

Mathematics of Data: From Theory to Computation

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Outline

► Today

1. Convex constrained optimization and motivating examples
2. Optimality condition
3. Conjugate functions
4. Monotone inclusion and monotone mixed variational inequality formulations
5. Chambolle-Pock's primal-dual method
6. Primal-dual hybrid gradient method
7. Splitting methods
8. Model-based excessive gap primal-dual method

► Next week

1. Disciplined convex programming

Motivation

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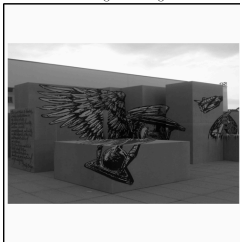
- ▶ **Primal-dual convex optimization methods** are **powerful** for solving **saddle point problems** as well as **constrained convex optimization problems**.
- ▶ This lecture **aims** at presenting some **emerging primal-dual methods** which have been recently used to solve **many practical problems** in signal/image processing, machine learning and statistics.
- ▶ This lecture is a **continuation** of Lecture 7.

Motivating example: image denoising via anisotropicTV-norm

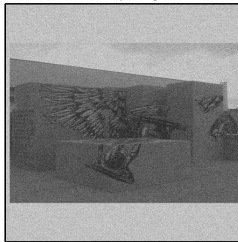
We consider an image denoising problem with anisotropic total variation norm:

- Given a noisy image $\mathbf{b} \in \mathbb{R}^{m \times n}$. The goal is to recover a clean image from \mathbf{b} using anisotropic total variation norm.

Original image



Noisy image

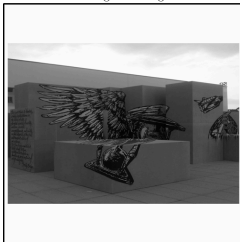


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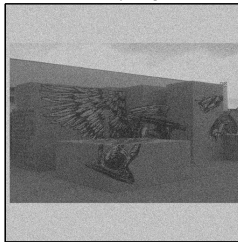
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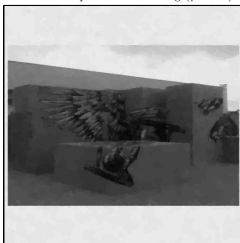
Original image



Noisy image



Anisotropic TV denoising ($\rho = 20$)



Anisotropic TV denoising ($\rho = 40$)



Motivating example: image denoising via anisotropicTV-norm

We consider an image denoising problem with anisotropic total variation norm:

- ▶ Given a noisy image $\mathbf{b} \in \mathbb{R}^{m \times n}$. The goal is to recover a clean image from \mathbf{b} using anisotropic total variation norm.
- ▶ This problem can be formulated as a convex optimization problem:

$$\min_{\mathbf{x} \in \mathbb{R}^{m \times n}} \frac{1}{2} \|\mathbf{x} - \mathbf{b}\|_F^2 + \rho \|\mathbf{D}\mathbf{x}\|_1, \quad (1)$$

where $\rho > 0$ is a regularization parameter and \mathbf{D} is a given matrix representing the total variation of \mathbf{x} .

There are different ways to reformulate problem (1), for example:

- ▶ Since $\|\mathbf{z}\|_1 = \max_{\|\mathbf{u}\|_\infty \leq 1} \mathbf{u}^T \mathbf{z}$, we can reformulate (1) as a saddle point problem:

$$\min_{\mathbf{x} \in \mathbb{R}^{m \times n}} \max_{\|\mathbf{u}\|_\infty \leq 1} \left\{ \frac{1}{2} \|\mathbf{x} - \mathbf{b}\|_F^2 + \rho \mathbf{u}^T \mathbf{D}\mathbf{x} \right\},$$

- ▶ We can also reformulate (1) as a constrained convex minimization problem:

$$\begin{aligned} \min_{\mathbf{x}, \mathbf{r}} \quad & \left\{ \frac{1}{2} \|\mathbf{x} - \mathbf{b}\|_F^2 + \rho \|\mathbf{r}\|_1 \right\}, \\ \text{s.t.} \quad & \mathbf{D}\mathbf{x} - \mathbf{r} = 0. \end{aligned}$$

In this lecture, we present several emerging methods to solve both the saddle point formulation and the constrained formulation of (1).

Mathematical form of constrained convex optimization

Constrained convex optimization setting

$$f^* := \begin{cases} \min_{\mathbf{x} \in \mathbb{R}^p} & f(\mathbf{x}) \\ \text{s.t.} & \mathbf{A}\mathbf{x} = \mathbf{b}. \end{cases} \quad (2)$$

- ▶ $f \in \mathcal{F}(\mathbb{R}^p)$ is a proper, closed and **convex** function (see Lecture 2).
- ▶ $\mathbf{A} \in \mathbb{R}^{n \times p}$ is full-row rank ($n < p$), $\mathbf{b} \in \mathbb{R}^n$.

We can **incorporating constraints** $\mathbf{x} \in \mathcal{X}$ for a given closed and convex set \mathcal{X} via its **indicator function** $\iota_{\mathcal{X}}$, i.e.:

$$\boxed{f(\mathbf{x}) \leftarrow f(\mathbf{x}) + \iota_{\mathcal{X}}(\mathbf{x})} \quad \text{where} \quad \iota_{\mathcal{X}}(\mathbf{x}) := \begin{cases} 0 & \text{if } \mathbf{x} \in \mathcal{X} \\ +\infty & \text{otherwise.} \end{cases}$$

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Common structures

As in Lecture 7, **methods** presented in this lecture also rely on the **two common structures**:

- ▶ **Decomposability** of f .
- ▶ **Tractable proximity** of f .

Structures of constrained convex optimization problems

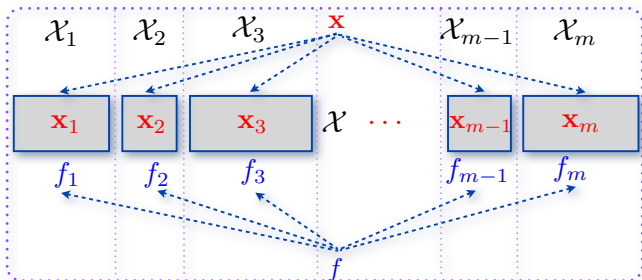
Decomposable structure

The function f can be decomposed as

$$f(\mathbf{x}) := \sum_{i=1}^m f_i(\mathbf{x}_i) \quad (3)$$

where $m \geq 1$ is the **number of components**, \mathbf{x}_i is a **sub-vector** (component) of \mathbf{x} , $f_i : \mathbb{R}^{p_i} \rightarrow \mathbb{R} \cup \{+\infty\}$ is **convex** and $\sum_{i=1}^m p_i = p$.

Special case: $m = 2$, which already covers many **important instances** (see Lecture 7).



Structures of constrained convex optimization problems

Tractable proximity

Each component f_i has a “tractably proximal operator” ($i = 1, \dots, m$), i.e.:

$$\text{prox}_{f_i}(\mathbf{x}_i) := \underset{\mathbf{z}_i \in \mathbb{R}^{p_i}}{\text{argmin}} \left\{ f_i(\mathbf{z}_i) + (1/2) \|\mathbf{z}_i - \mathbf{x}_i\|_2^2 \right\} \quad (4)$$

can be solved “efficiently”:

- ▶ (4) has a closed form solution (with low computational cost)
- ▶ (4) can be solved in polynomial time.

Example (Tractable proximity functions)

- ▶ One-variable functions
 - ▶ Smooth functions, e.g., $f(x) := x - 2 \log(1 + x)$
 - ▶ Nonsmooth functions, e.g., $f(x) := |x|$
- ▶ Separable functions, e.g., $f(\mathbf{x}) := \sum_{i=1}^p \|\mathbf{x}_i\|_2$, where $\mathbf{x} := (\mathbf{x}_1^T, \dots, \mathbf{x}_p^T)^T$.
- ▶ The indicator function ι of boxes, cones (\mathbb{R}_+^p , \mathbb{S}_+^p and Lorentz cone) and simplex.
- ▶ More examples can be found in Lectures 4 and 5.

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1. Disciplined convex programming

Solutions and solution set

Optimal solutions and optimal solution set

We define the feasible set of (2): (cf. $f^* := \min_{\mathbf{x} \in \mathbb{R}^p} \{f(\mathbf{x}) : \mathbf{Ax} = \mathbf{b}\}$) as:

$$\mathcal{D} := \{\mathbf{x} \in \mathbb{R}^p : \mathbf{Ax} = \mathbf{b}\}.$$

A feasible point $\mathbf{x}^* \in \mathcal{D}$ is called a **globally optimal solution** (or solution) of (2) if

$$f(\mathbf{x}^*) \leq f(\mathbf{x}), \quad \forall \mathbf{x} \in \mathcal{D}.$$

All **solutions** of (2) forms the **solution set** \mathcal{X}^* of (2).

- ▶ The solution set \mathcal{X}^* is **closed** and **convex**.
- ▶ **Numerical solution methods** often try to find an **approximation** \mathbf{x}_ϵ^* of **one solution** $\mathbf{x}^* \in \mathcal{X}^*$ in the following sense:

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Approximate solution

Given a **tolerance** $\epsilon \geq 0$, a point $\mathbf{x}_\epsilon^* \in \mathbb{R}^p$ is called an **ϵ -solution** of (2) if

$$\begin{cases} |f(\mathbf{x}_\epsilon^*) - f^*| \leq \epsilon & \text{(objective residual),} \\ \|\mathbf{Ax}_\epsilon^* - \mathbf{b}\| \leq \epsilon & \text{(feasibility gap).} \end{cases}$$

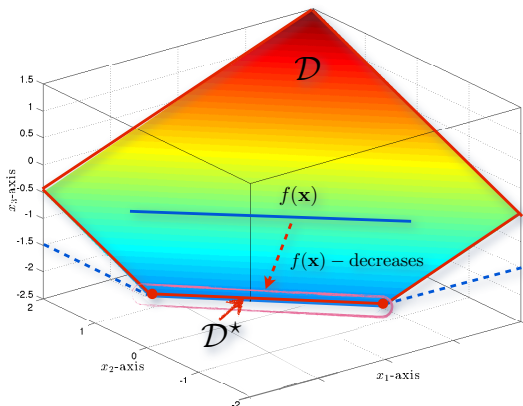
Remark: we can use different tolerances for the objective residual and feasibility gap.

Example: Feasible set and solution set

Consider a constrained convex problem:

$$\begin{array}{ll}\min_{\mathbf{x} \in \mathbb{R}^3} & \{f(\mathbf{x}) := x_3\}, \\ \text{s.t.} & 2x_3 - x_1 - x_2 = -1, \\ & \mathbf{x} \in \mathcal{X} := [-2, 2] \times [-2, 2] \times [-1.5, 1.5].\end{array}$$

The **feasible set** $\mathcal{D} := \{\mathbf{x} \in \mathbb{R}^3 : 2x_3 - x_1 - x_2 = -1, \mathbf{x} \in \mathcal{X}\}$ and the **solution set** \mathcal{D}^* of this problem are plotted in the figure below. \mathcal{D}^* is in fact a **segment** (many solutions), which is the lowest edge of the polytope \mathcal{D} .



Optimality condition

Lagrange function

$$\mathcal{L}(\mathbf{x}, \lambda) := f(\mathbf{x}) + \lambda^T (\mathbf{A}\mathbf{x} - \mathbf{b}).$$

Here, $\lambda \in \mathbb{R}^n$ is the vector of **Lagrange multipliers** (or **dual** variables).

Optimality condition

The **optimality condition** of (2) can be written as

$$\begin{cases} 0 \in \mathbf{A}^T \lambda^* + \partial f(\mathbf{x}^*), \\ 0 = \mathbf{A}\mathbf{x}^* - \mathbf{b}. \end{cases} \quad (5)$$

Here the subdifferential of f at \mathbf{x}^* is defined as (see Lecture 2):

$$\partial f(\mathbf{x}^*) := \{\mathbf{z} \in \mathbb{R}^p : f(\mathbf{y}) \geq f(\mathbf{x}^*) + \mathbf{z}^T(\mathbf{y} - \mathbf{x}^*), \forall \mathbf{y} \in \mathbb{R}^p\}.$$

- ▶ The condition (5) is the **KKT** (Karush-Kuhn-Tucker) condition.
- ▶ Any point $(\mathbf{x}^*, \lambda^*)$ satisfying (5) is called a **KKT point**.
- ▶ \mathbf{x}^* is called a **stationary point** and λ^* is the corresponding **multipliers**.

Conjugation of functions

- ▶ **Duality** is a **central concept** in optimization, especially in **convex optimization**.
- ▶ We review the notion of **Fenchel's conjugate function** and its basic properties which will be used to define the dual problem.
- ▶ We limit our definition to the class of **convex functions** $f \in \mathcal{F}(\mathbb{R}^p)$.

Definition

Let \mathcal{Q} be a **predefined Euclidean space** and \mathcal{Q}^* be its **dual space**. Given a **proper, closed and convex function** $f : \mathcal{Q} \rightarrow \mathbb{R} \cup \{+\infty\}$, the function $f^* : \mathcal{Q}^* \rightarrow \mathbb{R} \cup \{+\infty\}$ such that

$$f^*(\mathbf{y}) = \sup_{\mathbf{x} \in \text{dom}(f)} \left\{ \mathbf{y}^T \mathbf{x} - f(\mathbf{x}) \right\}$$

is called the **Fenchel conjugate** (or conjugate) of f .

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is called the **Fenchel conjugate** (or conjugate) of f .

- ▶ f^* is a **convex** and lower, semicontinuous function by construction (as the supremum of affine functions of \mathbf{y}).
- ▶ The **conjugate** of the **conjugate** of a convex function f is ... the same function f ; i.e., $f^{**} = f$ for $f \in \mathcal{F}(\mathcal{Q})$.

Two basic properties of the function and its conjugation

Lemma (Fenchel-Young inequality (Property 1))

Let $f : \mathcal{Q} \rightarrow \mathbb{R} \cup \{+\infty\}$ and $f^ : Q^* \rightarrow \mathbb{R} \cup \{+\infty\}$ be a function and its conjugation; here Q^* be the dual space of \mathcal{Q} . Then, the following inequality holds true:*

$$f(\mathbf{x}) + f^*(\mathbf{y}) \geq \mathbf{x}^T \mathbf{y}, \quad \forall \mathbf{x} \in \mathcal{Q}, \mathbf{y} \in Q^*.$$

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$$f(\mathbf{x}) + f^*(\mathbf{y}) \geq \mathbf{x}^T \mathbf{y}, \quad \forall \mathbf{x} \in \mathcal{Q}, \mathbf{y} \in Q^*.$$

- ▶ Since f^* is proper, closed and convex, its subdifferential ∂f^* exists for any \mathbf{y} in the relative interior of its domain.
- ▶ For $f \in \mathcal{F}(\mathcal{Q})$, if the subdifferential of f and f^* exists, then we have the following relation:

Lemma (Subgradient property (Property 2))

Let $\mathbf{y} \in \partial f(\mathbf{x})$ for some $\mathbf{x} \in \text{dom}(f)$. Then $\mathbf{y} \in \text{dom}(f^*)$ and vice versa. Moreover, we have

$$\mathbf{u} \in \partial f(\mathbf{x}) \Leftrightarrow \mathbf{x} \in \partial f^*(\mathbf{u}).$$

Conjugation of functions

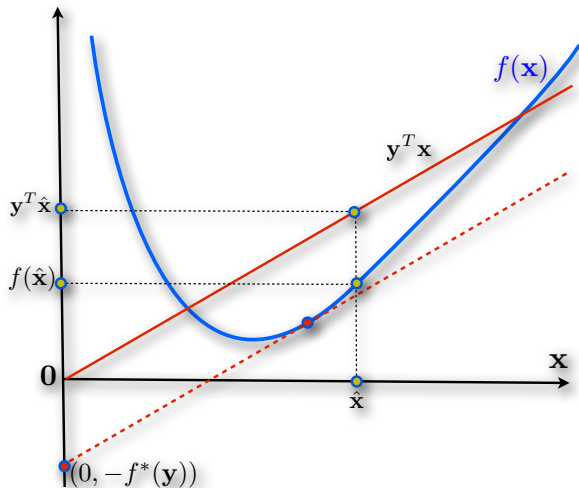


Figure: The conjugate function $f^*(y)$ is the maximum gap between the linear function $x^T y$ (red line) and $f(x)$, as shown in dashed line.

Example 1: Convex quadratic function

Example (Strictly convex quadratic function)

- ▶ Given a **symmetric positive definite** matrix Φ , ($\Phi \succ 0$) and a vector \mathbf{b} .
- ▶ Let $f(\mathbf{x}) := \frac{1}{2}\mathbf{x}^T \Phi \mathbf{x} + \mathbf{b}^T \mathbf{x}$ be a quadratic function for $\mathbf{x} \in \mathbb{R}^p$.
- ▶ It is clear that $\text{dom}(f) = \mathbb{R}^p$.
- ▶ By definition,

$$f^*(\mathbf{y}) := \sup_{\mathbf{x} \in \mathbb{R}^p} \left\{ \mathbf{y}^T \mathbf{x} - \frac{1}{2}\mathbf{x}^T \Phi \mathbf{x} - \mathbf{b}^T \mathbf{x} \right\}.$$

- ▶ Since this is an **unconstrained convex problem**, the **maximum** is attained when $\Phi \mathbf{x}^* + \mathbf{b} = \mathbf{y}$, which leads to $\mathbf{x}^* = \Phi^{-1}(\mathbf{y} - \mathbf{b})$.
- ▶ Hence, we have

$$f^*(\mathbf{y}) = \frac{1}{2}(\mathbf{y} - \mathbf{b})^T \Phi^{-1}(\mathbf{y} - \mathbf{b}) = \frac{1}{2}\mathbf{y}^T \Phi^{-1} \mathbf{y} - (\Phi^{-1} \mathbf{b})^T \mathbf{y} + \frac{1}{2}\mathbf{b}^T \Phi^{-1} \mathbf{b},$$

which is again a **convex quadratic function** with $\text{dom}(f^*) = \mathbb{R}^p$.

- ▶ Since $\nabla f(\mathbf{x}) = \Phi \mathbf{x} + \mathbf{b} := \mathbf{u}$ and $\nabla f^*(\mathbf{y}) = \Phi^{-1}(\mathbf{y} - \mathbf{b})$, we can see that $\mathbf{x} = \Phi^{-1}(\mathbf{u} - \mathbf{b}) = \nabla f^*(\mathbf{u})$.

Example 2: Log-determinant

Example (Log-determinant function)

- ▶ Let $f(\mathbf{X}) := -\log \det(\mathbf{X})$, where $\text{dom}(f) \equiv \mathbb{S}_{++}^p$.
- ▶ By definition, we have

$$f^*(\mathbf{Y}) = \sup_{\mathbf{X} \in \text{dom}(f)} \{\text{tr}(\mathbf{Y}\mathbf{X}) + \log \det(\mathbf{X})\},$$

- ▶ One can show that the above is unbounded above unless $\mathbf{Y} \prec 0$.
- ▶ To find the maximum of the above problem, we have:

$$\nabla (\text{tr}(\mathbf{Y}\mathbf{X}) + \log \det(\mathbf{X})) = 0 \Rightarrow \mathbf{X}^* = -\mathbf{Y}^{-1},$$

and thus,

$$f^*(\mathbf{Y}) = -\log \det(-\mathbf{Y}) - p, \quad \text{with } \text{dom}(f^*) = -\mathbb{S}_{++}^p.$$

- ▶ Since $\nabla f(\mathbf{X}) = -\mathbf{X}^{-1} := \mathbf{U}$ and $\nabla f^*(\mathbf{Y}) = -\mathbf{Y}^{-1}$, we have $\mathbf{X} = -\mathbf{U}^{-1} = \nabla f^*(\mathbf{U})$.

Conjugation of functions

Example

$f(\mathbf{x})$	$\text{dom}(f)/\text{dom}(f^*)$	$f^*(\mathbf{y})$
$f(a\mathbf{x})$ (where $a \neq 0$)	$\mathcal{Q}/\mathcal{Q}^*$	$f^*\left(\frac{\mathbf{y}}{a}\right)$
$f(\mathbf{x} + \beta)$	$\mathcal{Q}/\mathcal{Q}^*$	$f^*(\mathbf{y}) - \langle \beta, \mathbf{y} \rangle$
$\alpha f(\mathbf{x})$ (where $\alpha > 0$)	$\mathcal{Q}/\mathcal{Q}^*$	$\alpha f^*\left(\frac{\mathbf{y}}{\alpha}\right)$
$\frac{\ \mathbf{x}\ ^r}{r}$ (where $r > 1$)	$\mathbb{R}^p/\mathbb{R}^p$	$\frac{\ \mathbf{y}\ ^q}{q}$ (where $\frac{1}{r} + \frac{1}{q} = 1$)
$-\log(x)$	$\mathbb{R}_{++}/\mathbb{R}_{--}$	$-(1 + \log(y))$
e^x	\mathbb{R}/\mathbb{R}_+	$\begin{cases} y \log(y) - y & \text{if } y > 0 \\ 0 & \text{if } y = 0 \end{cases}$

Table: Legendre transforms (conjugations) for many common functions as well as a few useful properties.

