





Winter Conference in Statistics 2013

Compressed Sensing

LECTURE #7-8 Algorithms for low-dimensional models



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Convex Algorithms for Low-Dimensional Models



...And, this is how you solve huge-dimensional problems

The classical problem templates

Criteria seen above have the form

Basis pursuit (BP) [Chen, Donoho, Saunders, 1998] $\min_{x} \|x\|_1$ s.t. $\Phi x = u$

BP denoising (BPDN): [Chen, Donoho, Saunders, 1998] $\min_{x} \|x\|_1$ s.t. $\|\Phi x - u\|_2^2 \le \varepsilon$

Also well known: LASSO (least absolute shrinkage/selection operator): [Tibshirani, 1996]

$$\min_{x} \|\Phi x - u\|_{2}^{2} \text{ s.t. } \|x\|_{1} \le \tau$$

 $\|x\|_1 = \sum |[x]_i|$

All can be written as $\widehat{x} \in \arg\min_{x\in\mathbb{R}^N} f_1(x) + f_2(x)$

Convex optimization and proximal algorithms

$$\widehat{x} \in \arg\min_{x \in \mathbb{R}^N} f_1(x) + f_2(x)$$

 $f_1: \mathbb{R}^N o \mathbb{R}$ data fidelity term; convex, smooth.

typically:
$$f_1(x) = \frac{1}{2} \|\Phi x - u\|_2^2$$

 $f_2: \mathbb{R}^N \to \overline{\mathbb{R}} = \mathbb{R} \cup \{+\infty\}$ Convex regularizer (maybe non-smooth; e.g. ℓ_1) (non-convex, later...).

Difficulties: non-smoothness and large dimension ($N \gg 1$)

Constrained vs unconstrained formulations

Constrained optimization formulations

$$\widehat{x} \in \arg\min_{x \in \mathbb{R}^N} f_1(x) \quad (*) \quad \widehat{x} \in \arg\min_{x \in \mathbb{R}^N} f_2(x)$$
s.t. $h(x) \le \nu \quad (*) \quad \text{s.t. } g(x) \le \tau$

can be written as $\widehat{x} \in rgmin_{x\in\mathbb{R}^N} f_1(x) + f_2(x)$

...using indicator functions: $\iota_S(x) = \begin{cases} 0 & \Leftarrow & x \in S \\ +\infty & \Leftarrow & x \notin S \end{cases}$

Example: (*) same as

$$\widehat{x} \in \arg\min_{x\in\mathbb{R}^N} f_1(x) + \iota_{\{x:g(x)\leq\nu\}}(x)$$

Classical example: the LASSO: $\min \|\Phi x - u\|_2^2$ s.t. $\|x\|_1 \leq \tau$

Convex and strictly convex sets

 $\mathcal{S} \text{ is convex if } x, x' \in \mathcal{S} \ \Rightarrow \forall \lambda \in [0,1] \ \lambda x + (1-\lambda)x' \in \mathcal{S}$



 \mathcal{S} is strictly convex if $x, x' \in \mathcal{S} \Rightarrow \forall \lambda \in (0, 1) \ \lambda x + (1 - \lambda) x' \in \operatorname{int}(\mathcal{S})$



Convex and strictly convex functions

Extended real valued function: $f : \mathbb{R}^N \to \overline{\mathbb{R}} = \mathbb{R} \cup \{+\infty\}$ Domain of a function: $\operatorname{dom}(f) = \{x : f(x) \neq +\infty\}$ f is a convex function if $\forall \lambda \in [0, 1], x, x' \in \operatorname{dom}(f) \ f(\lambda x + (1 - \lambda)x') \leq \lambda f(x) + (1 - \lambda)f(x')$ f is a strictly convex function if $\forall \lambda \in (0, 1), x, x' \in \operatorname{dom}(f) \ f(\lambda x + (1 - \lambda)x') < \lambda f(x) + (1 - \lambda)f(x')$



Convexity, coercivity, and minima

$$f: \mathbb{R}^N \to \bar{\mathbb{R}} = \mathbb{R} \cup \{+\infty\}$$

f is coercive if $\lim_{\|x\|\to+\infty}f(x)=+\infty$

if f is coercive, then $G\equiv \arg\min_x f(x)\,$ is a non-empty set

if f is strictly convex, then $G\,$ has at most one element



Euclidean projections on convex sets

Our problem:
$$\widehat{x} \in \arg \min_{x \in \mathbb{R}^n} f_1(x) + f_2(x)$$

consider $f_2(x) = \iota_{\mathcal{S}}(x) = \begin{cases} 0 \iff x \in \mathcal{S} \\ +\infty \iff x \notin \mathcal{S} \end{cases}$
(convex if \mathcal{S} is convex)

and
$$f_1(x) = rac{1}{2} \|u-x\|_2^2$$
 (strictly convex)

Projected gradient algorithm

Our problem:
$$\widehat{x} \in \arg\min_{x\in\mathbb{R}^n} f_1(x) + f_2(x)$$

with $f_2(x) = \iota_{\mathcal{S}}(x)$ (\mathcal{S} is a convex set)
and f_1 some function, e.g., $f_1(x) = \frac{1}{2} \|\Phi x - u\|_2^2$

Projected gradient algorithm:

Detour: majorization-minimization (MM)

