) - \tau \mathbf{A}^\top \bar{\mathbf{y}}_{k+1})$$

Draw $i_{k+1} \in \{1, \dots, n\}$ randomly w.p. $\mathbb{P}(i_{k+1} = i) = p_i$

$$\mathbf{x}_{k+1}^{i_{k+1}} = \bar{\mathbf{x}}_{k+1}^{i_{k+1}}$$

$$\mathbf{x}_{k+1}^j = \mathbf{x}_k^j, \forall j \neq i_{k+1}$$

$$\mathbf{y}_{k+1}^j = \bar{\mathbf{y}}_{k+1}^j + \sigma_j \theta_j [\mathbf{A}(\mathbf{x}_{k+1} - \mathbf{x}_k)]_j, \forall j \in J(i_{k+1})$$

$$\mathbf{y}_{k+1}^j = \mathbf{y}_k^j, \forall j \notin J(i_{k+1})$$

end for

step size w. dense \mathbf{A}	iter. cost
$n\tau_i\sigma\ \mathbf{A}_i\ ^2 < 1$	$\text{nnz}(\mathbf{A}_i)$

¹ β_i are coordinate-wise Lipschitz constants of ∇f

*Experiments

- Datasets with varying sparsity levels, sparse, moderately sparse, and dense.
- Comparison with dense friendly SPDHG (Chambolle et al., 2018), sparse friendly VC-CD (Fercoq&Bianchi, 2019) with duplication².
- PURE-CD stays efficient in all cases, attaining best of both worlds.

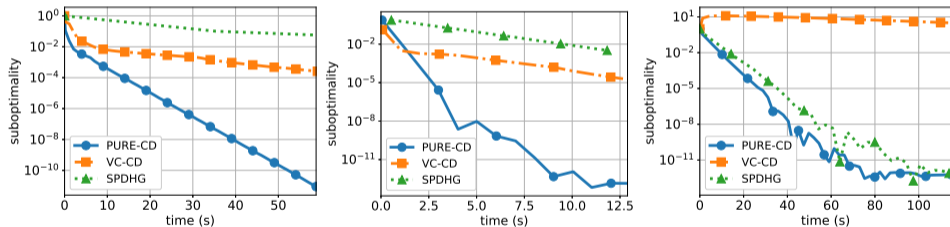


Figure: Lasso: Left: rcv1, $n = 20, 242$, $m = 47, 236$, density = 0.16%, $\lambda = 10$; Middle: w8a, $n = 49, 749$, $m = 300$, density = 3.9%, $\lambda = 10^{-1}$; Right: covtype, $n = 581, 012$, $m = 54$, density = 22.1%, $\lambda = 10$.

²Fercoq, Bianchi, A coordinate-descent primal-dual algorithm with large step size and possibly nonseparable functions, SIOPT, 2019.

*Experiments

- Strongly convex strongly concave ridge regression problems with varying regularization parameter.
- PURE-CD is competitive with state-of-the-art specialized methods for this problem.

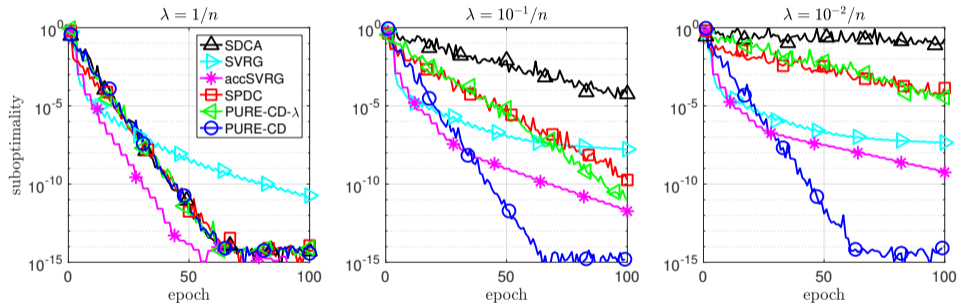


Figure: Ridge. a9a, $n = 32,561$, $m = 123$.

References I

- [1] Ahmet Alacaoglu, Olivier Fercoq, and Volkan Cevher.
On the convergence of stochastic primal-dual hybrid gradient.
arXiv preprint arXiv:1911.00799, 2019.
- [2] Dimitri P Bertsekas.
Necessary and sufficient conditions for a penalty method to be exact.
Mathematical programming, 9(1):87–99, 1975.
- [3] Volkan Cevher and B?ng Công Vũ.
A reflected forward-backward splitting method for monotone inclusions involving lipschitzian operators.
Set-Valued and Variational Analysis, pages 1–12, 2020.
- [4] A. Chambolle and T. Pock.
A first-order primal-dual algorithm for convex problems with applications to imaging.
Journal of Mathematical Imaging and Vision, 40(1):120–145, 2011.
- [5] L. Condat.
A primal–dual splitting method for convex optimization involving lipschitzian, proximable and linear composite terms.
J. Optim. Theory Appl., 158:460–479, 2013.