Dual problem

- ▶ In Lecture 7 we have used the [Lagrange duality](#) theory to present **methods of multipliers**.
- ▶ In this lecture, we use the [Fenchel duality](#) theory to define the **dual problem** formulation and develop **primal-dual methods** for solving (2).

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Dual formulation

- ▶ From the **optimality condition (5)**: $\begin{cases} 0 \in \mathbf{A}^T \lambda^* + \partial f(\mathbf{x}^*) \\ 0 = \mathbf{A} \mathbf{x}^* - \mathbf{b} \end{cases}$ we have

$$-\mathbf{A}^T \lambda^* \in \partial f(\mathbf{x}^*).$$

- ▶ **The last expression** is equivalent to $\mathbf{x}^* \in \partial f^*(-\mathbf{A}^T \lambda^*)$ (see Property 2).
- ▶ Since $\mathbf{A} \mathbf{x}^* - \mathbf{b} = 0$, using $\mathbf{x}^* \in \partial f^*(-\mathbf{A}^T \lambda^*)$, we have

$$0 \in \mathbf{A} \partial f^*(-\mathbf{A}^T \lambda^*) - \mathbf{b}. \quad (6)$$

(6) is exactly the **optimality condition** of

$$\max_{\lambda \in \mathbb{R}^n} \left\{ -f^*(-\mathbf{A}^T \lambda) - \mathbf{b}^T \lambda \right\}. \quad (7)$$

(7) is still a **convex problem** and equivalent to $\min_{\lambda} \{f^*(-\mathbf{A}^T \lambda) + \mathbf{b}^T \lambda\}$.

Decomposable structure

Dual problem of the decomposable objective function

If $f(\mathbf{x}) := \sum_{i=1}^m f_i(\mathbf{x}_i)$ then

$$\max_{\lambda \in \mathbb{R}^n} \left\{ - \sum_{i=1}^m f_i^*(\mathbf{A}_i^T \lambda) - \mathbf{b}^T \lambda \right\}.$$

where $\mathbf{A} \equiv [\mathbf{A}_1, \dots, \mathbf{A}_m]$.

Note: The **evaluation** of the **dual objective function** and its **gradient** can be computed in **parallel**.

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Dual formulation of empirical risk minimization

An **empirical risk minimization** problem can be expressed as

$$\min_{\mathbf{z} \in \mathbb{R}^n} \sum_{i=1}^m f_i(\mathbf{A}_i^T \mathbf{z} + \mathbf{b}_i).$$

Its **Fenchel dual problem** therefore can be written as:

$$\min_{\mathbf{x} \in \mathbb{R}^p} \sum_{i=1}^m \left\{ f_i^*(\mathbf{x}_i) - \mathbf{b}_i^T \mathbf{x}_i \right\} \quad \text{s.t.} \quad \sum_{i=1}^m \mathbf{A}_i \mathbf{x}_i = \mathbf{0}.$$

Example: Dual problem of the basis pursuit

Basis pursuit

Consider the following basis pursuit problem:

$$\begin{array}{ll} \min_{\mathbf{x} \in \mathbb{R}^p} & \{f(\mathbf{x}) := \|\mathbf{x}\|_1 = \sum_{i=1}^p |x_i|\} \\ \text{s.t.} & \mathbf{Ax} = \mathbf{b}. \end{array}$$

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Dual problem of basis pursuit

- ▶ The Fenchel dual function $f^*(\mathbf{u}) := \sup_{\mathbf{x} \in \mathbb{R}^p} \{\mathbf{u}^T \mathbf{x} - \|\mathbf{x}\|_1\}$ becomes

$$f^*(\mathbf{u}) = \iota_{\{\|\mathbf{u}\|_\infty \leq 1\}}(\mathbf{u}) = \begin{cases} 0 & \text{if } \|\mathbf{u}\|_\infty \leq 1, \\ +\infty & \text{otherwise.} \end{cases}$$

- ▶ The dual problem

$$\max_{\lambda} \left\{ -\iota_{\{\|\mathbf{A}^T \lambda\|_\infty \leq 1\}}(-\mathbf{A}^T \lambda) - \mathbf{b}^T \lambda \right\}$$

- ▶ Equivalent expression:

$$\min_{\lambda \in \mathbb{R}^n} \mathbf{b}^T \lambda \quad \text{s.t.} \quad \|\mathbf{A}^T \lambda\|_\infty \leq 1.$$

Set-valued mappings

In the previous slide, we have seen the **subdifferential** $\partial f(\mathbf{x})$ of a **convex function** f at a given point \mathbf{x} :

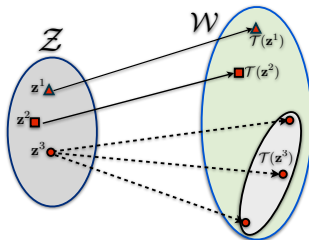
$\partial f(\mathbf{x}) :=$ set of all subgradients of f at \mathbf{x}

Example: Subdifferential of $|\mathbf{x}|$ in \mathbb{R} :

$$\mathcal{T}(\mathbf{x}) = \partial|\mathbf{x}| = \begin{cases} +1 & \text{if } x > 0 \text{ (single value),} \\ -1 & \text{if } x < 0 \text{ (single value),} \\ [-1, 1] & \text{if } x = 0 \text{ (multiple values).} \end{cases}$$

Set-valued mappings:

- ▶ For any **convex set** \mathcal{W} , we denote by $2^{\mathcal{W}}$ the **set** of **all subsets** of \mathcal{W} .
- ▶ $\mathcal{T} : \mathcal{Z} \rightrightarrows 2^{\mathcal{W}}$ is a **set-valued mapping** if for $\mathbf{z} \in \mathcal{Z}$, $\mathcal{T}(\mathbf{z})$ is a **subset** in \mathcal{W} .



Roughly speaking, a **mapping** that **produces more** than **one output values** for at **least one input** is called a **set-valued mapping**.

Monotone operators

For a **set-valued mapping** $\mathcal{T} : \mathcal{Z} \rightrightarrows 2^{\mathcal{Z}}$, we define

- ▶ The **domain** of \mathcal{T} as $\text{dom}(\mathcal{T}) := \{\mathbf{z} \in \mathcal{Z} : \mathcal{T}(\mathbf{z}) \neq \emptyset\}$.
- ▶ The **graph** of \mathcal{T} is $\text{graph}(\mathcal{T}) := \{(\mathbf{z}, \mathbf{v}) : \mathbf{v} \in \mathcal{T}(\mathbf{z}), \mathbf{z} \in \text{dom}(\mathcal{T})\}$.

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Monotonicity can be considered as an equivalent property of **convexity** acting on the differential ∇f or subdifferential ∂f of the function instead of the function value.

For smooth and convex function, monotonicity of ∇f means that

$$(\nabla f(\mathbf{x}) - \nabla f(\hat{\mathbf{x}}))^T (\mathbf{x} - \hat{\mathbf{x}}) \geq 0, \quad \forall \mathbf{x}, \hat{\mathbf{x}} \in \text{dom}(f).$$

This inequality is the sum of $f(\mathbf{x}) \geq f(\hat{\mathbf{x}}) + \nabla f(\hat{\mathbf{x}})^T (\mathbf{x} - \hat{\mathbf{x}})$ and $f(\hat{\mathbf{x}}) \geq f(\mathbf{x}) + \nabla f(\mathbf{x})^T (\hat{\mathbf{x}} - \mathbf{x})$.

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Monotonicity

Given a **convex set** \mathcal{Z} and a **set-valued mapping** $\mathcal{T} : \mathcal{Z} \rightrightarrows 2^{\mathcal{Z}}$.

- ▶ \mathcal{T} is called **μ -strongly monotone** on \mathcal{Z} if for any \mathbf{z} and $\hat{\mathbf{z}}$ in \mathcal{Z} :

$$(\mathbf{u} - \hat{\mathbf{u}})^T (\mathbf{z} - \hat{\mathbf{z}}) \geq \mu \|\mathbf{z} - \hat{\mathbf{z}}\|^2, \quad \forall \mathbf{u} \in \mathcal{T}(\mathbf{z}), \hat{\mathbf{u}} \in \mathcal{T}(\hat{\mathbf{z}}).$$

- ▶ If $\mu = 0$, then we say that \mathcal{T} is **monotone**.
- ▶ If $\mu > 0$, then we say that \mathcal{T} is **strongly monotone** with the **parameter** μ .

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If \mathcal{T} is **single-valued**, then the condition reduces to

$$(\mathbf{z} - \hat{\mathbf{z}})^T (\mathcal{T}(\mathbf{z}) - \mathcal{T}(\hat{\mathbf{z}})) \geq \mu \|\mathbf{z} - \hat{\mathbf{z}}\|^2, \quad \forall \mathbf{z}, \hat{\mathbf{z}} \in \mathcal{Z}.$$

Monotone function vs nonmonotone functions

$$(T(\hat{\mathbf{x}}) - T(\mathbf{x}))^T(\hat{\mathbf{x}} - \mathbf{x}) \geq 0 \quad (T(\hat{\mathbf{x}}) - T(\mathbf{x}))^T(\hat{\mathbf{x}} - \mathbf{x}) \leq 0 \quad (T(\hat{\mathbf{x}}) - T(\mathbf{x}))^T(\hat{\mathbf{x}} - \mathbf{x}) \leq 0$$

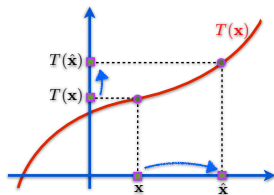
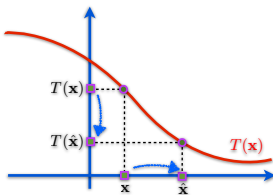


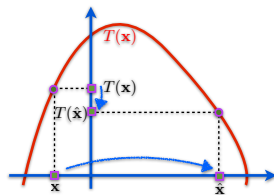
Figure: Monotone function

increasing function



Nonmonotone function

decreasing function



Nonmonotone function

Maximal monotone operators

Definition (Maximal monotonicity)

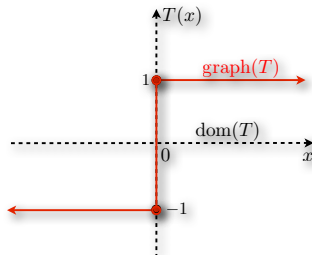
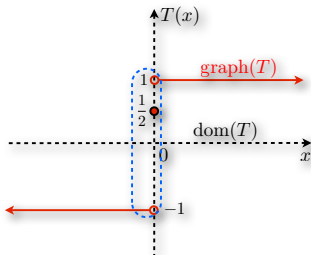
A set-valued mapping \mathcal{T} is called **maximal monotone** if:

- ▶ \mathcal{T} is **monotone**
- ▶ There exists **no** other **monotone mapping** $\tilde{\mathcal{T}}$ such that $\text{graph}(\mathcal{T}) \subset \text{graph}(\tilde{\mathcal{T}})$.

Example (Visualization of a maximal monotone operator)

Consider the mapping \mathcal{T} from \mathbb{R} to $2^{\mathbb{R}}$ as follows:

- ▶ **Nonmaximal monotone (left figure):** $T(x) = 1$ if $x > 0$, $T(x) = -1$ if $x < 0$ and $T(x) = 1/2$ if $x = 0 \Rightarrow T$ is **not maximal monotone**.
- ▶ **Maximal monotone (right figure):** $T(x) = 1$ if $x > 0$, $T(x) = -1$ if $x < 0$ and $T(x) = [-1, 1]$ if $x = 0 \Rightarrow T$ is **maximal monotone**.



Example of maximal monotone operators

- ▶ Affine mapping $\mathcal{T}(\mathbf{z}) := \mathbf{H}\mathbf{z} + \mathbf{h}$ is maximal monotone if \mathbf{H} is positive semidefinite.

- ▶ We have

$$(\mathcal{T}(\mathbf{z}) - \mathcal{T}(\hat{\mathbf{z}}))^T (\mathbf{z} - \hat{\mathbf{z}}) = (\mathbf{z} - \hat{\mathbf{z}})^T \mathbf{H} (\mathbf{z} - \hat{\mathbf{z}}) \geq \sigma_{\min}(\mathbf{H}) \|\mathbf{z} - \hat{\mathbf{z}}\|^2,$$

where $\mu = \sigma_{\min}(\mathbf{H})$ the smallest singular value of \mathbf{H}

- ▶ \mathcal{T} is strongly monotone if \mathbf{H} is positive definite.
- ▶ Any nondecreasing function $f : \mathbb{R} \rightarrow \mathbb{R}$ is monotone.

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- ▶ \mathcal{T} is strongly monotone if \mathbf{H} is positive definite.
- ▶ Any nondecreasing function $f : \mathbb{R} \rightarrow \mathbb{R}$ is monotone.
- ▶ The subdifferential ∂f of a convex function f is maximal monotone.
 - ▶ By definition, any $\mathbf{u} \in \partial f(\mathbf{x})$, one has $f(\hat{\mathbf{x}}) - f(\mathbf{x}) \geq \mathbf{u}^T(\hat{\mathbf{x}} - \mathbf{x})$ for any $\hat{\mathbf{x}}$.
 - ▶ Similarly, $\hat{\mathbf{u}} \in \partial f(\hat{\mathbf{x}})$, then $f(\mathbf{x}) - f(\hat{\mathbf{x}}) \geq \hat{\mathbf{u}}^T(\mathbf{x} - \hat{\mathbf{x}})$.
 - ▶ Summing up these inequalities, we obtain $(\mathbf{u} - \hat{\mathbf{u}})^T(\mathbf{x} - \hat{\mathbf{x}}) \geq 0$.
- ▶ If f is strongly convex with the convexity parameter μ , then ∂f is strongly monotone with monotonicity parameter μ .

Example of maximal monotone operators

- ▶ **Affine mapping** $\mathcal{T}(\mathbf{z}) := \mathbf{H}\mathbf{z} + \mathbf{h}$ is **maximal monotone** if \mathbf{H} is positive semidefinite.

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 - ▶ Summing up these inequalities, we obtain $(\mathbf{u} - \hat{\mathbf{u}})^T(\mathbf{x} - \hat{\mathbf{x}}) \geq 0$.
- ▶ If f is **strongly convex** with the **convexity parameter** μ , then ∂f is **strongly monotone** with monotonicity parameter μ .
- ▶ The **normal cone** $\mathcal{N}_{\mathcal{X}}$ of a **nonempty, closed and convex** set \mathcal{X} is also a **monotone mapping**
 - ▶ Since it is the subdifferential of the indicator function $\iota_{\mathcal{X}}$, which is **proper, closed and convex**.

Monotonicity of the normal cone

Given a **nonempty, closed and convex** set \mathcal{X} . The **normal cone** of \mathcal{X} at \mathbf{x} is defined as

$$\mathcal{N}_{\mathcal{X}}(\mathbf{x}) := \begin{cases} \{\mathbf{u} : \mathbf{u}^T(\mathbf{x} - \mathbf{y}) \geq 0, \forall \mathbf{y} \in \mathcal{X}\} & \text{if } \mathbf{x} \in \mathcal{X}, \\ \emptyset & \text{otherwise.} \end{cases}$$

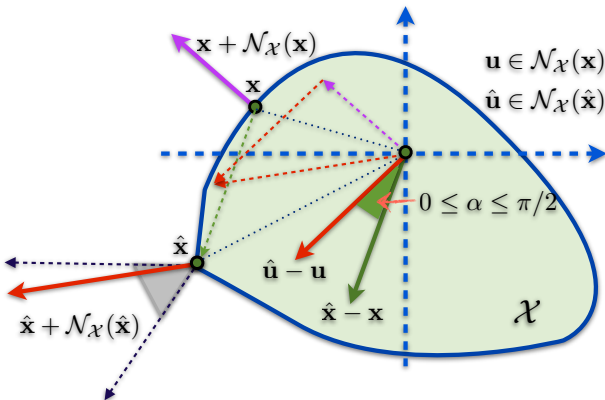
Then $\mathcal{N}_{\mathcal{X}}(\cdot)$ is a **set-valued mapping** and is **monotone**.

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Resolvent and relation to prox-operator

Resolvent of a maximal monotone mapping

Given a maximal monotone mapping \mathcal{T} .

- ▶ The **resolvent** $\mathcal{J}_{\mathcal{T}}$ of \mathcal{T} at \mathbf{w} is defined as a **solution** of the *inclusion* w.r.t. \mathbf{z} :

$$\mathbf{w} \in \mathbf{z} + \mathcal{T}(\mathbf{z}).$$

- ▶ Conventionally, we can write $\mathcal{J}_{\mathcal{T}}(\mathbf{w}) := (\mathbb{I} + \mathcal{T})^{-1}(\mathbf{w})$.

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Lemma (Well-definedness [12])

If \mathcal{T} is **maximal monotone** then $\mathcal{J}_{\mathcal{T}}(\mathbf{w})$ is **well-defined** and **single-valued**.

Remark: If \mathcal{T} is **not** maximal monotone, then $\mathcal{J}_{\mathcal{T}}(\mathbf{w})$ may **not** be well-defined.

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Relation to prox operator

Let $\mathcal{T} := \partial f$ the **subdifferential** of a **proper, closed and convex** function $f \in \mathcal{F}(\mathbb{R}^p)$.
Then \mathcal{T} is **maximal monotone** and

$$\mathcal{J}_{\partial f}(\cdot) \equiv \text{prox}_f(\cdot).$$

Example: Resolvent of the normal cone

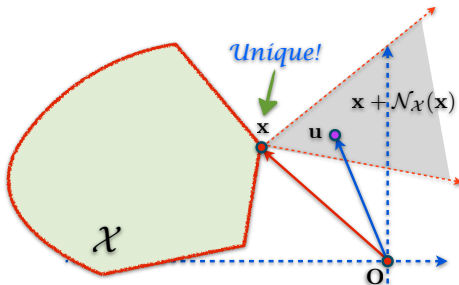
Consider the inclusion $\mathbf{u} \in \mathbf{x} + \mathcal{N}_{\mathcal{X}}(\mathbf{x})$. We can write $\mathbf{u} - \mathbf{x} \in \mathcal{N}_{\mathcal{X}}(\mathbf{x})$. By definition of $\mathcal{N}_{\mathcal{X}}(\mathbf{x})$ we have $(\mathbf{u} - \mathbf{x})^T \mathbf{x} \geq (\mathbf{u} - \mathbf{x})^T \mathbf{y}$ for all $\mathbf{y} \in \mathcal{X}$. Hence, we can write

$$(\mathbf{x} - \mathbf{u})^T (\mathbf{y} - \mathbf{u}) \geq \|\mathbf{x} - \mathbf{u}\|_2^2, \quad \forall \mathbf{y} \in \mathcal{X}.$$

This inequality shows that x is the solution of

$$\mathbf{x} = \underset{\mathbf{y} \in \mathcal{X}}{\operatorname{argmin}} \|\mathbf{y} - \mathbf{u}\|_2^2$$

which is indeed the projection of \mathbf{u} onto \mathcal{X} , i.e. $\mathcal{S}_{\mathcal{N}_{\mathcal{X}}(\cdot)}(\mathbf{u}) = \mathbf{x} = \pi_{\mathcal{X}}(\mathbf{u})$.



From equations to inclusions

Before presenting methods for solving (2), we review a notion in **convex analysis** called **inclusion**. Let us motivate this concept by starting from a **system of equations**.

For a **single-valued** mapping $J : \mathbb{R}^p \rightarrow \mathbb{R}^p$, we consider the **system of equations**:

$$\mathbf{x} = J(\mathbf{x}). \quad (8)$$

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Fixed point iteration scheme

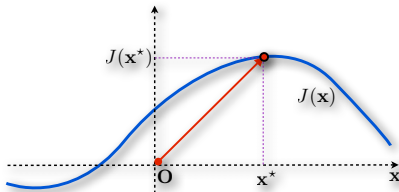
If the mapping J is **contractive**, i.e., $\exists \kappa \in [0, 1)$ such that

$$\|J(\mathbf{x}) - J(\hat{\mathbf{x}})\| \leq \kappa \|\mathbf{x} - \hat{\mathbf{x}}\|, \quad \forall \mathbf{x}, \hat{\mathbf{x}},$$

then, by the **Banach contraction mapping principle**, the sequence $\{\mathbf{x}^k\}$ generated by

$$\mathbf{x}^{k+1} := J(\mathbf{x}^k), \quad k \geq 0,$$

starting from \mathbf{x}^0 **converges** to the **unique fixed point** \mathbf{x}^* of J , i.e., $\mathbf{x}^* = J(\mathbf{x}^*)$.