Problem:
$$\widehat{x} \in \arg\min_{x\in\mathbb{R}^n} f(x)$$

 $Q(x,x_k)$ is a majorizer of f at x_k

 $Q(x, x_k) \ge f(x), \quad Q(x_k, x_k) = f(x_k)$



MM algorithm: $x_{k+1} = \arg\min_{x} Q(x, x_k)$

monotonicity:

$$f(x_{k+1}) \leq Q(x_{k+1}, x_k)$$
$$\leq Q(x_k, x_k)$$
$$= f(x_k)$$

Projected gradient from majorization-minimization

Our problem:
$$\hat{x} \in \arg \min_{x \in \mathbb{R}^n} f_1(x) + f_2(x)$$

with $f_2(x) = \iota_S(x)$ (S is a convex set)
and f_1 has L -Lipschitz gradient
 $\|\nabla f_1(x) - \nabla f_1(x')\| \le L \|x - x'\|$
e.g. $f_1(x) = \frac{1}{2} \|\Phi x - u\|_2^2 \Rightarrow L = \lambda_{\max}(\Phi^T \Phi) = \|\Phi\|_2^2$
Hessian of f_1

...a separable approximation of f_1

$$Q(x, x_k) = f_1(x_k) + (x - x_k)^T \nabla f_1(x_k) + \frac{1}{2\beta_k} ||x - x_k||_2^2$$

Projected gradient from majorization-minimization

Our problem:
$$\widehat{x} \in rgmin_{x\in\mathbb{R}^n} f_1(x) + \iota_\mathcal{S}(x)$$

Separable approximation of f_1

$$egin{aligned} Q(x,x_k) &= f_1(x_k) + (x-x_k)^T
abla f_1(x_k) + rac{1}{2eta_k} \|x-x_k\|_2^2 \ Q(x,x_k) & ext{ is a majorizer of } f_1 ext{, if } eta_k < rac{1}{L} \ Q(x,x_k) + \iota_\mathcal{S}(x) & ext{ is a majorizer } f_1(x) + \iota_\mathcal{S}(x) \end{aligned}$$

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MM algorithm:

$$\begin{split} x_{k+1} &= \arg\min_{x} Q(x, x_{k}) + \iota_{S}(x) \\ &= \arg\min_{x} \left. \frac{1}{2\beta_{k}} \right\| x - \left(x_{k} - \beta_{k} \nabla f_{1}(x_{k}) \right) \right\|_{2}^{2} + \iota_{S}(x) \\ &= P_{S} \left(x_{k} - \beta_{k} \nabla f_{1}(x_{k}) \right) \quad \text{...projected gradient.} \end{split}$$

Proximity operators

Our problem:
$$\widehat{x} \in \arg\min_{x\in\mathbb{R}^n} f_1(x) + f_2(x)$$

with $f_2\,$ a convex function

and
$$f_1(x) = rac{1}{2} \|u-x\|_2^2$$
 (strictly convex)

$$\widehat{x} = \arg\min_{x \in \mathbb{R}^n} \frac{1}{2} \|u - x\|_2^2 + f_2(x) \equiv \operatorname{prox}_{f_2}(u)$$

Proximity operator [Moreau 62], [Combettes 01].

Generalizes the notion of Euclidean projection.

Proximity operators (linear)

$$\operatorname{prox}_{f}(u) = \arg\min_{x \in \mathbb{R}^{n}} \frac{1}{2} \|u - x\|_{2}^{2} + f(x) \quad (\mathbb{R}^{N} \to \mathbb{R}^{N})$$

Classical cases: squared ℓ_2 regulizer $f(x) = \frac{\tau}{2} ||x||_2^2$

$$\operatorname{prox}_{f}(u) = \arg\min_{x \in \mathbb{R}^{n}} \frac{1}{2} \|u - x\|_{2}^{2} + \frac{\tau}{2} \|x\|_{2}^{2} = \frac{u}{1 + \tau}$$

squared ℓ_2 regularizer with "analysis" operator $f(x) = rac{ au}{2} \|Dx\|_2^2$

$$prox_f(u) = \arg\min_{x \in \mathbb{R}^n} \frac{1}{2} ||u - x||_2^2 + \frac{\tau}{2} ||Dx||_2^2$$
$$= (I + \tau D^T D)^{-1} u$$

if D is a circulant matrix, $O(N \log N)$ cost using the FFT

Proximity operator of the $\,\ell_1\,$ norm

$$\operatorname{prox}_{\tau \|\cdot\|_{1}}(u) = \arg\min_{x \in \mathbb{R}^{n}} \frac{1}{2} \|u - x\|_{2}^{2} + \tau \|x\|_{1}$$



Separable: solve w.r.t. each component: $\min \tau |x| + 0.5(x-u)^2$ Possible approach: write $|x| = \max_{|z| < 1} zx$ $\min_{x} \max_{|z| \le 1} \tau z x + 0.5(x-u)^2 = \max_{|z| \le 1} \min_{x} \tau z x + 0.5(x-u)^2$ $= \max_{|z| \le 1} -0.5\tau^2 z^2 + \tau z u \quad \text{(for } x = u - \tau z \text{)}$ $\arg\max_{|z|\leq 1} -0.5\tau^2 z^2 + \tau z u = \begin{cases} u/\tau & \Leftarrow & |u| \cdot \tau \\ 1 & \Leftarrow & u > \tau \\ -1 & \Leftarrow & u < -\tau \end{cases}$ $\arg\min_{x} \tau |x| + 0.5(x-u)^2 = \begin{cases} 0 & \Leftarrow & |u| \le \tau \\ u - \tau & \Leftarrow & u > \tau \\ u + \tau & \Leftarrow & u < -\tau \end{cases}$