References II

- [6] D. Davis and W. Yin.
A three-operator splitting scheme and its optimization applications.
Tech. Report., 2015.
- [7] Chi Jin, Praneeth Netrapalli, and Michael I Jordan.
What is local optimality in nonconvex-nonconcave minimax optimization?
arXiv preprint arXiv:1902.00618, 2019.
- [8] G. M. Korpelevic.
An extragradient method for finding saddle-points and for other problems.
Ākonom. i Mat. Metody., 12(4):747–756, 1976.
- [9] Tianyi Lin, Chi Jin, and Michael I Jordan.
On gradient descent ascent for nonconvex-concave minimax problems.
arXiv preprint arXiv:1906.00331, 2019.
- [10] Songtao Lu, Ioannis Tsaknakis, Mingyi Hong, and Yongxin Chen.
Hybrid block successive approximation for one-sided non-convex min-max problems: algorithms and applications.
IEEE Transactions on Signal Processing, 2020.

References III

- [11] Yura Malitsky and Matthew K Tam.
A forward-backward splitting method for monotone inclusions without cocoercivity.
SIAM Journal on Optimization, 30(2):1451–1472, 2020.
- [12] Aryan Mokhtari, Asuman Ozdaglar, and Sarath Pattathil.
Convergence rate of $\mathcal{O}(1/k)$ for optimistic gradient and extra-gradient methods in smooth convex-concave saddle point problems.
arXiv preprint arXiv:1906.01115, 2019.
- [13] Aryan Mokhtari, Asuman Ozdaglar, and Sarath Pattathil.
A unified analysis of extra-gradient and optimistic gradient methods for saddle point problems: Proximal point approach.
In *International Conference on Artificial Intelligence and Statistics*, pages 1497–1507. PMLR, 2020.
- [14] A. Nemirovskii.
Prox-method with rate of convergence $\mathcal{O}(1/t)$ for variational inequalities with Lipschitz continuous monotone operators and smooth convex-concave saddle point problems.
SIAM J. Op, 15(1):229–251, 2004.
- [15] J. Nocedal and S.J. Wright.
Numerical Optimization.
Springer Series in Operations Research and Financial Engineering. Springer, 2 edition, 2006.

References IV

- [16] Leonid Denisovich Popov.
A modification of the arrow-hurwicz method for search of saddle points.
Mathematical notes of the Academy of Sciences of the USSR, 28(5):845–848, 1980.
- [17] Sasha Rakhlin and Karthik Sridharan.
Optimization, learning, and games with predictable sequences.
In *Advances in Neural Information Processing Systems*, pages 3066–3074, 2013.
- [18] R. T. Rockafellar.
Augmented lagrangians and applications of the proximal point algorithm in convex programming.
Mathematics of Operations Research, 1:97–116, 1976.
- [19] Ohad Shamir.
Can we find near-approximately-stationary points of nonsmooth nonconvex functions?
arXiv preprint arXiv:2002.11962, 2020.
- [20] Quoc Tran-Dinh, Olivier Fercoq, and Volkan Cevher.
A smooth primal-dual optimization framework for nonsmooth composite convex minimization.
SIAM Journal on Optimization, 28(1):96–134, 2018.

References V

- [21] Bang Cong Vu.
A splitting algorithm for dual monotone inclusions involving cocoercive operators.
Advances in Computational Mathematics, 38(3):667–681, 2013.
- [22] Zi Xu, Huiling Zhang, Yang Xu, and Guanghui Lan.
A unified single-loop alternating gradient projection algorithm for nonconvex-concave and convex-nonconcave minimax problems.
arXiv preprint arXiv:2006.02032, 2020.
- [23] Ming Yan.
A new primal–dual algorithm for minimizing the sum of three functions with a linear operator.
Journal of Scientific Computing, 76(3):1698–1717, 2018.
- [24] Alp Yurtsever, Bang Công Vu, and Volkan Cevher.
Stochastic three-composite convex minimization.
In *Advances in Neural Information Processing Systems*, pages 4329–4337, 2016.
- [25] Jingzhao Zhang, Hongzhou Lin, Suvrit Sra, and Ali Jadbabaie.
On complexity of finding stationary points of nonsmooth nonconvex functions.
arXiv preprint arXiv:2002.04130, 2020.