From equations ...

Now, given a **single-valued** mapping $T : \mathbb{R}^p \rightarrow \mathbb{R}^p$, let us consider a **general system of equations**:

$$T(\mathbf{x}) = 0, \quad (9)$$

A simple way to transform (9) into (8), i.e., $\mathbf{x} = J(\mathbf{x})$, is:

$$\mathbf{x} = \mathbf{x} - \gamma T(\mathbf{x}) := J_T^\gamma(\mathbf{x}), \quad \gamma \neq 0.$$

Then, we can generate a **fixed-point scheme** for solving (9) as

$$\mathbf{x}^{k+1} := J_T^\gamma(\mathbf{x}^k). \quad (10)$$

If J_T^γ is **contractive** for given $\gamma \neq 0$, then $\{\mathbf{x}^k\}$ generated by (10) **converges** to a **solution** \mathbf{x}^* of (8).

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Example (Gradient method)

Let us consider the **optimality condition** of the **unconstrained smooth convex** problem:
 $f^* = \min_{\mathbf{x} \in \mathbb{R}^p} f(\mathbf{x})$ as

$$\nabla f(\mathbf{x}) = 0.$$

Since ∇f is **single-valued**, using the same trick as in the previous slide, we can write the fixed-point **iterative scheme** as

$$\mathbf{x}^{k+1} := \mathbf{x}^k - \gamma_k \nabla f(\mathbf{x}^k) = J_T^{\gamma_k}(\mathbf{x}^k).$$

If f is μ -**strongly convex** with $\mu > 0$ and ∇f is L_f -**Lipschitz continuous**, then J_T^γ is **contractive** for any $\gamma \in (0, 2\mu/L^2)$.

Proof of contractivity of J_T^γ

Proof.

1. If f is **smooth** and **strongly convex** with the strong convexity parameter μ . Then

$$(\nabla f(\mathbf{x}) - \nabla f(\hat{\mathbf{x}}))^T (\mathbf{x} - \hat{\mathbf{x}}) \geq \mu \|\mathbf{x} - \hat{\mathbf{x}}\|^2, \quad \forall \mathbf{x}, \hat{\mathbf{x}} \in \mathbb{R}^p.$$

2. Since $J_T^\gamma(\mathbf{x}) := \mathbf{x} - \gamma \nabla f(\mathbf{x})$, for any $\mathbf{x}, \hat{\mathbf{x}} \in \mathbb{R}^p$, we have

$$\begin{aligned} \|J_T^\gamma(\mathbf{x}) - J_T^\gamma(\hat{\mathbf{x}})\|^2 &= \|\mathbf{x} - \hat{\mathbf{x}} - \gamma(\nabla f(\mathbf{x}) - \nabla f(\hat{\mathbf{x}}))\|^2 \\ &= \|\mathbf{x} - \hat{\mathbf{x}}\|^2 - 2\gamma(\mathbf{x} - \hat{\mathbf{x}})^T (\nabla f(\mathbf{x}) - \nabla f(\hat{\mathbf{x}})) \\ &\quad + \gamma^2 \|\nabla f(\mathbf{x}) - \nabla f(\hat{\mathbf{x}})\|^2 \\ &\leq \|\mathbf{x} - \hat{\mathbf{x}}\|^2 - 2\gamma \underbrace{\mu \|\mathbf{x} - \hat{\mathbf{x}}\|^2}_{\text{strong convexity}} + \gamma^2 \underbrace{L^2 \|\mathbf{x} - \hat{\mathbf{x}}\|^2}_{\text{Lipschitz gradient}} \\ &\leq (1 - 2\gamma\mu + \gamma^2 L^2) \|\mathbf{x} - \hat{\mathbf{x}}\|^2. \end{aligned}$$

3. Hence, J_T^γ is **contractive** if $1 - 2\gamma\mu + \gamma^2 L^2 \in [0, 1)$, which implies $\gamma \in (0, 2\mu/L^2)$.
4. We note that $1 - 2\gamma\mu + \gamma^2 L^2$ is **minimized** if $\gamma_\star = \mu/L^2$ and hence $1 - 2\gamma\mu + \gamma^2 L^2 = 1 - \mu^2/L^2$.

□

... To inclusions

In the example presented previously, if we **no longer** assume f to be **smooth**, then the **optimality condition** turns out to be

$$0 \in \partial f(\mathbf{x}).$$

Since ∂f is a **set-valued mapping**, this condition is called an **inclusion**.

We can generalize this **inclusion** to any **set-valued mapping** \mathcal{T} from \mathbb{R}^p to $2^{\mathbb{R}^p}$ as

$$0 \in \mathcal{T}(\mathbf{x}). \tag{11}$$

In general, **solving the inclusion** (11) is **much more difficulty** than **solving the equation system** $T(\mathbf{x}) = 0$. Methods for solving (11) on the one hand can **inherit** from methods of solving equations, but on the other hand, require **new mathematical tools**.

... To inclusions

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We can generalize this **inclusion** to any **set-valued mapping** \mathcal{T} from \mathbb{R}^p to $2^{\mathbb{R}^p}$ as

$$0 \in \mathcal{T}(\mathbf{x}). \quad (11)$$

In general, **solving the inclusion** (11) is **much more difficulty** than **solving the equation system** $T(\mathbf{x}) = 0$. Methods for solving (11) on the one hand can **inherit** from methods of solving equations, but on the other hand, require **new mathematical tools**.

Additional mathematical tools

Since we are working with the possibly **nonsmooth** and **constrained convex** problem (2), where \mathcal{X} is **not specified**, its **optimality condition** will be reformulated as an **inclusion**. We will use **additional mathematical tools** from **variational analysis** such as:

- ▶ Monotone inclusions
- ▶ Monotone mixed variational inequalities
- ▶ Gap functions

Mixed variational inequality (MVI) formulation

Primal-dual mapping

We introduce a new primal-dual variable $\mathbf{z} := (\mathbf{x}^T, \lambda^T)^T \in \mathbb{R}^{p+n}$ and two mappings:

$$M(\mathbf{z}) := \begin{bmatrix} \mathbf{A}^T \lambda \\ \mathbf{b} - \mathbf{A} \mathbf{x} \end{bmatrix} \quad \text{and} \quad \mathcal{T}(\mathbf{z}) := \left\{ \begin{pmatrix} \xi \\ 0^n \end{pmatrix} \in \mathbb{R}^{p+n} : \xi \in \partial f(\mathbf{x}) \right\}. \quad (12)$$

- ▶ Then $M : \mathbb{R}^{p+n} \rightarrow \mathbb{R}^{p+n}$ is a **single-valued mapping** (linear mapping).
- ▶ If f is not differentiable, then $\mathcal{T} : \mathbb{R}^{p+n} \rightrightarrows \mathbb{R}^{p+n}$ is a **set-valued mapping**.

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Inclusion and MVI formulation

- ▶ The **optimality condition** (5) can be written as an **inclusion**:

$$0 \in \mathcal{P}(\mathbf{z}) := M(\mathbf{z}) + \mathcal{T}(\mathbf{z}).$$

- ▶ (5) can also be expressed as a **mixed variational inequality** (MVI):

$$f(\mathbf{x}) - f(\mathbf{x}^*) + M(\mathbf{z}^*)^T (\mathbf{z} - \mathbf{z}^*) \geq 0, \quad \forall \mathbf{z} \in \mathbb{R}^{p+n}. \quad (13)$$

Optimality condition as a monotone inclusion/VI problem

Lemma (Monotonicity of primal-dual mapping)

The mapping M and \mathcal{T} defined in (12):

$$M(\mathbf{z}) := \begin{bmatrix} \mathbf{A}^T \lambda \\ \mathbf{b} - \mathbf{A}\mathbf{x} \end{bmatrix} \quad \text{and} \quad \mathcal{T}(\mathbf{z}) := \left\{ \begin{pmatrix} \xi \\ 0^n \end{pmatrix} \in \mathbb{R}^{p+n} : \xi \in \partial f(\mathbf{x}) \right\}.$$

are **maximal monotone**. Consequently, $\mathcal{P} := M + \mathcal{T}$ is also **maximal monotone**.

To show the monotonicity of M , we can write $M(\mathbf{z})$ as

$$M(\mathbf{z}) := \mathbf{H}\mathbf{z} + \mathbf{h} \equiv \begin{bmatrix} 0 & \mathbf{A}^T \\ -\mathbf{A} & 0 \end{bmatrix} \mathbf{z} + \begin{bmatrix} 0 \\ \mathbf{b} \end{bmatrix}$$

It is clear that

$$\mathbf{z}^T \mathbf{H}\mathbf{z} = \mathbf{x}^T (\mathbf{A}^T \lambda) - \lambda^T \mathbf{A}\mathbf{x} = 0,$$

which shows that M is monotone.

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$$\mathbf{z}^T \mathbf{H}\mathbf{z} = \mathbf{x}^T (\mathbf{A}^T \lambda) - \lambda^T \mathbf{A}\mathbf{x} = 0,$$

which shows that M is monotone.

Monotone inclusion and monotone variational inequality

- ▶ The **inclusion** $0 \in \mathcal{P}(\mathbf{z}^*)$ is called a **monotone inclusion** if \mathcal{P} is **maximal monotone**.
- ▶ The **mixed variational inequality**

$$f(\mathbf{x}) - f(\mathbf{x}^*) + M(\mathbf{z}^*)^T (\mathbf{z} - \mathbf{z}^*) \geq 0, \quad \forall \mathbf{z} \in \mathbb{R}^{p+n}$$

is called **monotone** if f is **proper, closed and convex** and M is **maximal monotone**.

Gap function for the MVI problem

Gap function

Let us consider a **monotone MVI** problem of the form:

$$f(\mathbf{x}) - f(\mathbf{x}^*) + M(\mathbf{z}^*)^T(\mathbf{z} - \mathbf{z}^*) \geq 0, \quad \forall \mathbf{z} \in \mathbb{R}^{p+n}$$

The **gap function** associated with this problem is defined as follows

$$G(\mathbf{z}) := \max_{\hat{\mathbf{z}} \in \mathbb{R}^{p+n}} \left\{ f(\mathbf{x}) - f(\hat{\mathbf{x}}) + M(\mathbf{z})^T(\mathbf{z} - \hat{\mathbf{z}}) \right\}. \quad (14)$$

Gap function for the MVI problem

Gap function

Let us consider a **monotone MVIP** problem of the form:

$$f(\mathbf{x}) - f(\mathbf{x}^*) + M(\mathbf{z}^*)^T(\mathbf{z} - \mathbf{z}^*) \geq 0, \quad \forall \mathbf{z} \in \mathbb{R}^{p+n}$$

The **gap function** associated with this problem is defined as follows

$$G(\mathbf{z}) := \max_{\hat{\mathbf{z}} \in \mathbb{R}^{p+n}} \left\{ f(\mathbf{x}) - f(\hat{\mathbf{x}}) + M(\mathbf{z})^T(\mathbf{z} - \hat{\mathbf{z}}) \right\}. \quad (14)$$

Properties

- ▶ Computing G and its **gradient** require to solve the **convex problem** in (14).
- ▶ G is **nonnegative**, i.e.: $G(\mathbf{z}) \geq 0$ for all $\mathbf{z} \in \mathbb{R}^{p+n}$.
 - ▶ Indeed, we have
$$G(\mathbf{z}) = \max_{\hat{\mathbf{z}} \in \mathbb{R}^{p+n}} \left\{ f(\mathbf{x}) - f(\hat{\mathbf{x}}) + M(\mathbf{z})^T(\mathbf{z} - \hat{\mathbf{z}}) \right\} \geq f(\mathbf{x}) - f(\mathbf{x}) + M(\mathbf{z})^T(\mathbf{z} - \mathbf{z}) = 0,$$
with $\hat{\mathbf{z}} = \mathbf{z}$.
- ▶ $G(\mathbf{z}^*) = 0$ iff \mathbf{z}^* is a **KKT point** of (5).

Exercise: By writing the optimality condition of the maximization problem (14) and rearranging it, we obtain exactly the KKT condition (5).

Example: Gap function

Consider the following constrained convex problem:

$$\begin{array}{ll}\min_{\mathbf{x} \in \mathbb{R}^2} & \{f(\mathbf{x}) := |x_1| + x_2^2\}, \\ \text{s.t.} & x_1 + x_2 = 1.\end{array}$$

We have $\mathbf{z} = (x_1, x_2, \lambda)^T \in \mathbb{R}^3$ and $M(\mathbf{z}) := (\lambda, \lambda, 1 - x_1 - x_2)^T$.

The **gap function** associated with the **optimality condition** of this problem becomes:

$$\begin{aligned}G(\mathbf{z}) &:= \max_{(y_1, y_2, \eta)^T \in \mathbb{R}^3} \left\{ |x_1| + x_2^2 - |y_1| - y_2^2 + \eta(x_1 + x_2 - 1) - \lambda(y_1 + y_2 - 1) \right\} \\ &= \max_{(y_1, y_2)^T \in \mathbb{R}^2} \left\{ -|y_1| - y_2^2 - \lambda(y_1 + y_2 - 1) \right\} + \max_{\eta \in \mathbb{R}} \left\{ |x_1| + x_2^2 + \eta(x_1 + x_2 - 1) \right\} \\ &= \begin{cases} |x_1| + x_2^2 - d(\lambda) & \text{if } x_1 + x_2 - 1, \\ +\infty & \text{otherwise,} \end{cases}\end{aligned}$$

where $d(\lambda) := \min_{(y_1, y_2)^T \in \mathbb{R}^2} \left\{ |y_1| + y_2^2 + \lambda(y_1 + y_2 - 1) \right\}$ is the dual function.

Outline

► Today

1. Convex constrained optimization and motivating examples
2. Optimality condition
3. Conjugate functions
4. Monotone inclusion and monotone mixed variational inequality formulations
5. Chambolle-Pock's primal-dual method
6. Primal-dual hybrid gradient method
7. Splitting methods
8. Model-based excessive gap primal-dual method

► Next week

1. Disciplined convex programming

A special class of constrained convex problems (2)

We first consider a constrained reformulation of **composite convex minimization** problems considered in Lecture 5.

Constrained convex reformulation

We consider the following special case of (2):

$$\begin{cases} \min_{\mathbf{x} := (\mathbf{u}, \mathbf{v}) \in \mathbb{R}^{p_1 + p_2}} & \{ F(\mathbf{x}) := f(\mathbf{u}) + g(\mathbf{v}) \} \\ \text{s.t.} & \mathbf{K}\mathbf{u} - \mathbf{v} = 0. \end{cases} \quad (15)$$

where \mathbf{K} is a **linear operator**, $f \in \mathcal{F}(\mathbb{R}^{p_1})$ and $g \in \mathcal{F}(\mathbb{R}^{p_2})$ are two **convex functions**.

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where \mathbf{K} is a **linear operator**, $f \in \mathcal{F}(\mathbb{R}^{p_1})$ and $g \in \mathcal{F}(\mathbb{R}^{p_2})$ are two **convex functions**.

- By setting $\mathbf{A} := [\mathbf{K}, -\mathbf{I}]$ and $\mathbf{b} := \mathbf{0}^n$, we can formulate (15) into (2):

$$F^* := \begin{cases} \min_{\mathbf{x} \in \mathbb{R}^p} & F(\mathbf{x}) \\ \text{s.t.} & \mathbf{A}\mathbf{x} = \mathbf{b}. \end{cases}$$

- Problem (15) can be written as a composite convex minimization:

$$F^* := \min_{\mathbf{u} \in \mathbb{R}^{p_1}} \{ F(\mathbf{u}) := f(\mathbf{u}) + g(\mathbf{K}\mathbf{u}) \}. \quad (16)$$

Min-max formulation and dual problem

The min-max (saddle point) problem

By using the Fenchel conjugate g^* of g , we can write

$$g(\mathbf{K}\mathbf{u}) = \max_{\mathbf{v} \in \mathbb{R}^n} \{ \langle \mathbf{K}\mathbf{u}, \mathbf{v} \rangle - g^*(\mathbf{v}) \}.$$

Substituting this function into (16), we obtain:

$$\boxed{\min_{\mathbf{u} \in \mathbb{R}^{p_1}} \max_{\mathbf{v} \in \mathbb{R}^n} \{ \langle \mathbf{K}\mathbf{u}, \mathbf{v} \rangle + f(\mathbf{u}) - g^*(\mathbf{v}) \}} \quad (17)$$

where g^* is the conjugate of g .

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Dual problem

By exchanging the min-max in (17) and note that

$$\min_{\mathbf{u} \in \mathbb{R}^{p_1}} \{ \langle \mathbf{K}\mathbf{u}, \mathbf{v} \rangle + f(\mathbf{u}) \} = - \max_{\mathbf{u} \in \mathbb{R}^{p_1}} \{ \langle -\mathbf{K}^T \mathbf{v}, \mathbf{u} \rangle - f(\mathbf{u}) \} = -f^*(-\mathbf{K}^T \mathbf{v})$$

we have

$$\boxed{\max_{\mathbf{v} \in \mathbb{R}^q} \{ -f^*(-\mathbf{K}^T \mathbf{v}) - g^*(\mathbf{v}) \}}. \quad (18)$$

Chambolle-Pock's algorithm: the main idea

Optimality condition

First, we write the **optimality condition** of (17) as follows:

$$\begin{cases} \mathbf{K}\mathbf{u}^* & \in \partial g^*(\mathbf{v}^*) \\ -\mathbf{K}^T \mathbf{v}^* & \in \partial f(\mathbf{u}^*). \end{cases} \quad (19)$$

Since problem is **convex**, condition (19) is **necessary and sufficient** for $(\mathbf{u}^*, \mathbf{v}^*)$ to be **primal and dual optimal** to (17):

$$\min_{\mathbf{u} \in \mathbb{R}^p} \max_{\mathbf{v} \in \mathbb{R}^q} \left\{ \langle \mathbf{K}\mathbf{u}, \mathbf{v} \rangle + f(\mathbf{u}) - g^*(\mathbf{v}) \right\}.$$

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Fixed-point expression

Second, from (19), for any $\sigma > 0$ and $\tau > 0$, we can write

$$\mathbf{v}^* + \sigma \mathbf{K}\mathbf{u}^* \in (\mathbb{I} + \sigma \partial g^*)(\mathbf{v}^*) \text{ and } \mathbf{u}^* - \tau \mathbf{K}^T \mathbf{v}^* \in (\mathbb{I} + \tau \partial f)(\mathbf{u}^*).$$

Using the **proximal operator** of τf and σg^* , we can write the last expression as

$$\begin{cases} \mathbf{v}^* & = \text{prox}_{\sigma g^*}(\mathbf{v}^* + \sigma \mathbf{K}\mathbf{u}^*) \\ \mathbf{u}^* & = \text{prox}_{\tau f}(\mathbf{u}^* - \tau \mathbf{K}^T \mathbf{v}^*). \end{cases} \quad (20)$$

This relation shows that $\mathbf{x}^* := (\mathbf{u}^*, \mathbf{v}^*)$ is a **fixed point** of the **mapping** \mathcal{T} with:

$$\mathcal{T}(\mathbf{x}) := (\text{prox}_{\sigma g^*}(\mathbf{v} + \sigma \mathbf{K}\mathbf{u}), \text{prox}_{\tau f}(\mathbf{u} - \tau \mathbf{K}^T \mathbf{v})).$$

The Chambolle-Pock algorithm

The **Chambolle-Pock algorithm** is rooted from the [classical Arrow-Hurwicz method](#), which is based on the [fixed-point](#) expression (20).

Chambolle-Pock's algorithm (CPA) [2]

1. Choose $\tau > 0$, $\sigma > 0$, $\theta \in [0, 1]$, $\mathbf{u}^0 \in \mathbb{R}^p$ and $\mathbf{v}^0 \in \mathbb{R}^q$. Set $\hat{\mathbf{u}}^0 := \mathbf{u}^0$.
2. For $k = 0, 1, \dots$, perform:

$$\begin{cases} \mathbf{v}^{k+1} &:= \text{prox}_{\sigma g^*}(\mathbf{v}^k + \sigma \mathbf{K} \hat{\mathbf{u}}^k) \\ \mathbf{u}^{k+1} &:= \text{prox}_{\tau f}(\mathbf{u}^k - \tau \mathbf{K}^T \mathbf{v}^k) \\ \hat{\mathbf{u}}^{k+1} &:= \mathbf{u}^{k+1} + \theta(\mathbf{u}^{k+1} - \mathbf{u}^k). \end{cases} \quad (21)$$

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Remarks

- ▶ If $\theta = 0$, then $\hat{\mathbf{u}}^k = \mathbf{u}^k$ and (21) collapses to the [Arrow-Hurwicz method](#):

$$\begin{cases} \mathbf{v}^{k+1} &:= \text{prox}_{\sigma g^*}(\mathbf{v}^k + \sigma \mathbf{K} \mathbf{u}^k) \\ \mathbf{u}^{k+1} &:= \text{prox}_{\tau f}(\mathbf{u}^k - \tau \mathbf{K}^T \mathbf{v}^k). \end{cases}$$

- ▶ The [step sizes](#) σ and τ and the [parameter](#) θ can keep [constantly](#) or [adaptively updates](#).
- ▶ When $\mathbf{K} = \mathbf{I}$, the **Chambolle-Pock algorithm** is equivalent to [ADMM](#).