Proximity operator of the ℓ_1 norm: the "soft"

$$\arg\min_{x} \tau |x| + 0.5(x-u)^{2} = \begin{cases} 0 & \Leftarrow & |u| \leq \tau \\ u - \tau & \Leftarrow & u > \tau \\ u + \tau & \Leftarrow & u < -\tau \end{cases}$$
soft
thresholding
$$-\tau$$

$$\tau \quad z = \operatorname{sign}(u) \max\{0, |u| - \tau\}$$

$$\equiv \operatorname{soft}(u, \tau) = \operatorname{prox}_{\tau |\cdot|}$$

(for vectors, $\operatorname{soft}(u,\tau)$ is applied component-wise)

$$p$$
-th power of ℓ_p norms $\|x\|_p^p = \sum_i |[x]_i|^p$ closed form prox for $\ p \in \left\{1,2,\frac{4}{3},\frac{3}{2},3,4\right\}$ [Combettes, Wajs, 2005]

Dual norms, proximity operators, and projections

Dual norm: some norm, $\|\cdot\|: \mathbb{R}^N \to \mathbb{R}_+$

its dual norm:
$$||x||^* = \max_{||z|| \le 1} \langle x, z \rangle$$

Dual norm of $\|\cdot\|_p$ is $\|\cdot\|_q$, where $\frac{1}{p} + \frac{1}{q} = 1$ Hölder conjugates

... simple corollary of Hölder's inequality: $x^T z \leq \|x\|_p \|z\|_q$

Examples of Hölder conjugates: $(2,2), (1,+\infty), (3/2,3), ...$

These concepts are related through:

$$\operatorname{prox}_{\|\cdot\|}(u) = u - P_{\{x:\|x\|^* \le 1\}}(u)$$

[Combettes, Wajs, 2005]

Dual norms, proximity operators, and projections $\operatorname{prox}_{\tau \parallel \cdot \parallel}(u) = u - P_{\{x: \parallel x \parallel^* < \tau\}}(u)$

This relation underlies our earlier derivation of $\operatorname{Prox}_{\|\cdot\|_1}$ $\operatorname{prox}_{\tau\|\cdot\|_1}(u) = u - P_{\{x:\|x\|_{\infty} \leq \tau\}}(u)$ $\|x\|_{\infty} = \max\{|[x]_i|\}$

It's all separable,

$$prox_{\tau|\cdot|}(u) = u - P_{\{x:|x| \le \tau\}}(u)$$

$$= u - \begin{cases} u & \Leftarrow |u| \leq \tau \\ -\tau & \Leftarrow u < -\tau \\ \tau & \Leftarrow u > \tau \end{cases} \xrightarrow{-\tau} \begin{bmatrix} \operatorname{soft}(u, \tau) \\ & & \\ \tau & & \\ = \operatorname{soft}(u, \tau) \end{cases}$$

Dual norms, proximity operators, and projections

$$\operatorname{prox}_{\tau \|\cdot\|}(u) = u - P_{\{x: \|x\|^* \le \tau\}}(u)$$

This relation allows deriving $prox_{\|\cdot\|_{\infty}}$ and $\;prox_{\|\cdot\|_2}$

$$\begin{split} & \operatorname{prox}_{\|\cdot\|_{\infty}}(u) = u - P_{\{x:\|x\|_{1} \leq \tau\}}(u) & \text{projection on the } \ell_{1} \\ & \text{ball of radius } \tau \\ & P_{\{x:\|x\|_{2} \leq \tau\}}(u) \end{split} \quad & O(n \log n) \end{split}$$

$$= u - \begin{cases} u & \Leftarrow \|u\|_2 \leq \tau \\ \tau u / \|u\|_2 & \Leftarrow \|u\|_2 > \tau \\ = \frac{u}{\|u\|_2} \max\{0, \|u\|_2 - \tau\} \end{cases}$$

vector soft thresholding

Proximity operators of atomic norms

$$\operatorname{prox}_{\tau \|\cdot\|}(u) = u - P_{\{x: \|x\|^* \le \tau\}}(u)$$

These relation allows deriving prox operators of atomic norms:

$$||x||_{\mathcal{A}} = \inf\{t > 0 : x \in t \operatorname{conv}(\mathcal{A})\}$$

The dual of an atomic norm ball:

$$\|x\|_{\mathcal{A}}^{*} = \max_{\|z\|_{\mathcal{A}} \le 1} \langle z, x \rangle = \max_{z \in \operatorname{conv}(\mathcal{A})} \langle z, x \rangle$$
$$= \max\{\langle a, x \rangle, \ a \in \mathcal{A}\}$$

$$P_{\{x:\|x\|_{\mathcal{A}}^* \le \tau\}}(u) = \arg\min_{\langle a,x \rangle \le \tau, \,\forall_{a \in \mathcal{A}}} \|u - x\|_2^2$$

 $\operatorname{prox}_{\tau \|\cdot\|_{\mathcal{A}}}(u) = u - \arg\min_{\langle a,x \rangle \leq \tau, \, \forall_{a \in \mathcal{A}}} \|u - x\|_{2}^{2}$