Restricted gap function for optimality certification

Let us define $\mathbf{x} := (\mathbf{u}, \mathbf{v}) \equiv (\mathbf{u}^T, \mathbf{v}^T)^T$, $\phi(\mathbf{x}) := f(\mathbf{u}) + g^*(\mathbf{v})$ and

$$M(\mathbf{x}) := \begin{bmatrix} -\mathbf{K}^T \mathbf{v} \\ \mathbf{K} \mathbf{u} \end{bmatrix}.$$

Then (19): $\begin{cases} \mathbf{K} \mathbf{u}^* \in \partial g^*(\mathbf{v}^*) \\ -\mathbf{K}^T \mathbf{v}^* \in \partial f(\mathbf{u}^*) \end{cases}$ can be written as a **monotone MVI problem**:

$$\boxed{\phi(\mathbf{x}) - \phi(\mathbf{x}^*) + M(\mathbf{x}^*)^T (\mathbf{x} - \mathbf{x}^*) \geq 0, \quad \forall \mathbf{x} \in \mathbb{R}^{p_1+p_2}} \quad (22)$$

where $\mathbf{x}^* := (\mathbf{u}^*, \mathbf{v}^*)$.

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$$\phi(\mathbf{x}) - \phi(\mathbf{x}^*) + M(\mathbf{x}^*)^T (\mathbf{x} - \mathbf{x}^*) \geq 0, \quad \forall \mathbf{x} \in \mathbb{R}^{p_1+p_2} \quad (22)$$

where $\mathbf{x}^* := (\mathbf{u}^*, \mathbf{v}^*)$.

Definition (Restricted gap function)

Let $\mathcal{X} \subseteq \mathbb{R}^{p_1+p_2}$ be a **nonempty, closed, convex and bounded** set. We define a **restricted gap function** of (22) restricted on \mathcal{X} as

$$G_{\mathcal{X}}(\mathbf{x}) := \max_{\tilde{\mathbf{x}} \in \mathcal{X} \subseteq \mathbb{R}^{p_1+p_2}} \left\{ \phi(\mathbf{x}) - \phi(\tilde{\mathbf{x}}) + M(\mathbf{x})^T (\mathbf{x} - \tilde{\mathbf{x}}) \right\} \quad (23)$$

Convergence theorem

Theorem (Convergence)

Assumptions:

- ▶ (17): $\min_{\mathbf{u} \in \mathbb{R}^{p_1}} \max_{\mathbf{v} \in \mathbb{R}^{p_2}} \left\{ \langle \mathbf{K}\mathbf{u}, \mathbf{v} \rangle + f(\mathbf{u}) - g^*(\mathbf{v}) \right\}$ has a **saddle point** $\mathbf{z}^* := (\mathbf{u}^*, \mathbf{v}^*)$.
- ▶ $\{(\mathbf{u}^k, \mathbf{v}^k)\}$ be the sequence generated by the **Chambolle-Pock algorithm**.
- ▶ If we choose $\theta := 1$, $\sigma > 0$ and $\tau > 0$ such that $\tau\sigma\|\mathbf{K}\|^2 < 1$.

Conclusion:

- ▶ The sequence $\{\bar{\mathbf{x}}^k\}_{k \geq 0}$ defined by

$$\bar{\mathbf{x}}^k = (\bar{\mathbf{u}}^k, \bar{\mathbf{v}}^k) := \frac{1}{(k+1)} \sum_{j=0}^k (\mathbf{u}^j, \mathbf{v}^j)$$

satisfies

$$G_{\mathcal{X}}(\bar{\mathbf{x}}^k) \leq \frac{1}{k+1} \left[\max_{\mathbf{z} := (\mathbf{u}, \mathbf{v}) \in \mathcal{X}} \left\{ (1/(2\tau)) \|\mathbf{u} - \mathbf{u}^0\|^2 + (1/(2\sigma)) \|\mathbf{v} - \mathbf{v}^0\|^2 \right\} \right]. \quad (24)$$

- ▶ $\{\bar{\mathbf{z}}^k\}_{k \geq 0}$ **converges** to a **saddle point** \mathbf{z}^* of (17) at the $\mathcal{O}(1/k)$ **rate** w.r.t. the **restricted gap function** $G_{\mathcal{X}}$ (in the **ergodic** sense).

Convergence theorem: Remarks

$$G_{\mathcal{X}}(\bar{\mathbf{x}}^k) \leq \frac{1}{k+1} \left[\max_{\mathbf{x} := (\mathbf{u}, \mathbf{v}) \in \mathcal{X}} \left\{ (1/(2\tau)) \|\mathbf{u} - \mathbf{u}^0\|^2 + (1/(2\sigma)) \|\mathbf{v} - \mathbf{v}^0\|^2 \right\} \right]. \quad (24)$$

- ▶ The right-hand side of the estimate (24) depends on the choice of \mathcal{X} .
Theoretically, we need to choose \mathcal{X} such that $\mathcal{X}^* \subseteq \mathcal{X}$, where \mathcal{X}^* is the solution set of (17), which is **unknown**.
- ▶ The estimate (24) does not imply the convergence rate of $\{\bar{\mathbf{u}}^k\}$ and $\{\bar{\mathbf{v}}^k\}$ **separately**.

Acceleration - Case 1: f or g^* is strongly convex

If either f or g^* is **strongly convex**, the **Chambolle-Pock algorithm** can be **accelerated** to get **better convergence rate**.

Assumption A.1.

The function f is **strongly convex** with a strong convexity parameter $\sigma_f > 0$.

- ▶ Under Assumption A.1., the conjugate f^* is **smooth** and has the **Lipschitz gradient**.
- ▶ The **Chambolle-Pock algorithm** can be **accelerated** to get $\mathcal{O}(1/k^2)$ **convergence rate**.

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Chambolle-Pock's algorithm for strongly convex (CPA₁)

1. Choose $\tau > 0$, $\sigma > 0$ such that $\sigma_0 \tau_0 L^2 \leq 1$, $\mathbf{u}^0 \in \mathbb{R}^p$ and $\mathbf{v}^0 \in \mathbb{R}^q$. Set $\hat{\mathbf{u}}^0 := \mathbf{u}^0$.
2. For $k = 0, 1, \dots$, perform:

$$\begin{cases} \mathbf{v}^{k+1} &:= \text{prox}_{\sigma_k g^*}(\mathbf{v}^k + \sigma_k \mathbf{K} \hat{\mathbf{u}}^k) \\ \mathbf{u}^{k+1} &:= \text{prox}_{\tau_k f}(\mathbf{u}^k - \tau_k \mathbf{K}^T \mathbf{v}^k) \\ \theta_k &:= (1 + 2\sigma_f \tau_k)^{-1/2}, \\ \tau_{k+1} &:= \theta_k \tau_k, \\ \sigma_{k+1} &:= \theta_k^{-1} \sigma_k, \\ \hat{\mathbf{u}}^{k+1} &:= \mathbf{u}^{k+1} + \theta_k (\mathbf{u}^{k+1} - \mathbf{u}^k). \end{cases} \quad (25)$$

Acceleration - Case 2: Both f and g^* are strongly convex

If both f and g^* are **strongly convex**, we can **accelerate** the **Chambolle-Pock algorithm** to obtain the **linear convergence rate**.

Assumption A.2.

The functions f and g^* are strongly convex with strong convexity parameters $\sigma_f > 0$ and $\sigma_{g^*} > 0$, respectively.

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If both f and g^* are **strongly convex**, we can **accelerate** the **Chambolle-Pock algorithm** to obtain the **linear convergence rate**.

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The functions f and g^* are strongly convex with a strong convexity parameters $\sigma_f > 0$ and $\sigma_{g^*} > 0$, respectively.

Chambolle-Pock's algorithm for strongly convex f and g^* (CPA₂)

1. Choose $\omega \leq 2 \sqrt{\sigma_f \sigma_{g^*}} L$.
2. Set $\tau := \omega / (2\sigma_f)$, $\sigma := \omega / (2\sigma_{g^*})$ and choose $\theta \in [(1 + \omega)^{-1}, 1]$.
3. Find $\mathbf{u}^0 \in \mathbb{R}^p$ and $\mathbf{v}^0 \in \mathbb{R}^q$. Set $\hat{\mathbf{u}}^0 := \mathbf{u}^0$.
4. For $k = 0, 1, \dots$, perform:

$$\begin{cases} \mathbf{v}^{k+1} &:= \text{prox}_{\sigma g^*}(\mathbf{v}^k + \sigma \mathbf{K} \hat{\mathbf{u}}^k) \\ \mathbf{u}^{k+1} &:= \text{prox}_{\tau f}(\mathbf{u}^k - \tau \mathbf{K}^T \mathbf{v}^k) \\ \hat{\mathbf{u}}^{k+1} &:= \mathbf{u}^{k+1} + \theta(\mathbf{u}^{k+1} - \mathbf{u}^k). \end{cases} \quad (26)$$

Example: Strong convexity of f/g^*

Example (Strong convexity of f)

We consider the following **image ℓ_1 -TV denoising** problem:

$$\min_{\mathbf{u} \in \mathbb{R}^{p_1}} (1/2) \|\mathbf{u} - \mathbf{b}\|_F^2 + \rho \|\mathbf{D}\mathbf{u}\|_1. \quad (27)$$

Here \mathbf{b} is a **noisy image**, $\rho > 0$ is a **regularization parameter**, and \mathbf{D} is a given matrix. We can write this problem into the following **minmax** form:

$$\min_{\mathbf{u} \in \mathbb{R}^{p_1}} \max_{\|\mathbf{v}\|_\infty \leq 1} \left\{ \rho \mathbf{v}^T \mathbf{D}\mathbf{u} + (1/2) \|\mathbf{u} - \mathbf{b}\|_F^2 \right\}.$$

In this case, we have $f(\mathbf{u}) := (1/2) \|\mathbf{u} - \mathbf{b}\|_F^2$, which is **strongly convex** with the parameter $\mu_f = 1$, and $g^*(\mathbf{v}) = 0$.

Example (Strong convexity of both f and g^*)

If we apply **Nesterov's smoothing** technique to (27) with a **simple prox-function** $(1/2) \|\mathbf{v}\|_F^2$, we obtain the following problem:

$$\min_{\mathbf{u} \in \mathbb{R}^{p_1}} \max_{\|\mathbf{v}\|_\infty \leq 1} \left\{ \rho \mathbf{v}^T \mathbf{D}\mathbf{u} + (1/2) \|\mathbf{u} - \mathbf{b}\|_F^2 - (\gamma/2) \|\mathbf{v}\|_F^2 \right\}.$$

where $\gamma > 0$ is a smoothness parameter. In this case, we can denote $g^*(\mathbf{v}) := (\gamma/2) \|\mathbf{v}\|_F^2$, which is **strongly convex** with the parameter $\gamma > 0$.

Convergence of CPA₁ and CPA₂

► **Assumptions:**

- f is **strongly convex** with a **strong convexity parameter** $\sigma_f > 0$.
- Let $\{(\mathbf{u}^k, \mathbf{v}^k)\}_{k \geq 0}$ be the sequence generated by CPA₁.
- Let $\tau_0 > 0$ and $\sigma_0 := 1/(\tau_0 L^2)$.

Conclusion: Then for any $\epsilon > 0$, there exists K_0 (depending on ϵ and $\sigma_f \tau_0$) such that for any $k \geq K_0$,

$$\|\bar{\mathbf{u}}^k - \mathbf{u}^*\|^2 \leq \frac{1 + \epsilon}{(k + 1)^2} \left(\frac{\|\mathbf{u}^0 - \mathbf{u}^*\|^2}{\sigma_f^2 \tau_0^2} + \frac{L^2}{\sigma_f^2} \|\mathbf{v}^0 - \mathbf{v}^*\|^2 \right),$$

where $\bar{\mathbf{u}}^k := (k + 1)^{-1} \sum_{j=0}^k \mathbf{u}^j$. The sequence $\{\bar{\mathbf{u}}^k\}_{k \geq 0}$ converges to \mathbf{u}^* at the $\mathcal{O}(1/k^2)$ rate.

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► Assumptions:

- f and g^* are **strongly convex** with **strong convexity parameters** $\sigma_f > 0$ and $\sigma_{g^*} > 0$, respectively.
- Let $\{(\mathbf{u}^k, \mathbf{v}^k)\}_{k \geq 0}$ be the sequence generated by CPA₂.
- Let

$$c := \frac{1 + \theta}{2(1 + \sqrt{\sigma_f \sigma_{g^*}}/L)} < 1.$$

Conclusion: Then $\{(\mathbf{u}^k, \mathbf{v}^k)\}_{k \geq 0}$ converges to $(\mathbf{u}^*, \mathbf{v}^*)$ at **linear rate** $\mathcal{O}(c^k)$.

Example: Image inpainting

Mathematical formulation

Given a damaged image $\mathbf{b} \in \mathbb{R}^{m \times n}$, where the missed pixels b_{ij} are in certain region, i.e., $(i, j) \in \mathcal{M} \subset \mathcal{I} := \{1, \dots, m\} \times \{1, \dots, n\}$. The aim is to recover a **undamaged image** \mathbf{x} by using the **total variation operator**. This problem can be formulated as:

$$\min_{\mathbf{u} \in \mathbb{R}^{m \times n}} \|\mathbf{K}\mathbf{u}\|_1 + (\rho/2) \sum_{(i,j) \in \mathcal{I} \setminus \mathcal{M}} (u_{ij} - b_{ij})^2 \quad (28)$$

where $\rho > 0$ is a **regularization parameter** and \mathbf{K} is the total variation linear transform.

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How to apply the Chambolle-Pock algorithm?

- ▶ This problem is of the form $F^* := \min_{\mathbf{u} \in \mathbb{R}^{p_1}} \{F(\mathbf{u}) := f(\mathbf{u}) + g(\mathbf{K}\mathbf{u})\}$, where $f(\mathbf{u}) := (\rho/2) \sum_{(i,j) \in \mathcal{I} \setminus \mathcal{M}} (u_{ij} - b_{ij})^2$ and $g(\mathbf{v}) := \|\mathbf{v}\|_1$.
- ▶ Both f and g have closed form prox-operators:

$$\begin{aligned} \text{prox}_{\sigma g^*}(\mathbf{v}) &= \mathbf{v} ./ \max(1, |\mathbf{v}|) \\ \text{prox}_{\tau f}(\mathbf{u}) &= \begin{cases} u_{ij} & \text{if } (i, j) \in \mathcal{M} \\ \frac{u_{ij} + \tau \rho b_{ij}}{1 + \tau \rho} & \text{otherwise} \end{cases} \end{aligned}$$

Example: Image inpainting - configuration

We implement Chambolle-Pock's algorithm for solving the inpainting problem (28) using the following configurations:

- ▶ **Parameter selection:**

- ▶ $\sigma = 10$, $\tau = 0.01125$ and $\theta = 1$.
- ▶ The initial point $\mathbf{u}^0 := \mathbf{b}$ and $\mathbf{v}^0 := 0$.
- ▶ The tolerance $\epsilon = 10^{-5}$.

- ▶ **Stopping criterion:**

$$\|\mathbf{u}^{k+1} - \mathbf{u}^k\|_F \leq \epsilon \|\mathbf{u}^k\|_F.$$

- ▶ **Data generating:**

- ▶ We take a real gray image of size 255×255
- ▶ The image is damaged by a mask of 30 lines crossing from the left to the right.
- ▶ The regularization parameter ρ is chosen as $\rho = 1$ and $\rho = 0.75$ for two cases.

Convergence behavior: Left: $\rho = 1$ – Right: $\rho = 0.75$.

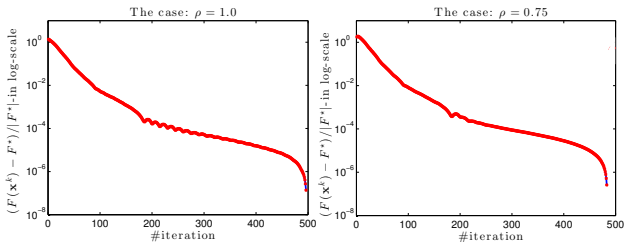


Image inpainting: outputs

Original image



Damaged image

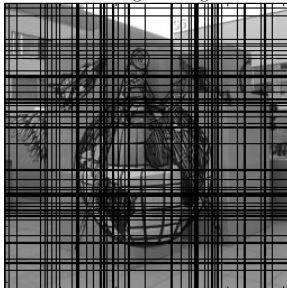
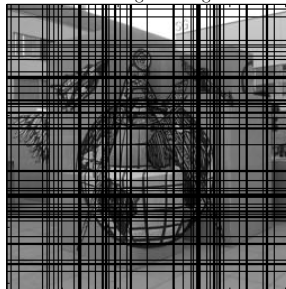


Image inpainting: outputs

Original image



Damaged image



- ▶ The objective value: 9.1323×10^3

- ▶ Relative error:

$$\|\mathbf{x}^k - \mathbf{x}^b\| / \|\mathbf{x}^b\| = 0.076398$$

where \mathbf{x}^b is the original image

- ▶ The number of iterations: 497
- ▶ The CPU time: 2.864s.

- ▶ The objective value: 5.9100×10^3

- ▶ Relative error:

$$\|\mathbf{x}^k - \mathbf{x}^b\| / \|\mathbf{x}^b\| = 0.078396$$

where \mathbf{x}^b is the original image

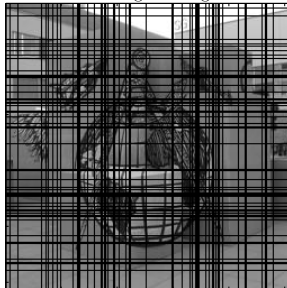
- ▶ The number of iterations: 484
- ▶ The CPU time: 2.707s.

Image inpainting: outputs

Original image



Damaged image



Recovered image ($\rho = 1$)



Recovered image ($\rho = 0.75$)



Outline

► Today

1. Convex constrained optimization and motivating examples
2. Optimality condition
3. Conjugate functions
4. Monotone inclusion and monotone mixed variational inequality formulations
5. Chambolle-Pock's primal-dual method
6. Primal-dual hybrid gradient method
7. Splitting methods
8. Model-based excessive gap primal-dual method

► Next week

1. Disciplined convex programming

Splitting methods

- ▶ **Splitting methods** have been widely used to solve **monotone inclusions** involving the **sum of two maximal monotone operators**.
- ▶ They can be used to solve the **constrained convex optimization problem** (2).

From the **first line** of (19) we have $\mathbf{v}^* \in \partial g(\mathbf{K}\mathbf{u}^*)$. Plug this into the **second line** of (19) to get

$$\boxed{0 \in \partial f(\mathbf{u}^*) + \mathbf{K}^T \partial g(\mathbf{K}\mathbf{u}^*)} \quad (29)$$

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Splitting monotone inclusion

Assumptions:

- ▶ Define $A(\mathbf{u}) := \mathbf{K}^T \partial g(\mathbf{K}\mathbf{u})$ and $B(\mathbf{u}) := \partial f(\mathbf{u})$.
- ▶ Assume that \mathbf{K} is **full rank**.

Conclusion:

- ▶ A and B are two **maximal monotone operators**.
- ▶ (29) can be expressed as:

$$0 \in A(\mathbf{u}^*) + B(\mathbf{u}^*) \quad (30)$$

Splitting methods

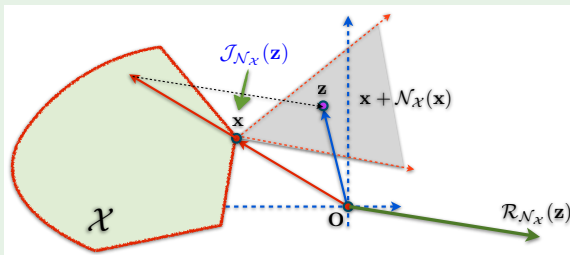
- ▶ **Splitting methods** have been widely used to solve **monotone inclusions** involving the **sum of two maximal monotone operators**.
- ▶ They can be used to solve the **constrained convex optimization problem** (2).