Proximity operators of atomic norms: ℓ_1

Deriving $\ prox_{\tau \|\cdot\|_1}$ from the atomic norm view



 $||x||_{\mathcal{A}}^* = \max\{\langle a, x \rangle, \ a \in \mathcal{A}\} = \max\{|[x]_i|\} = ||x||_{\infty}$ $\operatorname{prox}_{\tau \parallel \cdot \parallel_1}(u) = u - P_{\{x: \parallel x \parallel_{\infty} \le \tau\}}(u)$ $= \operatorname{soft}(x, \tau)$

Proximity operators of atomic norms: ℓ_∞

Deriving $prox_{\tau \|\cdot\|_{\infty}}$ from the atomic norm view

$$\begin{split} \|x\|_{\infty} &= \|x\|_{\mathcal{A}} \qquad \mathcal{A} = \left\{ \begin{bmatrix} 1\\1\\\vdots\\1 \end{bmatrix}, \begin{bmatrix} -1\\1\\\vdots\\1 \end{bmatrix}, \begin{bmatrix} -1\\-1\\\vdots\\1 \end{bmatrix}, \begin{bmatrix} -1\\-1\\\vdots\\1 \end{bmatrix}, \dots, \begin{bmatrix} -1\\-1\\\vdots\\1 \end{bmatrix} \right\} \\ &= \{-1, +1\}^{N} \\ \|\mathcal{A}\| = 2^{N} \\ \|\mathcal{A}\| = 2^{N} \\ \|x\|_{\mathcal{A}}^{*} &= \max\{\langle a, x \rangle, \ a \in \mathcal{A}\} = \sum_{i=1}^{N} |[x]_{i}| = \|x\|_{1} \\ \operatorname{prox}_{\tau\|\cdot\|_{\infty}}(u) = u - P_{\{x:\|x\|_{1} \leq \tau\}}(u) \end{split}$$

Proximity of atomic norms: matrix nuclear norm

Matrix nuclear norm:
$$\|X\|_* = \sum_i \sigma_i(X) = \sum_i \sqrt{\lambda_i(X^T X)}$$

$$||X||_* = ||X||_{\mathcal{A}} \qquad \mathcal{A} = \{Z : \operatorname{rank}(Z) = 1, ||Z||_F = 1\}$$
$$\operatorname{rank}(Z) = |\{\sigma_i(Z) \neq 0\}|$$

Frobenius norm
$$\|Z\|_F^2 = \sum_{ij} [Z]_{ij}^2 = \sum_i \sigma_i^2(Z)$$

$$\begin{split} \|X\|_{\mathcal{A}}^* &= \max\{\langle Z, X \rangle, \ Z \in \mathcal{A}\} \\ &= \max\left\{\sum_i \sigma_i(Z)\sigma_i(X), \ \operatorname{rank}(Z) = 1, \sum_i \sigma_i^2(Z) = 1\right\} \\ &= \sigma_{\max}(X) \ = \ \|X\|_2 \quad \text{spectral norm} \end{split}$$

Proximity of atomic norms: matrix nuclear norm

Euclidean matrix projection:
$$P_{\mathcal{S}}(X) = \arg\min_{Z \in \mathcal{S}} ||Z - X||_F^2$$

Note: for any unitary matrix U $(U^T U = I, UU^T = I)$ $\|UM\|_F^2 = \operatorname{trace} \left(M^T U^T U M\right) = \operatorname{trace} \left(M^T A\right) = \|M\|_F^2$

$$\operatorname{prox}_{\tau \, \|\cdot\|_{*}}(X) = X - P_{\{Z: \|Z\|_{2} \le \tau\}}(X)$$

singular value $= U\Lambda V^T - P_{\{Z:\sigma_{\max}(Z) \leq \tau\}}(U\Lambda V^T)$ diagonal matrix

[Lewis, Malick, 2009]

$$= U \operatorname{diag} \left(\operatorname{diag}(\Lambda) - P_{\{x: \|x\|_{\infty} \leq \tau\}} (\operatorname{diag}(\Lambda)) \right) V^{T}$$

 $= U \mathrm{soft}(\Lambda, au) V^T$ singular value thresholding (svt)

Proximity of atomic norms: matrix spectral norm

Matrix spectral norm:
$$||X||_2 = \sigma_{\max}(X)$$

$$||X||_2 = ||X||_{\mathcal{A}} \quad \mathcal{A} = \{Z : Z^T Z = I\} = \{Z : \sigma_i(Z) = 1, \forall_i\}$$

orthogonal matrices

$$\begin{split} \|X\|_{\mathcal{A}}^{*} &= \max\{\langle Z, X \rangle, \ Z \in \mathcal{A}\} \\ &= \max\left\{\sum_{i} \sigma_{i}(Z)\sigma_{i}(X), \ \sigma_{i}(Z) = 1, \ \forall_{i}\right\} \\ &= \sum_{i} \sigma_{i}(X) \ = \|X\|_{*} \quad \text{nuclear norm} \end{split}$$

Proximity of atomic norms: matrix spectral norm



$$\begin{aligned} \operatorname{prox}_{\tau \parallel \cdot \parallel_{2}}(X) &= X - P_{\{Z: \parallel Z \parallel_{*} \leq \tau\}}(X) \end{aligned}$$
$$= U\Lambda V^{T} - P_{\{Z: \parallel Z \parallel_{*} \leq \tau\}}(U\Lambda V^{T})$$
singular value
$$= U\left(\Lambda - P_{\{Z: \sum_{i} \sigma_{i}(Z) \leq \tau\}}(\Lambda)\right) V^{T}$$

$$= U \operatorname{diag} \left(\operatorname{diag}(\Lambda) - P_{\{x: \|x\|_1 \le \tau\}} (\operatorname{diag}(\Lambda)) \right) V^T$$

residual of projection of the singular values on an ℓ_1 ball of radius τ

Proximity and atomic sets: vectors vs matrices

vectors

matrices

norm	prox	atomic set	norm	prox	atomic set
$\ell_1 \ \ x\ _1$	component soft thresholding	$\mathcal{A} = \{\pm e_i\}$ $ \mathcal{A} = 2N$	nuclear $\ X\ _*$	singular value thresholding	$\mathcal{A}=$ set of all rank 1, norm 1 matrices
$\ell_{\infty} \\ \ x\ _{\infty}$	residual of projection on ℓ_1 ball	$\begin{aligned} \mathcal{A} &= \\ \{\pm 1\}^N \\ \mathcal{A} &= 2^N \end{aligned}$	spectral $\ X\ _2$	residual of s.v. proj. on ℓ_1 ball	$\mathcal{A}=$ set of all orthogonal matrices
$\ell_2 \ \ x\ _2$	vector soft thresholding	$\mathcal{A}=$ set of all vectors with norm 1 $ \mathcal{A} =\infty$	Frobenius $\ X\ _F$	matrix soft threshold.	$\mathcal{A}=$ all matrices of unit Frobenius norm.

Proximal algorithms

Back to the problem: $\widehat{x} \in \arg\min_{x\in\mathbb{R}^n} f_1(x) + f_2(x)$ with f_2 a proper convex function and f_1 has a L -Lipschitz gradient; e.g. $f_1(x) = \frac{1}{2} ||\Phi x - u||_2^2$ with $L = \lambda_{\max}(\Phi^*\Phi)$ separable majorizer $(\beta_k < 1/L)$ $Q(x, x_k) = f_1(x_k) + (x - x_k)^T \nabla f_1(x_k) + \frac{1}{2\beta_k} ||x - x_k||_2^2$

majorization-minimization algorithm

$$x_{k+1} = \arg\min_{x} Q(x, x_{k}) + f_{2}(x)$$

= $\arg\min_{x} \frac{1}{2\beta_{k}} \left\| x - (x_{k} - \beta_{k} \nabla f_{1}(x_{k})) \right\|_{2}^{2} + f_{2}(x)$
 $x_{k+1} = \operatorname{prox}_{\beta_{k} f_{2}} \left(x_{k} - \beta_{k} \nabla f_{1}(x_{k}) \right)$

Proximal algorithms: convergence

Problem:
$$\widehat{x} \in \arg\min_{x\in\mathbb{R}^n} \left[f_1(x) + f_2(x) \right]^{f(x)}$$

 f_1 has a L -Lipschitz gradient; e.g. $f_1(x) = \frac{1}{2} ||\Phi x - u||_2^2$
Iterative shrinkage/thresholding (IST) $L = \lambda_{\max}(\Phi^*\Phi)$

(or forward-backward)

$$x_{k+1} = \operatorname{prox}_{\beta_k f_2} \left(x_k - \beta_k \nabla f_1(x_k) \right)$$

f(x)

if $\beta_k < \frac{1}{\tau}$, IST is a majorization-minimization algorithm, thus $f(x_{k+1}) \le f(x_k)$

 $f(x)\geq 0$, thus $\ (f(x_1),\,f(x_2),\,...,f(x_k),\,...)$ converges. Attention: this does **not** imply convergence of $(x_1, ..., x_k, ...)$

Proximal algorithms: convergence

$$\begin{aligned} \widehat{x} \in G &= \arg\min_{x \in \mathbb{R}^n} f_1(x) + f_2(x) \\ \text{IST algorithm: } x_{k+1} &= \operatorname{prox}_{\beta_k f_2} \left(x_k - \beta_k \, \nabla f_1(x_k) \right) \\ \text{if } 0 < \beta_k < \frac{2}{L} \text{, then } (x_1, \, x_2, \, \dots, \, x_k, \, \dots) \\ \text{converges to a point in } G \\ \text{Inexact version: } & \left(\begin{array}{c} e^{\operatorname{errors}} \\ x_{k+1} &= \operatorname{prox}_{\beta_k f_2} \left(x_k - (\beta_k \, \nabla f_1(x_k) + b_k) \right) + a_k \end{aligned} \end{aligned}$$

convergence still guaranteed if

$$\sum_{k=1}^{\infty} \|a_k\| < \infty \qquad \sum_{k=1}^{\infty} \|b_k\| < \infty$$

Results and proofs in [Combettes and Wajs, 2005]

Proximal algorithms: convergence

Convergence rates (for function values) [Beck, Teboulle, 2009]:

$$f(x_k) - f(\widehat{x}) \le \frac{L \|x_0 - \widehat{x}\|_2^2}{2k}$$

Convergence rate for the iterates require further assumptions on f

Proximal algorithms: convergence of iterates

Small
$$l \Rightarrow \rho \lesssim 1 \Rightarrow$$
 slow convergence! [F, Bioucas-Dias, 2007]

Proximal algorithms: convergence of iterates

$$\widehat{x} \in G = \arg\min_{x} \frac{1}{2} \|\Phi x - u\|_{2}^{2} + \tau \|x\|_{1}$$