Reflection operator of the resolvent

Let $\mathcal{J}_{\mathcal{T}}(\cdot) := (\mathbb{I} + \mathcal{T})^{-1}(\cdot)$ be the **resolvent** of a **maximal monotone operator** \mathcal{T} . We define the **reflection operator** of $\mathcal{J}_{\mathcal{T}}$ as

$$\mathcal{R}_{\mathcal{T}}(\mathbf{z}) := 2\mathcal{J}_{\mathcal{T}}(\mathbf{z}) - \mathbf{z}.$$

Example (The reflection operator of the normal cone $\mathcal{N}_{\mathcal{X}}$)



Deriving a fixed-point formulation

- ▶ A splitting method **generates** an **iterative sequence** $\{\mathbf{u}^k\}$ by using a **fixed-point derivation** of the inclusion (30): $0 \in A(\mathbf{u}) + B(\mathbf{u})$.
- ▶ In addition, it **splits** the **computations** such that one can exploit the **individual computations** of A and B **separately**.

Deriving a fixed-point formulation

Starting from $0 \in A(\mathbf{u}) + B(\mathbf{u})$, we can rewrite

$$2\mathbf{u} \in (\mathbb{I} + A)(\mathbf{u}) + (\mathbb{I} + B)(\mathbf{u}). \quad (31)$$

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Using the resolvent of B , we can express the last inclusion as

$$\mathbf{u} = \mathcal{J}_A(2\mathcal{J}_B(\mathbf{z}) - \mathbf{z}) = \mathcal{J}_A(\mathcal{R}_B(\mathbf{z})).$$

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Using the reflection operator of A we finally get

$$\mathbf{z} = \mathcal{R}_A(\mathcal{R}_B(\mathbf{z})).$$

Fixed-point characterization and splitting computation

Fixed-point formulation

- If \mathbf{u}^* is a solution of (30) (cf. $0 \in A(\mathbf{u}^*) + B(\mathbf{u}^*)$) then

$$\boxed{\mathbf{z}^* = \mathcal{R}_A(\mathcal{R}_B(\mathbf{z}^*)) \text{ and } \mathbf{u}^* = \mathcal{J}_B(\mathbf{z}^*)}. \quad (32)$$

- Alternatively, if \mathbf{u}^* is a solution of (30) then for any $\beta \neq 0$, we have

$$\boxed{\mathbf{z}^* = (1 - \beta)\mathbf{z}^* + \beta \mathcal{R}_A(\mathcal{R}_B(\mathbf{z}^*)) \text{ and } \mathbf{u}^* = \mathcal{J}_B(\mathbf{z}^*)}. \quad (33)$$

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Let assume that our iterative scheme is based on the fixed point formulation (33) as:

$$\mathbf{z}^{k+1} := (1 - \beta)\mathbf{z}^k + \beta \mathcal{R}_A(\mathcal{R}_B(\mathbf{z}^k)).$$

We can split this formula by using the definition of \mathcal{R} and \mathcal{J} .

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Splitting: Douglas-Rachford's method

Fixed-point iteration

The **Douglas-Rachford method** bases on the **fixed-point formulation** (33) to generate an **iterative sequence** as:

$$\mathbf{z}^{k+1} := (1 - \beta_k)\mathbf{z}^k + \beta_k \mathcal{R}_A \left(\mathcal{R}_B(\mathbf{z}^k) \right) \quad \text{and} \quad \mathbf{u}^k = \mathcal{J}_B(\mathbf{z}^k).$$

By **splitting the computation** as in the previous slide, we can summarize this scheme as:

$$\begin{cases} \mathbf{u}^k &:= \mathcal{J}_B(\mathbf{z}^k) \\ \mathbf{v}^k &:= \mathcal{J}_A(2\mathbf{u}^k - \mathbf{z}^k) \\ \mathbf{z}^{k+1} &:= \mathbf{z}^k + \eta_k(\mathbf{v}^k - \mathbf{u}^k) \end{cases}$$

for $\eta_k := 2\beta_k \neq 0$.

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$$\mathbf{z}^{k+1} := (1 - \beta_k)\mathbf{z}^k + \beta_k \mathcal{R}_A \left(\mathcal{R}_B(\mathbf{z}^k) \right) \quad \text{and} \quad \mathbf{u}^k = \mathcal{J}_B(\mathbf{z}^k).$$

By **splitting the computation** as in the previous slide, we can summarize this scheme as:

$$\begin{cases} \mathbf{u}^k &:= \mathcal{J}_B(\mathbf{z}^k) \\ \mathbf{v}^k &:= \mathcal{J}_A(2\mathbf{u}^k - \mathbf{z}^k) \\ \mathbf{z}^{k+1} &:= \mathbf{z}^k + \eta_k(\mathbf{v}^k - \mathbf{u}^k) \end{cases}$$

for $\eta_k := 2\beta_k \neq 0$.

Douglas-Rachford's algorithm (DRA)

1. Given $\mathbf{z}^0 \in \text{dom}(B)$ as an initial point and $\eta_0 \neq 0$.
2. For $k = 0, 1, \dots$, perform:

$$\begin{cases} \mathbf{u}^k &:= \mathcal{J}_B(\mathbf{z}^k) \\ \mathbf{v}^k &:= \mathcal{J}_A(2\mathbf{u}^k - \mathbf{z}^k) \\ \mathbf{z}^{k+1} &:= \mathbf{z}^k + \eta_k(\mathbf{v}^k - \mathbf{u}^k) \end{cases}$$

and update η_k if required.

Douglas-Rachford method for convex problem (15)

In order to apply DRA for solving (15), we define $A(\mathbf{u}) := \tau \mathbf{K}^T \partial g(\mathbf{K}\mathbf{u})$ and $B(\mathbf{u}) := \tau \partial f(\mathbf{u})$ for a given scaling factor $\tau > 0$.

Douglas-Rachford's method for solving (15)

1. Given $\mathbf{z}^0 \in \text{dom}(f)$ as an initial point. Choose $\tau_0 > 0$ and $\eta_0 > 0$.
2. For $k = 0, 1, \dots$, perform:

$$\begin{cases} \mathbf{u}^k &:= \text{prox}_{\tau_k f}(\mathbf{z}^k), \\ \mathbf{v}^k &:= \underset{\mathbf{v}}{\text{argmin}} \left\{ g(\mathbf{K}\mathbf{v}) + (1/(2\tau_k)) \|\mathbf{v} - 2\mathbf{u}^k + \mathbf{z}^k\|_2^2 \right\}, \\ \mathbf{z}^{k+1} &:= \mathbf{z}^k + \eta_k (\mathbf{v}^k - \mathbf{u}^k). \end{cases}$$

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and update τ_k and η_k if necessary.

Remark: Quadratic loss and diagonalizable operator

► **Assumptions:** $g(\mathbf{K}\mathbf{v}) := (1/2)\|\mathbf{K}\mathbf{v} - \mathbf{b}\|_2^2$ and $\mathbf{K}^T \mathbf{K} = \Sigma$, where Σ is **diagonal**.

► **Conclusion:** $\mathbf{v}^k := (\tau_k^{-1} \mathbb{I} + \Sigma)^{-1} (\mathbf{K}^T \mathbf{b} + \tau_k^{-1} \mathbf{s}^k)$.

► Let $\mathbf{s}^k := 2\mathbf{u}^k - \mathbf{z}^k$. We can write the optimality condition of \mathbf{v}^k as

$$(\tau_k^{-1} \mathbb{I} + \mathbf{K}^T \mathbf{K}) \mathbf{v}^k = \mathbf{K}^T \mathbf{b} + \tau_k^{-1} \mathbf{s}^k.$$

► Since $\mathbf{K}^T \mathbf{K} = \Sigma$, we can compute \mathbf{v}^k explicitly as

$$\mathbf{v}^k = (\tau_k^{-1} \mathbb{I} + \Sigma)^{-1} (\mathbf{K}^T \mathbf{b} + \tau_k^{-1} \mathbf{s}^k).$$

Convergence of splitting methods

Theorem (Convergence of Douglas-Rachford's method [4])

Assume that the **solution set** \mathcal{U}^* of (30) is **nonempty** and the sequence $\{\eta_k\}$ is chosen such that

$$\eta_k \in [0, 2] \quad \text{and} \quad \sum_{k=0}^{+\infty} \eta_k(2 - \eta_k) = \infty. \quad (34)$$

Then, the sequence $\{\mathbf{u}^k\}$ generated by the **Douglas-Rachford algorithm** converges to a **solution** \mathbf{u}^* in \mathcal{U}^* .

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Remarks

- ▶ We can choose η_k as a constant step, e.g., $\eta_k = 1$ for $k \geq 0$.
- ▶ If $\eta_k = 1$ for all $k \geq 0$, then **Douglas-Rachford's algorithm** coincides with **Peaceman-Rachford's method** [10].
- ▶ When **Douglas-Rachford's algorithm** is **applied to solve (15)**, it coincides with **ADMM** [3].

Alternative derivation

Assumptions

- ▶ The **solution set** \mathcal{U}^* of $0 \in A(\mathbf{u}) + B(\mathbf{u})$ is **nonempty**
- ▶ For simplicity of discussion, we assume that B is **single-valued** and $\beta = 1/2$.

By using the definition of $\mathcal{R}_A = 2\mathcal{J}_A - \mathbb{I}$ and $\mathcal{R}_B = 2\mathcal{J}_B - \mathbb{I}$, the **Douglas-Rachford** iterative scheme

$$\begin{cases} \mathbf{u}^k &:= \mathcal{J}_B(\mathbf{z}^k), \\ \mathbf{z}^{k+1} &:= (1/2)\mathbf{z}^k + (1/2)\mathcal{R}_A(\mathcal{R}_B(\mathbf{z}^k)). \end{cases}$$

can be expressed as

$$\mathbf{u}^{k+1} := \mathcal{J}_B \left(\mathcal{J}_A(2\mathcal{J}_B(\mathcal{J}_B^{-1}(\mathbf{u}^k)) - \mathcal{J}_B^{-1}(\mathbf{u}^k)) + \mathcal{J}_B^{-1}(\mathbf{u}^k) - \mathcal{J}_B(\mathcal{J}_B^{-1}(\mathbf{u}^k)) \right). \quad (35)$$

- ▶ First, it is obvious that $\mathcal{J}_B(\mathcal{J}_B^{-1}(\mathbf{u}^k)) = \mathbf{u}^k$
- ▶ Second, by definition of \mathcal{J}_B , we also have $\mathcal{J}_B^{-1}(\mathbf{u}^k) = \mathbf{u}^k + B(\mathbf{u}^k)$.

Substituting these relations into (35), we obtain

$$\mathbf{u}^{k+1} := \mathcal{J}_B \left(\mathcal{J}_A(\mathbf{u}^k - B(\mathbf{u}^k)) + B(\mathbf{u}^k) \right) \quad (36)$$

We can rewrite (36) as

$$(\mathbb{I} + B)(\mathbf{u}^{k+1}) \in (\mathbb{I} + B)(\mathbf{u}^k) - e(\mathbf{u}^k) \quad (37)$$

where $e(\mathbf{u}^k) := \mathbf{u}^k - \mathcal{J}_A(\mathbf{u}^k - B(\mathbf{u}^k))$.

Convergence rate

Facts:

- ▶ It is obvious that $\mathbf{u}^* = \mathcal{J}_A(\mathbf{u}^* - B(\mathbf{u}^*))$ for any $\mathbf{u}^* \in \mathcal{U}^*$.
- ▶ We define $e(\mathbf{u}) := \mathbf{u} - \mathcal{J}_A(\mathbf{u} - B(\mathbf{u}))$ the **residual operator** at \mathbf{u} . Then $e(\mathbf{u}^*) = 0$.
- ▶ Let $\mathcal{T} := \mathbb{I} + A$. The **Douglas-Rachford** scheme (37) can be written as

$$\boxed{\mathcal{T}(\mathbf{u}^{k+1}) \in \mathcal{T}(\mathbf{u}^k) - e(\mathbf{u}^k)} \quad (38)$$

Theorem (Convergence rate [7])

Assume that B is **single-valued** and the **solution set** \mathcal{U}^* of $0 \in A(\mathbf{u}) + B(\mathbf{u})$ is **nonempty**. Let $\{\mathbf{u}^k\}_{k \geq 0}$ be the sequence generated by the **Douglas-Rachford algorithm** with $\eta_k = 1$ for all $k \geq 0$. Then for any **solution** \mathbf{u}^* of $0 \in A(\mathbf{u}) + B(\mathbf{u})$, we have

$$\|e(\mathbf{u}^k)\|^2 \leq \frac{1}{k+1} \|(\mathbf{u}^0 - \mathbf{u}^*) + (B(\mathbf{u}^0) - B(\mathbf{u}^*))\|^2. \quad (39)$$

The sequence $\{\|e(\mathbf{u}^k)\|^2\}$ **converges** to zero at the **$\mathcal{O}(1/k)$ rate**. Consequently, $\{\mathbf{u}^k\}$ **converges** to a solution \mathbf{u}^* of $0 \in A(\mathbf{u}) + B(\mathbf{u})$.

Remarks

- ▶ Since we assume that B is **single-valued**, the right-hand side (39) is **bounded**.
- ▶ In the case B is a **set-valued mapping**, the **right-hand side** may be **no longer bounded**. For example, if $B = \mathcal{N}_{\mathcal{Z}}$ the normal cone of a convex set \mathcal{Z} .

Example: Image denoising

TV-denoising

Given a **noisy image** $\mathbf{b} \in \mathbb{R}^{m \times n}$, we want to recover a **clean image** \mathbf{x} by using the **total variation operator**. This problem can be formulated as:

$$\min_{\mathbf{x} \in \mathbb{R}^{m \times n}} \frac{1}{2} \|\mathbf{x} - \mathbf{b}\|^2 + \rho \|\mathbf{x}\|_{\text{TV}},$$

where $\rho > 0$ is a **regularization parameter** and

$$\|\mathbf{x}\|_{\text{TV}} := \begin{cases} \sum_{i,j} |x_{i,j+1} - x_{i,j}| + |x_{i,j+1} - x_{i,j}| & \text{anisotropic (ATV)} \\ \sum_{i,j} \sqrt{(x_{i,j+1} - x_{i,j})^2 + (x_{i,j+1} - x_{i,j})^2} & \text{isotropic (ITV)} \end{cases}$$

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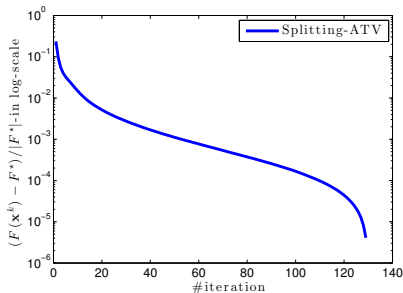
$$\|\mathbf{x}\|_{\text{TV}} := \begin{cases} \sum_{i,j} |x_{i,j+1} - x_{i,j}| + |x_{i,j+1} - x_{i,j}| & \text{anisotropic (ATV)} \\ \sum_{i,j} \sqrt{(x_{i,j+1} - x_{i,j})^2 + (x_{i,j+1} - x_{i,j})^2} & \text{isotropic (ITV)} \end{cases}$$

How to apply the **splitting algorithm**?

- ▶ By letting $\mathbf{z} = \mathbf{D}\mathbf{x}$, we can convert the **TV-denoising problem** into (2), where \mathbf{D} is a matrix representing the **total variation**.
- ▶ **Splitting algorithm** is now applied to the **resulting problem**.
- ▶ We choose $\tau_k = \eta_k = 1$ in our test.

Example: Image denoising - convergence behavior

- ▶ The convergence of the splitting method using ATV and ITV norms on a gray image of size 512×512 .
- ▶ The regularization parameter $\rho = 20$.

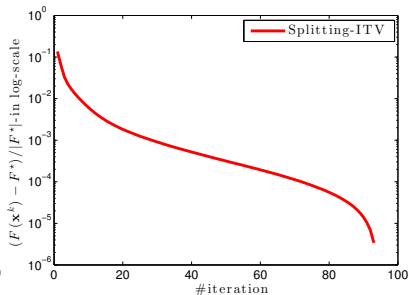


- ▶ The objective value: 9.9319×10^7
- ▶ Relative error:

$$\|\mathbf{x}^k - \mathbf{x}^{\natural}\| / \|\mathbf{x}^{\natural}\| = 0.0727258$$

where \mathbf{x}^{\natural} is the original image

- ▶ The number of iterations: 131
- ▶ The CPU time: 30.6s.
- ▶ The PSNR = 23.1983.



- ▶ The objective value: 9.4897×10^7
- ▶ Relative error:

$$\|\mathbf{x}^k - \mathbf{x}^{\natural}\| / \|\mathbf{x}^{\natural}\| = 0.0676706$$

where \mathbf{x}^{\natural} is the original image

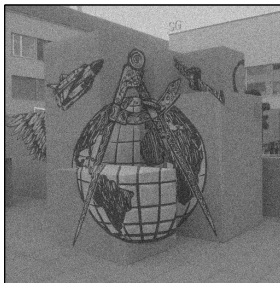
- ▶ The number of iterations: 95
- ▶ The CPU time: 26.0s.
- ▶ The PSNR = 23.6192.

Example: Image denoising

Original



Noisy

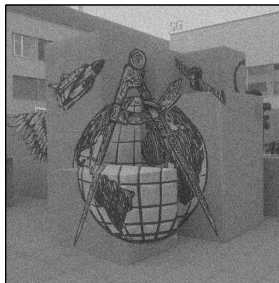


Example: Image denoising

Original



Noisy



Anisotropic TV denoising



Isotropic TV denoising



Primal-dual hybrid gradient (PDHG) algorithm

The idea of PDHG

- ▶ Originally, PDHG is a combination of a primal and dual proximal-gradient descent step applying to the min-max problem (17) [17]:

$$\min_{\mathbf{u} \in \mathbb{R}^p} \max_{\mathbf{v} \in \mathbb{R}^q} \left\{ \langle \mathbf{K}\mathbf{u}, \mathbf{v} \rangle + f(\mathbf{u}) - g^*(\mathbf{v}) \right\}.$$

- ▶ First, PDHG performs a primal proximal-gradient step on the minimization problem w.r.t. \mathbf{u} given \mathbf{v}^k to compute \mathbf{u}^{k+1} .
- ▶ Second, PDHG performs a dual proximal-gradient step on the maximization problem w.r.t. \mathbf{v} given \mathbf{u}^{k+1} to compute \mathbf{v}^{k+1} .
- ▶ We can also add an intermediate step $\bar{\mathbf{u}}^{k+1}$ before performing the dual step.

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Primal-Dual Hybrid Gradient algorithm (PDHG) for (15)

- Given $\mathbf{u}^0 \in \mathbb{R}^p$ and $\mathbf{v}^0 \in \mathbb{R}^n$ as an initial point.
- Choose $\tau_0 > 0$, $\sigma_0 > 0$ and $\eta_0 \neq 0$.
- For $k = 0, 1, \dots$, perform:

$$\begin{cases} \mathbf{u}^{k+1} &:= \text{prox}_{\tau_k f}(\mathbf{u}^k - \tau_k \mathbf{K}^T \mathbf{v}^k), \\ \bar{\mathbf{u}}^{k+1} &:= \mathbf{u}^{k+1} + \eta_k (\mathbf{u}^{k+1} - \mathbf{u}^k), \\ \mathbf{v}^{k+1} &:= \text{prox}_{\sigma_k g^*}(\mathbf{v}^k + \sigma_k \mathbf{K} \bar{\mathbf{u}}^{k+1}). \end{cases}$$

and update τ_k , σ_k and η_k if necessary.

Connection to Chambolle-Pock's algorithm and enhancement

Connection to Chambolle-Pock's algorithm

- ▶ PDHG is **very similar** to **Chambolle-Pock's algorithm**.
 - ▶ **Chambolle-Pock's algorithm** performs the **dual step** on \mathbf{v} and then the **primal step** on \mathbf{u} .
 - ▶ **PDHG** performs the **primal step** on \mathbf{u} and then the **dual step** on \mathbf{v} .
 - ▶ Symmetrically, we can say that both methods are **equivalent**.
- ▶ The **convergence theory** of **Chambolle-Pock's algorithm** is **applicable** to **PDHG**.

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Adaptive PDHG

- ▶ If we select the **parameters** τ_k and σ_k based on the **convergence theory**, **PDHG** often has **poor performance** in practice.
- ▶ We can **enhance** the performance of **PDHG** by **adaptively updating** τ_k and σ_k [5].
 - ▶ Define the **primal** and **dual residuals** ($\mathbf{p}^{k+1}, \mathbf{d}^{k+1}$):
$$\begin{cases} \mathbf{p}^{k+1} &:= \tau_k^{-1}(\mathbf{u}^k - \mathbf{u}^{k+1}) - \mathbf{K}^T(\mathbf{v}^k - \mathbf{v}^{k+1}), \\ \mathbf{d}^{k+1} &:= \sigma_k^{-1}(\mathbf{v}^k - \mathbf{v}^{k+1}) + \mathbf{K}(\mathbf{u}^k - \mathbf{u}^{k+1}). \end{cases}$$
 - ▶ **Adaptively update** τ_k and σ_k by **trading-off** the **primal residual** $\|\mathbf{p}^{k+1}\|$ and **dual residual** $\|\mathbf{d}^{k+1}\|$ at each iteration.