With $L = \lambda_{\max}(\Phi^*\Phi)$; using a step-size $\beta < 2/L$,

$$x_{k+1} = \operatorname{soft}\left(x_k - \beta \,\Phi^T(\Phi x_k - u), \beta \tau\right)$$

$$\mathcal{Z} \subseteq \{1,2,...,n\}$$
 such that $\widehat{x} \in G \Rightarrow [\widehat{x}]_{\mathcal{Z}} = 0$

Then, after a finite number of iterations: $[x_k]_\mathcal{Z} = [\widehat{x}]_\mathcal{Z} = 0$

After this, Q-linear convergence: $l = \lambda_{\min}(\Phi_{\bar{z}}^* \Phi_{\bar{z}}) > 0$ Optimal choice $\beta = \frac{2}{L+l}$, $\int \rho = \frac{1-\kappa}{1+\kappa}$ $\|x_{k+1} - \hat{x}\| \le \rho \|x_k - \hat{x}\|$ [Hale, Yin, Zhang, 2008]

Slowness and acceleration of IST

Problem:
$$\widehat{x} \in G = \arg \min_{x} \frac{1}{2} \|\Phi x - u\|_{2}^{2} + \tau \|x\|_{1}$$

IST algorithm:
$$x_{k+1} = ext{soft} \Big(x_k - eta \Phi^T (\Phi x_k - u), eta au \Big)$$

IST is **slow**, if Φ is very ill-conditioned and/or au is very small!

Several proposals for accelerated variants of IST

Methods with memory (TwIST, FISTA)

Quasi-Newton methods (SpaRSA)

Continuation, i.e., use a varying au (FPC, SpaRSA)

Memory-based variants of IST: FISTA

Fast IST algortihm (FISTA); based on Nesterov's work (1980's) [Beck, Teboulle, 2009]

 $t_{k+1} = \frac{1 + \sqrt{1 + 4t_k^2}}{2}$ **FISTA** $z_{k+1} = x_k + \frac{t_k - 1}{t_{k+1}} \left(x_k - x_{k-1} \right)$ $x_{k+1} = \operatorname{soft}\left(z_k - \beta \,\Phi^T (\Phi z_k - u), \beta \tau\right)$ IST: $f(x_k) - f(\hat{x}) = O\left(\frac{1}{k}\right) \qquad \left(\le \frac{L \|x_0 - \hat{x}\|_2^2}{2k} \right)$ FISTA: $f(x_k) - f(\widehat{x}) = O\left(\frac{1}{k^2}\right)$
Memory-based variants of IST: twist

Inspired by 2-step methods for linear systems [Frankel, 1950], [Axelsson, 1996]

TwIST (two-step IST): [Bioucas-Dias, F, 2007] $x_{k+1} = (\alpha - \beta)x_k + (1 - \alpha)x_{k-1} + \beta \operatorname{prox}_{f_2} (x_k - \Phi^T (\Phi x_k - u))$



Memory-based variants of IST: twist



Quasi-newton acceleration of IST: SpaRSA

IST:
$$x_{k+1} = \operatorname{prox}_{\beta_k f_2} \left(x_k - \beta_k \nabla f_1(x_k) \right)$$

A Newton step (instead of gradient descent) would be:

$$x_{k+1} = \operatorname{prox}_{\beta_k f_2} \left(x_k - \left[H(x_k) \right]^{-1} \nabla f_1(x_k) \right)$$

Ve! (matrix of second derivatives)

...computationally too expensive! Barzilai-Borwein approach: [Barzilai-Borwein, 1988], [Wright, Nowak, F, 2009]

$$\boxed{\frac{1}{\beta_k}I\simeq H(x_k)}$$

$$\begin{aligned} &\frac{1}{\beta_k} = \arg\min_{\alpha} \|\alpha(x_k - x_{k-1}) - (\nabla f(x_k) - \nabla f(x_{k-1}))\|_2^2\\ &\text{If } f_1(x) = \frac{1}{2} \|\Phi x - u\|_2^2, \text{ then } \beta_k = \frac{\|x_k - x_{k-1}\|_2^2}{\|\Phi(x_k - x_{k-1})\|_2^2} \end{aligned}$$

Acceleration via continuation

IST:
$$x_{k+1} = \operatorname{soft}\left(x_k - \beta \Phi^T(\Phi x_k - u), \beta \tau\right)$$

Slow, if au is small.

Observation: IST (as SpaRSA) benefits from "warm-starting" (being initialized *close* to the minimizer)

Continuation: start with large τ slowly decrease τ while tracking the solution. [F, Nowak, Wright, 2007], [Hale, Yin, Zhang, 2007]

IST + continuation = fixed point continuation (FPC) [Hale, Yin, Zhang, 2007]

Acceleration via continuation



Some speed comparisons

from [Lorenz, 2011]
$$\widehat{x} = rgmin_x rac{1}{2} \|\Phi x - u\|_2^2 + au \|x\|_1$$

 $\Phi = \begin{bmatrix} I \ U \ R \end{bmatrix}$ au = 0.1(512 × 1536) \widehat{x} with 120 non-zeros



IST, GPSR, SpaRSA, FISTA, YALL1, NESTA, fpc

Proximal algorithms for matrices

$$\widehat{M} \in \arg\min_{M \in \mathbb{R}^{n \times n}} \frac{1}{2} \| \Phi(M) - U \|_F^2 + \mu \| M \|_*$$

linear operator ...its adjoint

The proximal algorithm (IST) is as before:

$$X_{k+1} = \operatorname{svt}_{\mu\beta_k} \left(X_k - \beta_k \Phi^* (\Phi(X_k) - U) \right)$$