Outline

► Today

1. Convex constrained optimization and motivating examples
2. Optimality condition
3. Conjugate functions
4. Monotone inclusion and monotone mixed variational inequality formulations
5. Chambolle-Pock's primal-dual method
6. Primal-dual hybrid gradient method
7. Splitting methods
8. Model-based excessive gap primal-dual method

► Next week

1. Disciplined convex programming

Primal-dual method using model-based excessive gap technique

Problem restatement

We consider again problem (2) with **additional convex constraint** as:

$$f^* := \begin{cases} \min_{\mathbf{x} \in \mathbb{R}^p} & f(\mathbf{x}) \\ \text{s.t.} & \mathbf{Ax} = \mathbf{b}, \quad \mathbf{x} \in \mathcal{X} \end{cases} \quad (40)$$

where f , \mathbf{A} and \mathbf{b} are defined as in (2) and \mathcal{X} is a **nonempty, closed, convex and bounded** set in \mathbb{R}^p .

We recall the **variational inequality** presenting the **optimality condition** of (40) as

$$f(\mathbf{x}) - f(\mathbf{x}^*) + M(\mathbf{z}^*)^T (\mathbf{z} - \mathbf{z}^*) \geq 0, \quad \forall \mathbf{z} \in \mathcal{X} \times \mathbb{R}^n \quad (41)$$

where $\mathbf{z}^* := (\mathbf{x}^*, \lambda^*) \in \mathcal{X} \times \mathbb{R}^n$ is a primal-dual solution of (40) and

$$M(\mathbf{z}) := \begin{bmatrix} \mathbf{A}^T \lambda \\ \mathbf{b} - \mathbf{Ax} \end{bmatrix}.$$

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Gap function

The **gap function** of (41) is redefined as (different from (14) at $\hat{\mathbf{z}} \in \mathcal{X} \times \mathbb{R}^n$):

$$G(\mathbf{z}) := \max_{\hat{\mathbf{z}} \in \mathcal{X} \times \mathbb{R}^n} \left\{ f(\mathbf{x}) - f(\hat{\mathbf{x}}) + M(\mathbf{z})^T(\mathbf{z} - \hat{\mathbf{z}}) \right\} \quad (42)$$

Smoothed gap function

- ▶ Let b be a **prox-function** of \mathcal{X} with **strong convexity parameter** $\sigma_b = 1$.
 - ▶ i.e., b is a **smooth, strongly convex** function with strong convexity parameter $\sigma_b = 1$.
- ▶ We define ξ the **Bregman distance** corresponding to b as

$$\xi(\mathbf{x}, \hat{\mathbf{x}}) := b(\mathbf{x}) - b(\hat{\mathbf{x}}) - \nabla b(\hat{\mathbf{x}})^T(\mathbf{x} - \hat{\mathbf{x}}).$$

- ▶ For $\mathbf{x}_c \in \mathbb{R}^p$ and two **positive parameters** γ and β we define

$$\xi_{\gamma\beta}(\mathbf{z}) := \gamma\xi(\mathbf{z}, \mathbf{x}_c) + (\beta/2)\|\mathbf{z}\|_2^2 \quad (43)$$

a **smoother** for the **gap function** (42).

- ▶ $\xi_{\gamma\beta}$ is **strongly convex** and satisfies $\xi_{\gamma\beta}(\mathbf{z}) \geq (\gamma/2)\|\mathbf{z} - \mathbf{x}_c\|_2^2 + (\beta/2)\|\mathbf{z}\|_2^2$.

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We define a **smoothed version** for the **gap function** G given by (42) as follows:

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Smoothed gap function

We define a smoothed version for the gap function G given by (42) as follows:

$$G_{\gamma\beta}(\mathbf{z}) := \max_{\hat{\mathbf{z}} \in \mathcal{X} \times \mathbb{R}^n} \left\{ f(\mathbf{z}) - f(\hat{\mathbf{z}}) + M(\mathbf{z})^T(\mathbf{z} - \hat{\mathbf{z}}) - \xi_{\gamma\beta}(\hat{\mathbf{z}}) \right\} \quad (44)$$

Properties of $G_{\gamma\beta}$

- ▶ Evaluating $G_{\gamma\beta}$ and its gradient requires to solve a strongly convex program (42)
- ▶ $G_{\gamma\beta}(\mathbf{z}) \rightarrow G(\mathbf{z})$ as γ and β go to zero for all $\mathbf{z} \in \mathbb{R}^{p+n}$.

Comments on the gap function G and its smooth version $G_{\gamma\beta}$

- ▶ The gap function G defined by (42):

$$G(\mathbf{z}) := \max_{\hat{\mathbf{z}} \in \mathcal{X} \times \mathbb{R}^n} \left\{ f(\mathbf{x}) - f(\hat{\mathbf{x}}) + M(\mathbf{z})^T (\mathbf{z} - \hat{\mathbf{z}}) \right\}$$

can be written as

$$G(\mathbf{z}) := \max_{\hat{\mathbf{z}} \in \mathbb{R}^p \times \mathbb{R}^n} \left\{ (f(\mathbf{x}) + \iota_{\mathcal{X}}(\mathbf{x})) - (f(\hat{\mathbf{x}}) + \iota_{\mathcal{X}}(\hat{\mathbf{x}})) + M(\mathbf{z})^T (\mathbf{z} - \hat{\mathbf{z}}) \right\}$$

which is exactly the **gap function** G in the unconstrained form (14), where $\iota_{\mathcal{X}}$ is the **indicator function** of \mathcal{X} .

- ▶ If we define

$$\begin{aligned} d_{\gamma}(\lambda) &:= \min_{\mathbf{x} \in \mathcal{X}} \left\{ f(\mathbf{x}) + \lambda^T (\mathbf{A}\mathbf{x} - \mathbf{b}) + \gamma \xi(\mathbf{x}; \mathbf{x}_c) \right\} \\ f_{\beta}(\mathbf{x}) &:= f(\mathbf{x}) + \max_{\lambda \in \mathbb{R}^n} \left\{ \lambda^T (\mathbf{A}\mathbf{x} - \mathbf{b}) - (\beta/2) \|\lambda\|^2 \right\} \\ &= f(\mathbf{x}) + (1/(2\beta)) \|\mathbf{A}\mathbf{x} - \mathbf{b}\|^2 \end{aligned}$$

then:

- ▶ d_{γ} is an **approximation** of the **dual function** d
- ▶ f_{β} is an **approximation** of the **objective function** f .

Moreover, we have

$$G_{\gamma\beta}(\mathbf{z}) = f_{\beta}(\mathbf{x}) - d_{\gamma}(\lambda).$$

Illustration of smooth gap function $G_{\gamma\beta}$

- ▶ The **objective function** $f(\mathbf{x})$ **does not depend on** β
- ▶ The **dual function** $d(\lambda)$ **does not depend on** γ :

$$d(\bar{\lambda}) := \min_{\mathbf{x} \in \mathcal{X}} \{f(\mathbf{x}) + \bar{\lambda}^T (\mathbf{A}\mathbf{x} - \mathbf{b})\}.$$

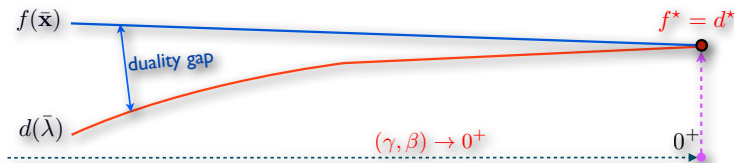
- ▶ $f(x)$ is **decreasing** over \mathbf{x} and $d(\lambda)$ is **increasing** over λ .

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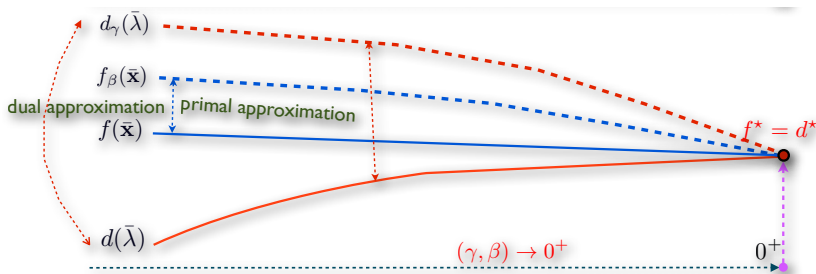
- ▶ The **duality gap** is defined as $G(\bar{\mathbf{z}}) := f(\bar{\mathbf{x}}) - d(\bar{\lambda}) \geq 0$.
- ▶ At the **optimal solution** $\mathbf{z}^* := (\mathbf{x}^*, \lambda^*)$, one has $f(\mathbf{x}^*) = d(\lambda^*)$ and $G(\mathbf{z}^*) = 0$.

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- ▶ $f(x)$ is **decreasing** over \mathbf{x} and $d(\lambda)$ is **increasing** over λ .



- ▶ The **augmented function** f_{β} **approximates** f : $f_{\beta}(\mathbf{x}) = f(\mathbf{x}) + (1/(2\beta))\|\mathbf{A}\mathbf{x} - \mathbf{b}\|^2$
- ▶ The **smoothed dual function** d_{γ} **approximates** d :

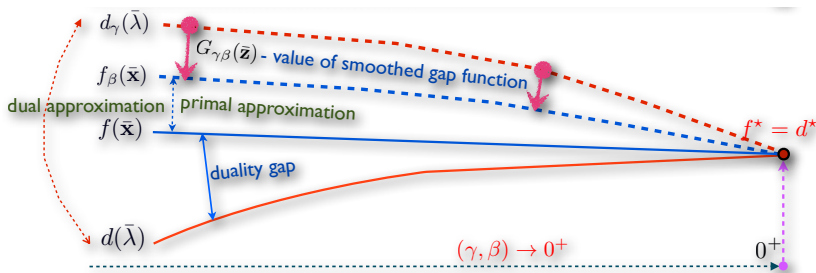
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$$G_{\gamma_k\beta_k}(\bar{\mathbf{z}}^k) \rightarrow G(\mathbf{z}^*) = 0 \text{ as } \gamma_k\beta_k \rightarrow 0^+.$$

Model-based excessive gap technique

What is the smoothed gap function used for?

Aim: To generate a primal-dual sequence $\{\bar{\mathbf{z}}^k\}_{k \geq 0}$ with $\bar{\mathbf{z}}^k := (\bar{\mathbf{x}}^k, \bar{\lambda}^k)$ such that

$$G_{\gamma_k \beta_k}(\bar{\mathbf{z}}^k) \rightarrow 0^+$$

by controlling γ_k and $\beta_k \rightarrow 0^+$.

► When γ_k and β_k go to zero, we have $G_{\gamma_k \beta_k}(\cdot) \rightarrow G(\cdot)$.

► **Consequence:** $G(\mathbf{z}^k) \rightarrow 0^+ \Rightarrow \bar{\mathbf{z}}^k \rightarrow \mathbf{z}^* = (\mathbf{x}^*, \lambda^*)$ (primal-dual solution).

Model-based excessive gap condition

A sequence $\{\bar{\mathbf{z}}^k\}_{k \geq 0} \subset \mathcal{X} \times \mathbb{R}^n$ is said to satisfy the model-based excessive gap condition if

$$G_{\gamma_{k+1} \beta_{k+1}}(\bar{\mathbf{z}}^{k+1}) \leq (1 - \tau_k) G_{\gamma_k \beta_k}(\bar{\mathbf{z}}^k) - \psi_k \quad (45)$$

where $\psi_k \geq 0$, $\tau_k \in (0, 1)$ and $\gamma_k \beta_{k+1} < \gamma_k \beta_k$ for $k \geq 0$.

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where $\psi_k \geq 0$, $\tau_k \in (0, 1)$ and $\gamma_k \beta_{k+1} < \gamma_k \beta_k$ for $k \geq 0$.

Let $\bar{G}_k := G_{\gamma_k \beta_k}(\bar{\mathbf{z}}^k)$. By induction, we have

$$\bar{G}_{k+1} \leq \prod_{j=0}^k (1 - \tau_j) \bar{G}_0 - \left[\psi_0 + \sum_{j=1}^{k-1} \prod_{l=0}^{j-1} (1 - \tau_l) \psi_j \right].$$

\Rightarrow The convergence rate of $\{\bar{G}_k\}$ depends on the convergence rate of $\{\tau_k\}$.

Key estimates

- ▶ For a **bounded set** \mathcal{X} and $\hat{\mathbf{x}} \in \mathcal{X}$, the quality $D_{\mathcal{X}}$ defined below is **finite**

$$D_{\mathcal{X}} := \max_{\mathbf{x} \in \mathcal{X}} \xi(\mathbf{x}, \hat{\mathbf{x}}) < +\infty.$$

- ▶ Denote

$$\omega_k := \prod_{j=0}^k (1 - \tau_j) \quad \text{and} \quad \Psi_k := \psi_0 + \sum_{j=1}^{k-1} \prod_{l=0}^{j-1} (1 - \tau_l) \psi_j.$$

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Theorem (Bounds on the objective residual and primal feasibility)

Assume that $\{\bar{\mathbf{z}}^k\}_{k \geq 0}$ is a sequence satisfying (45). Then

$$\begin{cases} -\|\lambda^*\| \|\mathbf{A}\bar{\mathbf{x}}^k - \mathbf{b}\| \leq f(\bar{\mathbf{x}}^k) - f^* \leq C_k, \\ \|\mathbf{A}\bar{\mathbf{x}}^k - \mathbf{b}\| \leq \beta_k \left[\|\lambda^*\| + \sqrt{\|\lambda^*\|^2 + 2\beta_k^{-1} C_k} \right], \end{cases} \quad (46)$$

where $C_k := \omega_{k-1} G_{\gamma_0 \beta_0}(\bar{\mathbf{z}}^0) + \gamma_k D_{\mathcal{X}} - \Psi_{k-1}$, provided that $\beta_k \|\lambda^*\| + 2C_k \geq 0$.

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As a consequence, we have

$$\begin{cases} |f(\bar{\mathbf{x}}^k) - f^*| \leq \max \left\{ \gamma_k D_{\mathcal{X}}, \left(2\beta_k D_{\Lambda^*} + \sqrt{2\gamma_k \beta_k D_{\mathcal{X}}} \right) D_{\Lambda^*} \right\}, \\ \|\mathbf{A}\bar{\mathbf{x}}^k - \mathbf{b}\| \leq 2\beta_k D_{\Lambda^*} + \sqrt{2\gamma_k \beta_k D_{\mathcal{X}}}, \end{cases} \quad (47)$$

where $D_{\Lambda^*} := \min\{\|\lambda^*\| : \lambda^* \in \Lambda^*\}$ the **norm** of the **minimum norm solution** of the **dual problem**.

Sketch of proof

- From the **saddle point inequalities**, we have $f^* = \mathcal{L}(\mathbf{x}^*, \lambda^*) \leq \mathcal{L}(\mathbf{x}, \lambda^*)$. Hence,

$$d(\lambda) \leq f^* \leq f(\mathbf{x}) + (\mathbf{Ax} - \mathbf{b})^T \lambda^* \leq f(\mathbf{x}) + \|\mathbf{Ax} - \mathbf{b}\| \|\lambda^*\|, \quad \forall \mathbf{x} \in \mathcal{X}.$$

We finally get $-\|\lambda^*\| \|\mathbf{Ax} - \mathbf{b}\| \leq f(\mathbf{x}) - f^* \leq f(\mathbf{x}) - d(\lambda)$ for all $\mathbf{x} \in \mathcal{X}$.

- Since $\xi(\mathbf{x}, \mathbf{x}_c) \geq 0$ and $d(\lambda) = \min_{\mathbf{x} \in \mathcal{X}} \{f(\mathbf{x}) + \lambda^T (\mathbf{Ax} - \mathbf{b})\}$, we have

$$d(\lambda) \leq d_\gamma(\lambda) \leq d(\lambda) + \gamma D_{\mathcal{X}}$$

- Putting things together, we get

$$\begin{aligned} -\|\lambda^*\| \|\mathbf{Ax} - \mathbf{b}\| &\leq f(\mathbf{x}) - f^* \leq f(\mathbf{x}) - d(\lambda) \\ &\leq f_\beta(\mathbf{x}) - d_\gamma(\lambda) + \gamma D_{\mathcal{X}} - (1/(2\beta)) \|\mathbf{Ax} - \mathbf{b}\|^2 \\ &= G_{\gamma\beta}(\mathbf{z}) + \gamma D_{\mathcal{X}} - (1/(2\beta)) \|\mathbf{Ax} - \mathbf{b}\|^2. \end{aligned} \quad (48)$$

- Since $G_{\gamma_k\beta_k}(\bar{\mathbf{z}}^k) \leq \omega_{k-1} G_{\gamma_0\beta_0}(\bar{\mathbf{z}}^0) - \Psi_{k-1}$ due to (45), we obtain from (48)

$$-\|\lambda^*\| \|\mathbf{Ax}^k - \mathbf{b}\| \leq f(\bar{\mathbf{x}}^k) - f^* \leq \omega_{k-1} G_{\gamma_0\beta_0}(\bar{\mathbf{z}}^0) - \Psi_{k-1} + \gamma_k D_{\mathcal{X}} = C_k. \quad (49)$$

which is the **first inequality of (46)**.

- Let $s := \|\mathbf{Ax}^k - \mathbf{b}\|$. From (49) and (48) we have $s^2 - 2\beta_k \|\lambda^*\| s - 2\beta_k C_k \leq 0$. Solving this in equation, we obtain the **second inequality of (46)**.

Evaluating the smoothed gap function $G_{\gamma\beta}$

Evaluation of $G_{\gamma\beta}$

In order to evaluate $G_{\gamma\beta}$, we need to solve the maximization problem:

$$G_{\gamma\beta}(\mathbf{z}) := \max_{\hat{\mathbf{z}} \in \mathcal{X} \times \mathbb{R}^n} \left\{ f(\mathbf{x}) - f(\hat{\mathbf{x}}) + M(\mathbf{z})^T(\mathbf{z} - \hat{\mathbf{z}}) - d_{\gamma\beta}(\hat{\mathbf{z}}) \right\}$$

The solution $\mathbf{z}_{\gamma\beta}^*(\mathbf{z}) := (\mathbf{x}_{\gamma}^*(\lambda), \lambda_{\beta}^*(\mathbf{x}))$ of this problem is given as

$$\begin{cases} \mathbf{x}_{\gamma}^*(\lambda) &:= \underset{\mathbf{x} \in \mathcal{X}}{\operatorname{argmin}} \left\{ f(\mathbf{x}) + (\mathbf{A}^T \lambda)^T \mathbf{x} + \gamma d(\mathbf{x}, \mathbf{x}_c) \right\} \\ \lambda_{\beta}^*(\mathbf{x}) &:= \beta^{-1}(\mathbf{A}\mathbf{x} - \mathbf{b}). \end{cases} \quad (50)$$

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Given $\bar{\mathbf{z}}^k := (\bar{\mathbf{x}}^k, \bar{\lambda}^k)$ and (γ_k, β_k) . The **idea of the algorithms** is to:

- **Update** $\bar{\mathbf{z}}^{k+1} := (\bar{\mathbf{x}}^{k+1}, \bar{\lambda}^{k+1})$ from $\bar{\mathbf{z}}^k$ and $\mathbf{z}_{\gamma_k\beta_k}^*(\mathbf{z})$.
- **Decrease** the parameters $(\gamma_{k+1}, \beta_{k+1})$ such that $\gamma_{k+1}\beta_{k+1} < \gamma_k\beta_k$.

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Proximal-gradient step

In our algorithms, we need to compute $\bar{\mathbf{x}}^{k+1}$ using the following mapping:

$$\text{prox}_{\beta f}(\mathbf{x}, \lambda) := \arg\min_{\hat{\mathbf{x}} \in \mathcal{X}} \left\{ f(\hat{\mathbf{x}}) + (\mathbf{A}^T \lambda)^T \hat{\mathbf{x}} + (\|\mathbf{A}\|^2 / (2\beta)) \|\hat{\mathbf{x}} - \mathbf{x}\|_2^2 \right\} \quad (51)$$

How can we generate $\{\bar{\mathbf{z}}^k\}$?

The **main idea** of generating the **sequence** $\{\bar{\mathbf{z}}^k\}$ such that $\{G_{\gamma_k\beta_k}(\bar{\mathbf{z}}^k)\}$ **decreases** come from the following **observations**:

- ▶ Since $G_{\gamma\beta}(\mathbf{z}) = f_\beta(\mathbf{x}) - d_\gamma(\lambda)$, then $G_{\gamma\beta}(\mathbf{z})$ **decreases** if **at least** $f_\beta(\mathbf{x})$ **decrease** or $d_\gamma(\lambda)$ **increases**.