Matrix completion: $\Phi(X) = X$ (subset of entries) $|\Omega| = p$

ISTAPG (FISTA)
$$n/r$$
 p p/d_r μ iter $\#sv$ erroriter $\#sv$ error100/10566638.21e-037723611.88e-01655131.06e-03200/101566541.05e-0212180962.45e-01812121.02e-03500/104947151.21e-02109002035.91e-011132167.63e-04

Unknown M				 FPC (continuation) - 			APG + continuation		
n/r	p	p/d_r	μ	iter	#sv	error	iter	#sv	error
100/10	5666	3	8.21e-03	429	32	1.06e-03	74	10	1.46e-04
200/10	15665	4	1.05e-02	278	49	4.38e-04	73	10	1.02e-04
500/10	49471	5	1.21e-02	484	125	5.50e-04	72	10	8.06e-05

from [Toh, Yun, 2009]

... the importance of acceleration!

Another class of methods: augmented Lagrangian

The problem: $\min_{x} f(x)$ s.t. $\Phi x = u$

Penalty parameter

The augmented Lagrangian (AL)

$$L_{\mu}(x,\lambda) = f(x) + \lambda^{T}(\Phi x - u) + \frac{\mu}{2} \|\Phi x - u\|_{2}^{2}$$

The "AL method" (ALM) (a.k.a. method of multipliers) [Hestenes, Powell, 1969]

$$x_{k+1} = \arg\min_{x} L_{\mu}(x, \lambda_k)$$
$$\lambda_{k+1} = \lambda_k + \mu(\Phi x_{k+1} - u)$$

Can be written as:

$$x_{k+1} = \arg\min_{x} f(x) + \frac{\mu}{2} \|\Phi x - u - d_k\|_2^2$$

$$d_{k+1} = d_k - (\Phi x_{k+1} - u)$$

Similar to Bregman method [Osher, Burger, Goldfarb, Xu, Yin, 2005] [Yin, Osher, Goldfarb, Darbon, 2008]

Augmented Lagrangian for variable splitting

The problem:
$$\min_x f_1(\Phi x) + f_2(x)$$

Equivalent constrained formulation

$$\min_{x} \quad f_1(z) + f_2(x)$$

s.t.
$$\Phi x - z = 0$$

$$(x_{k+1}, z_{k+1}) = \arg\min_{x, z} f_1(z) + f_2(x) + \frac{\mu}{2} \|\Phi x - z - d_k\|_2^2$$
$$d_{k+1} = d_k - (\Phi x_{k+1} - z_{k+1})$$

Augmented Lagrangian for variable splitting

It may be hard to solve

$$(x_{k+1}, z_{k+1}) = \arg\min_{x,z} f_1(z) + f_2(x) + \frac{\mu}{2} ||\Phi x - z - d_k||_2^2$$
Alternative:

$$x_{k+1} = \arg\min_x f_2(x) + \frac{\mu}{2} ||\Phi x - z_k - d_k||_2^2$$

$$z_{k+1} = \arg\min_z f_1(z) + \frac{\mu}{2} ||\Phi x_{k+1} - z - d_k||_2^2$$

$$d_{k+1} = d_k - (\Phi x_{k+1} - z_{k+1})$$

Alternating directions method of multipliers (ADMM) [Glowinsky, Marrocco, 1975], [Gabay, Mercier, 1976], [Eckstein, Bertsekas, 1992]

When applied to
$$\hat{x} = \arg \min_{x} \frac{1}{2} \|\Phi x - u\|_{2}^{2} + \tau \|x\|_{1}$$

split augmented Lagrangian shrinkage algorithm (SALSA)
[F, Bioucas-Dias, Afonso, 2009]

Augmented Lagrangian for variable splitting

Testing ADMM/SALSA on a typical image deblurring problem





$$\widehat{x} \in \arg\min_{x \in \mathbb{R}^{n}} \frac{1}{2} || B\Psi x - u ||_{2}^{2} + \tau || x ||_{1}$$
Objective function
$$Objective function
Objective fun$$

Handling more than two functions

$$\widehat{x} \in \arg\min_{x\in\mathbb{R}^n} f_0(x) + f_1(x) + \dots + f_n(x)$$

 f_0 has a L-Lipschitz gradient $f_1,...,f_n$ are convex

Possible uses: multiple regularizers, positivity constraints, ...

Generalized forward-backward algorithm [Raguet, Fadili, Peyré, 2011] Parameters: $\omega_1, ..., \omega_n \in (0, 1), \text{ s.t. } \sum_j \omega_j = 1$ Initialization: $k = 0, \ z_0^1, ..., z_0^n, \ x_0 = \sum_{j=1}^n \omega_j \ z_0^j$ repeat until convergence for i = 1:n $\begin{bmatrix} z_{k+1}^i = z_k^i + \operatorname{prox}_{\beta_k f_i / \omega_i} \left(2x_k - z_k^i - \beta_k \nabla f_1(x_k) \right) - x_k \\ x_{k+1} = \sum_{i=1}^n \omega_i z_{k+1}^i \\ k \leftarrow k+1 \end{bmatrix}$

Handling more than two functions

$$\widehat{x} \in rg\min_{x\in\mathbb{R}^n} f_1(x) + \dots + f_n(x)$$

 $f_1, \dots, f_n \,\,$ arbitrary convex functions

ADMM-based method [F and Bioucas-Dias, 2009], [Setzer, Steidl, Teuber, 2009]