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- ▶ We note on the one hand that:
 - ▶ d_γ is a **smoothed version** of the **dual function** d .
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In order to **increase** d_γ as much as possible, one can implement a **Nesterov's accelerated gradient ascent scheme**.

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 - ▶ f_β is the sum of a **convex** function f and $g(\mathbf{x}) := (1/(2\beta))\|\mathbf{Ax} - \mathbf{b}\|^2$.
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- ▶ Every iteration, we can perform **one** scheme or **both** in order to **decrease** $G_{\gamma\beta}(\mathbf{z})$, while **simultaneously decrease** the **product** $\gamma\beta$.

This observations lead to the following **two update schemes** in the next slides.

How can we generate $\{\bar{\mathbf{z}}^k\}$?

Expression of the idea

We denote by \mathbf{s} a **variable** standing for either \mathbf{x} or λ . The **main step** of the **algorithmic scheme** consists of

- ▶ For given \mathbf{s} , \mathbf{s}^* and $\tau \in (0, 1]$ at the **current iteration**. We perform **one interpolation step**:

$$\hat{\mathbf{s}} := (1 - \tau)\mathbf{s} + \tau\mathbf{s}^*.$$

- ▶ When $\hat{\mathbf{s}}$ is **available**, we perform a **proximal-gradient** or **gradient step** to compute the **next iteration** \mathbf{s}^+ :

$$\mathbf{s}^+ := \text{prox}_{\psi}(\hat{\mathbf{s}} - (1/L)\nabla\varphi(\hat{\mathbf{s}})).$$

where

- ▶ $\varphi := d_{\gamma}$ the **smoothed dual function** and $\psi = f$ in the **primal step**
- ▶ $\varphi := (1/(2\beta))\|\mathbf{A}\mathbf{x} - \mathbf{b}\|^2$ and $\psi = 0$ and in the **dual scheme**.
- ▶ L is the Lipschitz constant of $\nabla\varphi$.

Remarks

- ▶ The above scheme looks **very similar** to FISTA [1] in the context of **Nesterov's accelerating method**.

How can we generate $\{\bar{\mathbf{z}}^k\}$?

We propose **two schemes** to generate the sequence $\{\bar{\mathbf{z}}^k\}$

- ▶ The **primal-dual** scheme with **two primal steps and one dual step** (2P1D);
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$$\begin{cases} \hat{\mathbf{x}}^k & := (1 - \tau_k)\bar{\mathbf{x}}^k + \tau_k \mathbf{x}_{\gamma_k}^* (\bar{\lambda}^k) \\ \bar{\mathbf{x}}^{k+1} & := \text{prox}_{\beta_{k+1}f}(\hat{\mathbf{x}}^k, \lambda_{\beta_{k+1}}^* (\hat{\mathbf{x}}^k)) \\ \bar{\lambda}^{k+1} & := (1 - \tau_k)\bar{\lambda}^k + \tau_k \lambda_{\beta_{k+1}}^* (\hat{\mathbf{x}}^k) \end{cases} \quad (2P1D)$$

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where $\alpha_k := \gamma_{k+1} \|\mathbf{A}\|^{-2}$.

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where $\alpha_k := \gamma_{k+1} \|\mathbf{A}\|^{-2}$.

- ▶ The parameters β_k and γ_k are updated as ($c_k \in (-1, 1]$ given):

$$\gamma_{k+1} := (1 - c_k \tau_k) \gamma_k \quad \text{and} \quad \beta_{k+1} = (1 - \tau_k) \beta_k \quad (52)$$

Remarks on the computational complexity of both schemes

- ▶ (2P1D) requires **two primal steps**: **one** to compute $\mathbf{x}_{\gamma_k}^*(\bar{\lambda}^k)$ and **one** to compute $\bar{\mathbf{x}}^{k+1}$.

- ▶ The **first step** corresponds to solving

$$\mathbf{x}_{\gamma_k}^*(\bar{\lambda}^k) := \arg \min_{\mathbf{x} \in \mathcal{X}} \left\{ f(\mathbf{x}) + (\mathbf{A}^T \bar{\lambda}^k)^T \mathbf{x} + \gamma_k d(\mathbf{x}, \mathbf{x}_c) \right\}.$$

- ▶ The **second step** corresponds to solving

$$\bar{\mathbf{x}}^{k+1} := \arg \min_{\hat{\mathbf{x}} \in \mathcal{X}} \left\{ f(\hat{\mathbf{x}}) + (\mathbf{A}^T \lambda_{\beta_{k+1}}^*(\hat{\mathbf{x}}))^T \hat{\mathbf{x}} + (\|\mathbf{A}\|^2 / (2\beta)) \|\hat{\mathbf{x}} - \hat{\mathbf{x}}^k\|_2^2 \right\}$$

- ▶ If b is a **quadratic prox-function** and \mathcal{X} is **absent**, then solving **both problems** corresponds to computing the **proximal operator** of f .
- ▶ (1P2D) only requires **one primal step** to compute $\mathbf{x}_{\gamma_{k+1}}^*(\hat{\lambda}^k)$.
- ▶ (1P2D) requires **two dual steps** corresponding to two **matrix-vector multiplications** $\mathbf{A}\bar{\mathbf{x}}^k$ and $\mathbf{A}\mathbf{x}_{\gamma_{k+1}}^*(\hat{\lambda}^k)$.
- ▶ (2P1D) requires only **one** $\mathbf{A}\hat{\mathbf{x}}^k$.

Updating step-size

The **key point** in both schemes (1P2D) and (2P1D) is to update the **step size** τ_k :

- ▶ The **model-based excessive gap** condition (45) shows that $G_{\gamma_k \beta_k}(\bar{\mathbf{z}}^k) \rightarrow 0^+$.
- ▶ The convergence rate of $\{\beta_k\}$ and $\{\gamma_k\}$ depends on the convergence rate of $\{\tau_k\}$ and hence, the convergence rate of the algorithms.

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Theorem (Key condition)

Let $\{\bar{\mathbf{z}}^k\}$ be the sequence generated by **either** (1P2D) or (2P1D) and $G_{\gamma_k \beta_k}(\bar{\mathbf{z}}^k) \leq 0$. Then, under the condition:

$$\tau^2 \|\mathbf{A}\|^2 \leq \gamma_{k+1} \beta_{k+1} \quad (53)$$

we have $G_{\gamma_{k+1} \beta_{k+1}}(\bar{\mathbf{z}}^{k+1}) \leq 0$.

Condition (53) and the update rules $\gamma_{k+1} := (1 - c_k \tau_k) \gamma_k$ and $\beta_{k+1} = (1 - \tau_k) \beta_k$ allow us to derive the update rule for τ_k :

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Update the step-size τ_k

- ▶ **Initialization:** $\tau_0 := a_0^{-1}$, where $a_0 := (1 + c_0 + [4(1 - c_0) + (1 + c_0)^2]^{1/2})/2$.
- ▶ **Update:** τ_{k+1} is updated from τ_k as

$$\tau_k = a_k^{-1}, \quad a_{k+1} := (1 + c_{k+1} + \sqrt{4a_k^2 + (1 - c_{k+1})^2})/2$$

Primal-dual framework using model-based excessive gap technique

Putting all ingredients together, we can describe the complete algorithm as below:

Primal-dual method using model-based excessive gap technique (PDM)
Initialization
<p>1.1. Given $\gamma_0 > 0$, $c_0 \in (-1, 1]$ and $\bar{L}_d := \ \mathbf{A}\ ^2$.</p> <p>1.2. $a_0 := (1 + c_0 + \sqrt{4(1 - c_0) + (1 + c_0)^2})/2$, $\tau_0 := a_0^{-1}$ and $\beta_0 := \bar{L}_g \gamma^{-1}$.</p> <p>1.3. Compute a starting point $\bar{\mathbf{z}}^0 := (\bar{\mathbf{x}}^0, \bar{\lambda}^0)$.</p>

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Iterations: For $k = 0, 1, \dots, K$, perform:
<p>2.1. Given $(\bar{\mathbf{x}}^k, \bar{\lambda}^k)$, compute $(\bar{\mathbf{x}}^{k+1}, \bar{\lambda}^{k+1})$ by either (2P1D) or (1P2D).</p> <p>2.2. Update $\gamma_{k+1} := (1 - c_k \tau_k) \gamma_k$ and $\beta_{k+1} := (1 - \tau_k) \beta_k$.</p> <p>2.3. Update c_{k+1} from c_k if necessary (optional).</p> <p>2.4. Update $a_{k+1} := (1 + c_{k+1} + \sqrt{4a_k^2 + (1 - c_{k+1})^2})/2$ and $\tau_{k+1} = a_{k+1}^{-1}$.</p>

Convergence guarantee

Theorem (Convergence)

Let $\{\bar{\mathbf{z}}^k := (\bar{\mathbf{x}}^k, \bar{\lambda}^k)\}$ be the sequence generated by **PDM** after $k \geq 1$ iterations.

a) If (2P1D) is used and $\gamma_0 := \|\mathbf{A}\|$ and $c_k = 1$ for all $k \geq 0$, then

$$\left\{ \begin{array}{lll} \|\mathbf{A}\bar{\mathbf{x}}^k - \mathbf{b}\| & \leq & \frac{\|\mathbf{A}\|(2D_{\Lambda^*} + \sqrt{D_{\mathcal{X}}})}{k+1}, \\ -D_{\Lambda^*} \|\mathbf{A}\bar{\mathbf{x}}^k - \mathbf{b}\| & \leq f(\bar{\mathbf{x}}^k) - f^* & \leq \frac{2\|\mathbf{A}\|D_{\mathcal{X}}}{k+1}. \end{array} \right.$$

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The *worst-case complexity* of **PDM** to reach an ϵ -solution \mathbf{x}^* of (40) is $\mathcal{O}\left(\frac{\|\mathbf{A}\|R}{\epsilon}\right)$, where $R := \max\{D_{\mathcal{X}}, D_{\Lambda^*} + \sqrt{D_{\mathcal{X}}}\}$.

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Remarks:

- If (1P2D) is used, then $\gamma_0 := \frac{2\sqrt{2}\|\mathbf{A}\|}{K+1}$, which requires to fix the **number of iterations priori**

Strongly convex case

PDM can be **accelerated** from $\mathcal{O}(1/k)$ to $\mathcal{O}(1/k^2)$ if f is **strongly convex**.

Assumption A.2.

The objective function f is **strongly convex** with the **convexity parameter** $\mu_f > 0$.

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We define the **dual function** of (40) as

$$d(\lambda) := \min_{\mathbf{x} \in \mathcal{X}} \left\{ f(\mathbf{x}) + \lambda^T (\mathbf{A}\mathbf{x} - \mathbf{b}) \right\}. \quad (54)$$

- ▶ Let $\mathbf{x}^*(\lambda)$ be the **solution** of (54)
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Properties of d

- ▶ d is **concave** and **smooth**.
- ▶ Gradient of d is given by $\nabla d(\lambda) := \mathbf{A}\mathbf{x}^*(\lambda) - \mathbf{b}$.
- ▶ ∇d is **Lipschitz continuous** with a **Lipschitz constant** $\bar{L}_d := \frac{\|\mathbf{A}\|^2}{\mu_f}$.

Algorithm

When specifying **PDM** to solve the **strongly convex case**, some steps in the algorithm are changed:

- ▶ Only **one smoothness parameter** β_k is **updated**.
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Primal-dual method for strongly convex case (PDM_μ)

Initialization:

1.1. Compute $\bar{L}_d := \mu_f^{-1} \|\mathbf{A}\|^2$, $\tau_0 := (\sqrt{5} - 1)/2$ and $\beta_0 := \sqrt{\bar{L}_d}$.

1.2. Compute a starting point $\bar{\mathbf{z}}^0 := (\bar{\mathbf{x}}^0, \bar{\lambda}^0)$ as:

$$\bar{\mathbf{x}}^0 := \mathbf{x}^*(0^n) \quad \text{and} \quad \bar{\lambda}^0 := \beta_0^{-1}(\mathbf{A}\bar{\mathbf{x}}^0 - \mathbf{b}).$$

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2.2. Update $\beta_{k+1} := (1 - \tau_k)\beta_k$ and $\tau_{k+1} := \tau_k(\sqrt{\tau_k^2 + 4} - \tau_k)/2$.

Convergence guarantee

Theorem (Convergence guarantee)

Assumptions:

- ▶ f is **strongly convex** with a strong convexity parameter $\mu_f > 0$.
- ▶ $\{\bar{\mathbf{z}}^k\}$ is generated by **PDM** $_{\mu}$.

Conclusions:

- ▶ We have estimates:

$$\left\{ \begin{array}{lll} -D_{\Lambda^*} \|\mathbf{A}\bar{\mathbf{x}}^k - \mathbf{b}\| & \leq f(\bar{\mathbf{x}}^k) - f^* & \leq 0 \\ & \|\mathbf{A}\bar{\mathbf{x}}^k - \mathbf{b}\| & \leq \frac{\|\mathbf{A}\|^2}{(k+2)^2 \mu_f} D_{\Lambda^*} \\ & \|\bar{\mathbf{x}}^k - \mathbf{x}^*\| & \leq \frac{\|\mathbf{A}\|}{(k+2)^2 \mu_f} D_{\Lambda^*} \end{array} \right.$$

- ▶ The bounds **do not** depend on $\mathcal{D}_{\mathcal{X}}$ the prox-diameter of \mathcal{X} .
- ▶ $\{\mathbf{x}^k\}$ converges to the **unique solution** \mathbf{x}^* of (40) at $\mathcal{O}(1/k^2)$ **rate**.

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Remarks:

- ▶ We always have $f(\bar{\mathbf{x}}^k) \leq f^*$ in **PDM** $_{\mu}$, which is different from the **unconstrained case**, i.e. $f(\mathbf{x}^k) \geq f^*$.
- ▶ The **convergence rate** is **optimal** in the sense of **black-box first order methods**.

ADMM variant

ADMM was originally developed to solve a special case of (40):

$$f^* := \min_{\mathbf{x} \in \mathcal{X}} \{f(\mathbf{x}) : \mathbf{A}\mathbf{x} = \mathbf{b}\},$$

where f and \mathcal{X} is decomposable with $g = 2$.

Problem setting: When f and \mathcal{X} are 2-decomposable

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Augmented Lagrangian smoother

- ▶ When $\tilde{\mathbf{x}}_1^k$ and $\tilde{\mathbf{x}}_2^{k+1}$ are available, we use

$$d_{\gamma\beta}(\mathbf{z}) := \frac{\gamma}{2} [\|\mathbf{A}_1\mathbf{x}_1 + \mathbf{A}_2\tilde{\mathbf{x}}_2^k - \mathbf{b}\|^2 + \|\mathbf{A}_1\tilde{\mathbf{x}}_1^{k+1} + \mathbf{A}_2\mathbf{x}_2 - \mathbf{b}\|^2] + \frac{\beta}{2} \|\lambda\|^2.$$

to smooth the gap function $G(\mathbf{z}) := \max_{\hat{\mathbf{z}} \in \mathcal{X} \times \mathbb{R}^n} \{f(\mathbf{x}) - f(\hat{\mathbf{x}}) + M(\mathbf{z})^T(\mathbf{z} - \hat{\mathbf{z}})\}.$

- ▶ Modify the (1P2D) scheme to obtain a new variant of ADMM by alternating the computation of $\mathbf{x}_\gamma^*(\lambda)$.

ADMM variant

By **alternating** the step of $\mathbf{x}_\gamma^*(\lambda)$, the main step of the **new ADMM variant** becomes:

New ADMM scheme (ADMM₁)

$$\left\{ \begin{array}{ll} \hat{\lambda}^k & := (1 - \tau_k) \bar{\lambda}^k + \tau_k \lambda_{\beta_k}^* (\bar{\mathbf{x}}^k) \\ \tilde{\mathbf{x}}_1^{k+1} & := \operatorname{argmin}_{\mathbf{x}_1 \in \mathcal{X}_1} \left\{ f_1(\mathbf{x}_1) + (\mathbf{A}_1^T \hat{\lambda}^k)^T \mathbf{x}_1 + (\gamma_k/2) \|\mathbf{A}_1 \mathbf{x}_1 + \mathbf{A}_2 \tilde{\mathbf{x}}_2^k - \mathbf{b}\|^2 \right\} \\ \tilde{\mathbf{x}}_2^{k+1} & := \operatorname{argmin}_{\mathbf{x}_2 \in \mathcal{X}_2} \left\{ f_2(\mathbf{x}_2) + (\mathbf{A}_2^T \hat{\lambda}^k)^T \mathbf{x}_2 + (\gamma_k/2) \|\mathbf{A}_1 \tilde{\mathbf{x}}_1^{k+1} + \mathbf{A}_2 \mathbf{x}_2 - \mathbf{b}\|^2 \right\} \\ \bar{\mathbf{x}}^{k+1} & := (1 - \tau_k) \bar{\mathbf{x}}^k + \tau_k \tilde{\mathbf{x}}^{k+1}, \text{ where } \tilde{\mathbf{x}}^{k+1} := (\tilde{\mathbf{x}}_1^{k+1}, \tilde{\mathbf{x}}_2^{k+1}) \\ \bar{\lambda}^{k+1} & := \hat{\lambda}^k + (\gamma_k/2) (\mathbf{A} \tilde{\mathbf{x}}^{k+1} - \mathbf{b}). \end{array} \right.$$

ADMM variant

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Convergence of ADMM₁ [14]

Assumptions:

- ▶ Let $\{(\mathbf{x}^k, \lambda^k)\}_{k \geq 0}$ be the sequence generated by **PDM** using **ADMM₁**.
- ▶ Let $\gamma_k = \gamma_0 := \frac{2\sqrt{2}\|\mathbf{A}\|}{K+3}$ and $\beta_{k+1} := (1 - \tau_k)\beta_k$ for $k = 0, \dots, K$.

Conclusion:

$$\left\{ \begin{array}{lll} \|\mathbf{A} \bar{\mathbf{x}}^K - \mathbf{b}\| & \leq & \frac{2\sqrt{2}\|\mathbf{A}\|(\bar{D}_{\Lambda^*} + \bar{D}_{\mathcal{X}})}{K+3} \\ -D_{\Lambda^*} \|\mathbf{A} \bar{\mathbf{x}}^K - \mathbf{b}\| \leq f(\bar{\mathbf{x}}^K) - f^* & \leq & \frac{2\sqrt{2}\|\mathbf{A}\|}{K+3} (\bar{D}_{\mathcal{X}})^2 \end{array} \right.$$

where $\bar{D}_{\mathcal{X}} := 2 \max_{\mathbf{x}, \hat{\mathbf{x}} \in \mathcal{X}} \|\mathbf{A}(\mathbf{x} - \hat{\mathbf{x}})\|$.

Preconditioned ADMM variant

When f_1 and f_2 are **proximally tractable** and \mathcal{X}_1 and \mathcal{X}_2 are **absent**

- ▶ We can **linearize** the **quadratic terms** in lines 2 and 3 of ADMM₁.
- ▶ Then, by using the **gradient step**, we obtain a **preconditioned ADMM** variant.

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- ▶ Then, by using the **gradient step**, we obtain a **preconditioned ADMM** variant.

Preconditioned ADMM variant (PADMM₁)

$$\begin{cases} \tilde{\mathbf{x}}_1^{k+1} &:= \text{prox}_{\gamma_k^{-1} \alpha_{1k} f_1}(\mathbf{g}_1^k + \gamma_k^{-1} \mathbf{A}_1^T \lambda^k) \\ \tilde{\mathbf{x}}_2^{k+1} &:= \text{prox}_{\gamma_k^{-1} \alpha_{2k} f_2}(\mathbf{g}_2^k + \gamma_k^{-1} \mathbf{A}_2^T \lambda^k) \end{cases}$$

where \mathbf{g}_1^k and \mathbf{g}_2^k are the **gradient step** of the **quadratic term** computed as

$$\begin{cases} \mathbf{g}_1^k &:= \tilde{\mathbf{x}}_1^k - \alpha_{1k} \mathbf{A}_1^T (\mathbf{A}_1 \tilde{\mathbf{x}}_1^k + \mathbf{A}_2 \tilde{\mathbf{x}}_2^k - \mathbf{b}) \\ \mathbf{g}_2^k &:= \tilde{\mathbf{x}}_2^k - \alpha_{2k} \mathbf{A}_2^T (\mathbf{A}_1 \tilde{\mathbf{x}}_1^{k+1} + \mathbf{A}_2 \tilde{\mathbf{x}}_2^k - \mathbf{b}). \end{cases}$$

Here α_{1k} and α_{2k} are given step-sizes.

There are at least two ways of computing the step-sizes:

- ▶ **Constant step size:** We can take $\alpha_{1k} := \|\mathbf{A}_1\|^{-1}$ and $\alpha_{2k} := \|\mathbf{A}_2\|^{-1}$.
- ▶ **Adaptive step-size:** α_{1k} and α_{2k} are computed from the **exact line-search condition** of the form:

$$\alpha := \arg \min_{\alpha > 0} \xi(\mathbf{u}^k - \alpha \nabla \xi(\mathbf{u}^k))$$

where \mathbf{u} can be \mathbf{x}_1 or \mathbf{x}_2 , and ξ is the quadratic function of \mathbf{x}_1 or \mathbf{x}_2 in PADMM₁.

Convergence of PAMMM

Convergence of PADMM₁ [14]

Assumptions:

- ▶ Let $\{(\mathbf{x}^k, \lambda^k)\}_{k \geq 0}$ be the sequence generated by PDM using PADMM₁.
- ▶ Let $\gamma_k = \gamma_0 := \frac{2\sqrt{2}\|\mathbf{A}\|}{K+3}$ and $\beta_{k+1} := (1 - \tau_k)\beta_k$ for $k = 0, \dots, K$.

Conclusion:

$$\begin{cases} \|\mathbf{A}\bar{\mathbf{x}}^K - \mathbf{b}\| & \leq \frac{2\sqrt{2}\|\mathbf{A}\|(D_{\Lambda^*} + \bar{D}_{\mathcal{X}})}{K+3} \\ -D_{\Lambda^*}\|\mathbf{A}\bar{\mathbf{x}}^K - \mathbf{b}\| \leq f(\bar{\mathbf{x}}^K) - f^* & \leq \frac{2\sqrt{2}\|\mathbf{A}\|}{K+3}(\bar{D}_{\mathcal{X}})^2 \end{cases}$$

where $\bar{D}_{\mathcal{X}} := 4 \max_{\mathbf{x}, \hat{\mathbf{x}} \in \mathcal{X}} \|\mathbf{x} - \hat{\mathbf{x}}\|$.

Enhancements of the PDM algorithm:

- ▶ There is a **freedom** of choosing the **center point** \mathbf{x}_c for computing $\mathbf{x}_{\gamma}^*(\lambda)$.
 - ▶ \mathbf{x}_c can be fixed for all the iterations.
 - ▶ One can choose \mathbf{x}_c as the previous iteration, i.e., $\mathbf{x}_c := \mathbf{x}_{\gamma_k}^*(\lambda^{k-1})$.
 - ▶ Or choose \mathbf{x}_c adaptively as in the PADMM variant.
- ▶ The **smoothness parameter** γ can be increased as long as the objective values does not increase substantially.
 - ▶ When \mathbf{x}_c is adaptively chosen, we can slightly increase γ_k as $\gamma_{k+1} := c\gamma_k$, for e.g., $c_k := 1.05$.

Comparison

We summarize the convergence rate of 5 different methods and the assumptions where the methods use in the following table:

- ▶ The average sequence $\{\hat{\mathbf{x}}^k\}$ is computed as $\hat{\mathbf{x}}^k := (k+1)^{-1} \sum_{j=0}^k \mathbf{x}^j$.
- ▶ Convergence guarantee using this sequence is referred to as an ergodic convergence.

Method name	Assumptions	Convergence	References
ADMM	≤ 2 -decomposable	$\mathcal{O}(1/k)$ on the joint $(\mathbf{x}^k, \mathbf{y}^k)$ using a gap function	[2, 8, 9]
[Fast] ADMM	≤ 2 -decomposable and f_1 or $f_2 \in \mathcal{F}_\mu$	$[\mathcal{O}(1/k^2)] \mathcal{O}(1/k)$ on the dual-objective	[6]
Decomposition methods with 1P2D and 2P1D	p -decomposable	$ f(\mathbf{x}^k) - f^* \leq \mathcal{O}(1/k)$ and $\ \mathbf{A}\mathbf{x}^k - \mathbf{b}\ _2 \leq \mathcal{O}(1/k)$ (<i>non-ergodic</i>)	[14]
	p -decomposable and $f_i \in \mathcal{F}_\mu$	$ f(\mathbf{x}^k) - f^* \leq \mathcal{O}(1/k^2)$, $\ \mathbf{A}\mathbf{x}^k - \mathbf{b}\ _2 \leq \mathcal{O}(1/k^2)$, and $\ \mathbf{x}^k - \mathbf{x}^*\ _2 \leq \mathcal{O}(1/k)$ (<i>non-ergodic</i>)	
New ADMM	≤ 2 -decomposable	$ f(\mathbf{x}^k) - f^* \leq \mathcal{O}(1/k)$ and $\ \mathbf{A}\mathbf{x}^k - \mathbf{b}\ _2 \leq \mathcal{O}(1/k)$ (<i>non-ergodic</i>)	[14]
New preconditioned ADMM	≤ 2 -decomposable	$ f(\mathbf{x}^k) - f^* \leq \mathcal{O}(1/k)$ and $\ \mathbf{A}\mathbf{x}^k - \mathbf{b}\ _2 \leq \mathcal{O}(1/k)$ (<i>non-ergodic</i>)	[14]

Example 1: Group sparse recovery

Sparse recovery

- ▶ Let $\mathcal{I} := \{1, \dots, p\}$ be the **set of indices**. Let $\mathcal{G} := \{\mathcal{G}_1, \dots, \mathcal{G}_g\}$ be the **set of g groups** $\mathcal{G}_i \subseteq \mathcal{I}$ and $\mathcal{I} \subseteq \cup_{i=1}^g \mathcal{U}_i$.
- ▶ For given **group** \mathcal{G}_i , and a vector $\mathbf{x} \in \mathbb{R}^p$, we use $\mathbf{x}_{\mathcal{G}_i} = \{x_j : j \in \mathcal{G}_i\}$.
- ▶ For fixed **group structure** \mathcal{G} , $\mathbf{x} \in \mathbb{R}^p$ is called **group sparse vector** if the **number of groups** in \mathcal{G} is **small**.
- ▶ Given a **linear operator** \mathbf{A} and an **observed/measurement** vector $\mathbf{b} \in \mathbb{R}^n$. We want to recover the **group sparse** input vector $\mathbf{x} \in \mathbb{R}^p$ such that $\mathbf{b} = \mathbf{A}\mathbf{x}$.

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- ▶ For given **group** \mathcal{G}_i , and a vector $\mathbf{x} \in \mathbb{R}^p$, we use $\mathbf{x}_{\mathcal{G}_i} = \{x_j : j \in \mathcal{G}_i\}$.
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Optimization formulation

$$\begin{aligned} \min_{\mathbf{x} \in \mathbb{R}^p} \quad & \sum_{\mathcal{G}_i \in \mathcal{G}} \|\mathbf{x}_{\mathcal{G}_i}\|_2 \\ \text{s.t.} \quad & \mathbf{A}\mathbf{x} = \mathbf{b}. \end{aligned} \tag{56}$$

Here, $f(\mathbf{x}) := \sum_{\mathcal{G}_i \in \mathcal{G}} \|\mathbf{x}_{\mathcal{G}_i}\|_2$ and $\mathcal{X} := \mathbb{R}^p$. This problem possesses two common structures: **decomposability** and **tractable proximity**.

When $g = p$ and $\mathcal{G}_i = \{i\}$, (56) reduces to the well-known **linear sparse recovery problem** (basis pursuit):

$$\min_{\mathbf{x} \in \mathbb{R}^p} \|\mathbf{x}\|_1 \quad \text{s.t.} \quad \mathbf{A}\mathbf{x} = \mathbf{b}. \tag{57}$$

Example 1: Group sparse recovery - Numerical results

Algorithm configuration:

- ▶ Assume that (56) is **constrained** by a **boxed constraint** $\mathbf{x} \in \mathcal{X} := [\mathbf{l}, \mathbf{u}]$.
- ▶ The **Bregman distance** is chosen as $d(\mathbf{x}, \mathbf{x}_c) := (1/2)\|\mathbf{x} - \mathbf{x}_c\|_2^2$ and $\mathbf{x}_c = \mathbf{0} \in [\mathbf{l}, \mathbf{u}]$.
- ▶ $\beta_0 = \gamma_0 = \|\mathbf{A}\|$ in 2P1D and $\gamma_0 := 2\sqrt{2}\|\mathbf{A}\|/(K+1)$ with $K = 10^4$.

Data generation:

- ▶ $p = 1024$, $n = 341$ and $g = 128$.
- ▶ \mathbf{A} is a random matrix generated via the **standard Gaussian distribution**.
- ▶ $\mathbf{b} := \mathbf{A}\mathbf{x}^{\natural}$, where \mathbf{x}^{\natural} is a 128-group sparse vector.
- ▶ The group \mathfrak{G} is also generated **randomly**.
- ▶ $\mathbf{l} := \min(\mathbf{x}^{\natural})$ and $\mathbf{u} := \max(\mathbf{x}^{\natural})$.

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Data generation:

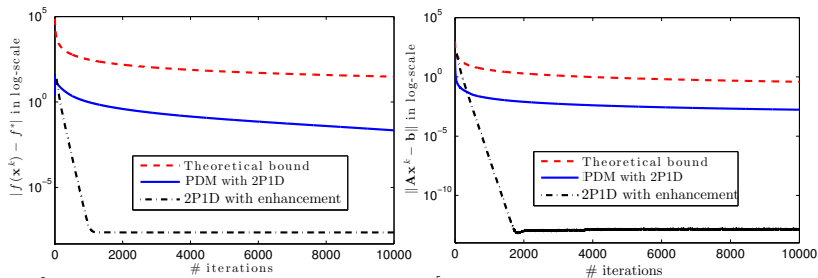
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Comparison We compare the following three quantities:

- ▶ The **theoretical bounds** given in the right-hand side of the convergence theorem
- ▶ The **PDM algorithm** with 2P1D or 1P2D - i.e., follow the theory.
- ▶ The 2P1D or 1P2D with **enhancement** - i.e., updating the parameter γ_k by $\gamma_{k+1} := 1.05\gamma_k$ and using the adaptive center point \mathbf{x}_c as in **PADMM**.

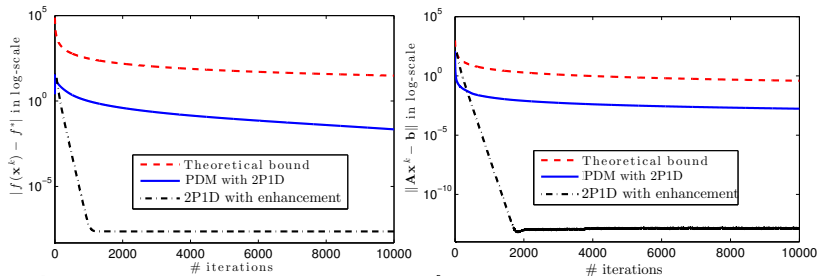
Example 1: Theoretical bounds vs actual performance

The performance of two variants of PDM using 2P1D

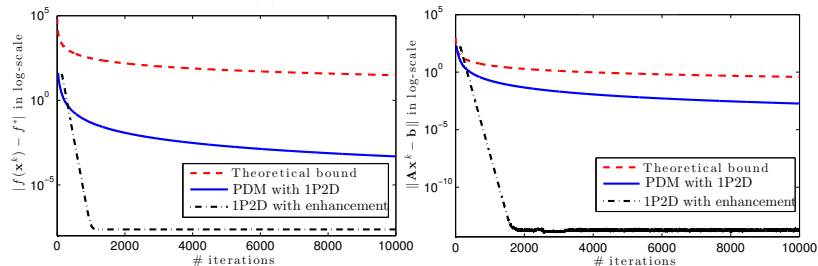


Example 1: Theoretical bounds vs actual performance

The performance of two variants of PDM using 2P1D



The performance of two variants of PDM using 1P2D



Example 2: Image processing

Problem (Imaging denoising/deblurring)

Our goal is to obtain a clean image \mathbf{x} given “dirty” observations $\mathbf{b} \in \mathbb{R}^{n \times 1}$ via $\mathbf{b} = \mathcal{A}(\mathbf{x}) + \mathbf{w}$, where \mathcal{A} is a linear operator, which, e.g., captures camera blur as well as image subsampling, and \mathbf{w} models Gaussian perturbations.

Optimization formulation

$$\text{Gaussian : } \min_{\mathbf{z} \in \mathcal{Z}} \left\{ (1/2) \|\mathcal{A}(\mathbf{z}) - \mathbf{b}\|_2^2 + \rho \|\mathbf{z}\|_{\text{TV}} \right\} \quad (58)$$

where $\|\mathbf{z}\|_{\text{TV}} := \sum_{i,j} |z_{i,j+1} - z_{i,j}| + |z_{i+1,j} - z_{i,j}|$, $\rho > 0$ is a regularization parameter and $\mathcal{Z} := [0, 255]^{n \times p}$.

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Reformulation: Since $\|\mathbf{z}\|_{\text{TV}} = \|\mathbf{D}\mathbf{z}\|_1$ for a given matrix \mathbf{D} . By letting $\mathbf{r} = \mathbf{D}\mathbf{z}$, we can reformulate (58) as

$$\begin{aligned} \min_{\mathbf{z} \in \mathcal{Z}, \mathbf{r}} \quad & \left\{ (1/2) \|\mathcal{A}(\mathbf{z}) - \mathbf{b}\|_2^2 + \rho \|\mathbf{r}\|_1 \right\} \\ \text{s.t.} \quad & \mathbf{D}\mathbf{z} - \mathbf{r} = 0. \end{aligned}$$

This problem is a **constrained convex minimization problem** with **2-decomposable objective** $f(\mathbf{x}) := (1/2) \|\mathcal{A}(\mathbf{z}) - \mathbf{b}\|_2^2 + \rho \|\mathbf{r}\|_1$ and $\mathbf{x} := (\mathbf{z}, \mathbf{r})$.

Example 2: Image processing - Input data

The original image and Gaussian noise image

original image



input: noise image



Data generation:

- ▶ The original image is **filtered** with a **multidimensional filter** H of size 9×9 (circulant).
- ▶ 0.5% Gaussian noise is **added** to the output.

Parameter configuration:

- ▶ The number of iterations: 200 and the relative tolerance: 10^{-8} .

Example 2: Image processing - Numerical results

The performance of the **new ADMM variant** of PDM for two values of ρ .

output: 1P2D with $\rho = 100$



output: 1P2D with $\rho = 50$



- ▶ $f(\mathbf{x}^h) = 138097.919259$ and 77284.828237 , where \mathbf{x}^h is the **original image**.
- ▶ The **objective values**: $f(\mathbf{x}^k) := 122557.13880$ and $f(\mathbf{x}^k) := 64456.44963$
- ▶ **Relative error** between **original image** to **clean image**: 0.089152 and 0.089167
- ▶ PSNR: 26.011 and 25.994.

Example 3: Binary linear support vector machine

Problem (Binary classification)

Given a sample vector $\mathbf{a} \in \mathbb{R}^p$ and a binary class label vector $\mathbf{b} \in \{-1, +1\}^n$. The *goal* is to find a separating hyperplane $\varphi(\mathbf{a}, \mathbf{z}) := \mathbf{a}^T \mathbf{z} + \mu$ such that

$$b_i = \begin{cases} +1 & \text{if } \varphi(\mathbf{a}, \mathbf{z}) \geq 0 \\ -1 & \text{otherwise} \end{cases}$$

where $\mathbf{z} \in \mathbb{R}^p$ is a weight vector, μ is called a bias.

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Optimization formulation

$$\min_{\mathbf{z} \in \mathbb{R}^p} \left\{ \frac{1}{n} \sum_{i=1}^n \mathcal{H}(b_i, \mathbf{a}_i^T \mathbf{z} + \mu) + \rho \|\mathbf{z}\|_1 \right\} \quad (59)$$

where \mathbf{a}_i is the i -th row of the observed data matrix \mathbf{A} in $\mathbb{R}^{n \times p}$, $\rho > 0$ is a regularization parameter, and \mathcal{H} is the Hinge loss function $\mathcal{H}(s, \tau) := \max\{0, 1 - s\tau\}$.

Constrained reformulation: By introducing a slack variable $\mathbf{r} := \mathbf{A}\mathbf{z} + \mu$, we have

$$\begin{aligned} \min_{\mathbf{z} \in \mathbb{R}^p, \mathbf{r}} \quad & \left\{ \frac{1}{n} \sum_{i=1}^n \mathcal{H}(b_i, r_i) + \rho \|\mathbf{z}\|_1 \right\} \\ \text{s.t.} \quad & \mathbf{A}\mathbf{z} + \mu - \mathbf{r} = \mathbf{0}. \end{aligned}$$

Example 3: Binary linear support vector machine

Testing data

- ▶ Test problems: Two real-world problems a1a and news20 from <http://www.csie.ntu.edu.tw/~cjlin/libsvmtools/datasets/>
- ▶ The data size of a1a: $p = 119$ features and $n = 1605$ data points
- ▶ The data size of news20: $p = 1355191$ features and $n = 19996$ data points
- ▶ The parameter ρ changes from $\rho^{-1} = 10^{-3}$ to $\rho^{-1} = 10^3$.

Comparison: We compare the new PADMM variant with LibSVM.

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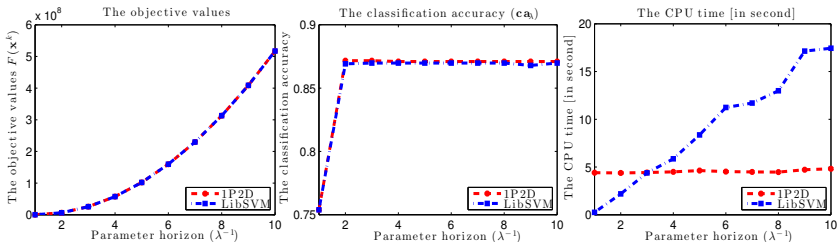
Comparison: We compare the new PADMM variant with **LibSVM**.

Problem	The parameter values									
λ^{-1}	10^{-3}	111.1	222.2	333.3	444.4	555.6	666.7	777.8	888.9	10^3
The accuracy of problem a1a										
(1P2D)	0.7539	0.8717	0.8717	0.8710	0.8710	0.8710	0.8710	0.8710	0.8710	0.8710
LibSVM	0.7539	0.8692	0.8698	0.8698	0.8698	0.8698	0.8698	0.8698	0.8679	0.8698
The CPU time [in second] of problem a1a										
(1P2D)	4.4045	4.3769	4.4246	4.4941	4.6238	4.5175	4.4836	4.4719	4.7179	4.8097
LibSVM	0.2549	2.1909	4.3884	5.8583	8.3662	11.2350	11.7036	12.9832	17.1424	17.4362
The accuracy of problem news20										
(1P2D)	0.5001	0.9987	0.9987	0.9987	0.9987	0.9987	0.9987	0.9987	0.9987	0.9987
LibSVM	0.5001	0.9987	0.9987	0.9987	0.9987	0.9988	0.9988	0.9988	0.9988	0.9988
The CPU time [in second] of problem news20										
(1P2D)	762.31	1023.22	994.64	1043.06	984.24	989.70	1064.33	1073.94	984.47	1018.35
LibSVM	890.26	1440.28	1449.23	1439.77	1434.27	1518.56	1560.38	1557.48	1535.19	1530.71

- ▶ The **accuracy** is computed as $\text{ca}_\rho := 1 - n^{-1} \sum_{j=1}^n [\text{sign}(\mathbf{A}\mathbf{z}^k + \mu)_i \neq b_i]$.

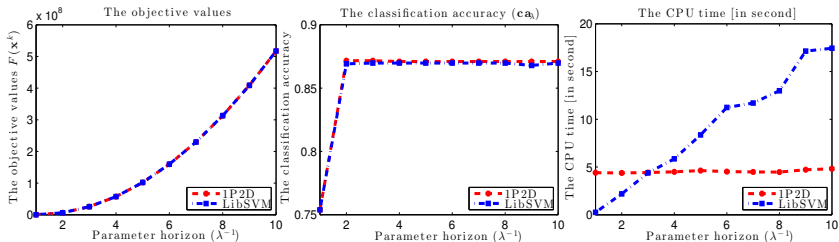
Example 3: Binary linear support vector machine

The results of **two algorithms** on the real-world problem a1a

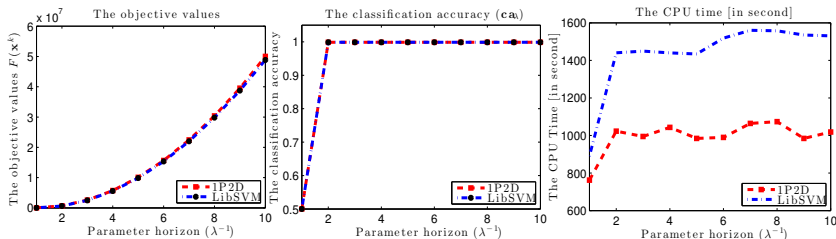


Example 3: Binary linear support vector machine

The results of **two algorithms** on the real-world problem a1a



The results of **two algorithms** on the real-world problem news20



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