Parameter: γ Initialization: $k = 0, z_0^1, ..., z_0^n, y_0^1, ..., y_0^n$ repeat until convergence $x_{k+1} = (1/n) \sum_{i=1}^{n} (y_k^i - z_k^i)$ for i = 1 : n $y_{k+1}^i = \operatorname{prox}_{\gamma f_i} (x_k - z_k^i)$ $z_{k+1}^i = z_k^i + x_k - y_{k+1}^i$ $k \leftarrow k+1$

Non-Convex Algorithms for Low-Dimensional Models



Discrete descriptions of low-dimensional models $x = \sum_{i=1}^{|\mathcal{A}|} a_i c_i \qquad a_i \in \mathcal{A}, \|c_i\|_0 \leq K$ $x = \sum_{i=1}^{|\mathcal{A}|} a_i c_i \qquad a_i : \text{ atoms}$ $\mathcal{A}: \text{ atomic set}$ $\mathcal{A} = \{A: \operatorname{rank}(A) = 1, \|A\|_F = 1\}$

Example: reflectivity of Lambertian surfaces

[Basri and Jacobs 2001]



 $K \leq 9$

Intensity $= \rho \max\{\langle n, l \rangle, 0\}$

Discrete descriptions of low-dimensional models



Discrete descriptions of *structure* in low-dimensional models



Non-convex criteria beyond atomic norms

• 1-bit compressive sensing

$$u = \operatorname{sign}\left(\Phi x\right)$$

- optimization criteria $\arg\min_{x:\|x\|_0\leq K} f(x)$

$$f(x) = -\langle u, \operatorname{sign}(\Phi x) \rangle$$

• Compressible signals in weak $\ell_q, q < 1$

$$|x|_{(i)} \le Ri^{-1/q}$$

- optimization criteria $\arg \min_{x:u=\Phi x} \|x\|_q$



[Boufounos and Baraniuk 2008]



[Chartrand and Yin 2008]

Non-convexity in this tutorial

- Anything not convex <> too big to cover *convexity is in general a rare condition*
 Active research topic with great depth [Attouch et al. 2010] *Key lesson: convergence of the projected gradient-descent algorithm*
- This tutorial

<>

a special subset

 $\widehat{x} \in \arg \min_{x \in \mathbb{R}^N} f_1(x) + f_2(x)$ with $f_2(x) = \begin{cases} g(x) \iff x \in S \\ +\infty \iff x \notin S \end{cases}$

(${\mathcal S}$ is non-convex)



Assumptions:

- 1. access to the gradient of convex f_1
- 2. tractable/approximate prox of non-convex f_2

Running examples

- Sparse signal: only K out of N coordinates nonzero
 - model: union of all K-dimensional subspaces aligned w/ coordinate axes
- Structured sparse signal: reduced set of subspaces (or model-sparse)
 - model: a particular union of subspaces
 ex: clustered or dispersed sparse patterns





 \mathbf{R}^N

 $x \in$

Running examples

- Sparse signal: only K out of N coordinates nonzero
 - model: union of all K-dimensional subspaces aligned w/ coordinate axes
- Structured sparse signal: reduced set of subspaces (or model-sparse)

 \mathbf{R}^N

 $x \in \Sigma_K$

model: a particular union of subspaces
 ex: clustered or dispersed sparse patterns



$$\widehat{x} = \arg\min_{x \in \mathbb{R}^n} \frac{1}{2} \|y - x\|_2^2 + f_2(x) \equiv \operatorname{prox}_{f_2}(y)$$

• Analysis of the prox for *structured* sparse sets g(x) = 0

$$\operatorname{prox}_{f_2}(y) = \arg\min_{x:x\in\Sigma_{\mathcal{M}_K}} \|x-y\|$$

support of the solution <> modular approximation problem

$$\operatorname{supp}\left(\arg\min_{x:\operatorname{supp}(x)\in\mathcal{M}_{K}}\|x-y\|_{2}^{2}\right) = \arg\min_{\mathcal{S}:\mathcal{S}\in\bar{\mathcal{M}}_{K}}\|(y)_{\mathcal{S}}-y\|_{2}^{2}$$
indexing set

$$\widehat{x} = \arg\min_{x \in \mathbb{R}^n} \frac{1}{2} \|y - x\|_2^2 + f_2(x) \equiv \operatorname{prox}_{f_2}(y)$$

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 $= \operatorname{arg\,max}_{\mathcal{S}:\mathcal{S}\in\bar{\mathcal{M}}_{K}} \left\{ \|y\|^{2} - \|(y)_{\mathcal{S}} - y\|_{2}^{2} \right\}$

$$\widehat{x} = \arg\min_{x \in \mathbb{R}^n} \frac{1}{2} \|y - x\|_2^2 + f_2(x) \equiv \operatorname{prox}_{f_2}(y)$$

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 $\operatorname{supp}\left(\operatorname{arg\,min}_{x:\operatorname{supp}(x)\in\mathcal{M}_{K}}\|x-y\|_{2}^{2}\right) = \operatorname{arg\,min}_{\mathcal{S}:\mathcal{S}\in\bar{\mathcal{M}}_{K}}\|(y)_{\mathcal{S}}-y\|_{2}^{2}$

 $= \operatorname{arg\,max}_{\mathcal{S}:\mathcal{S}\in\bar{\mathcal{M}}_{K}} \left\{ \|y\|^{2} - \|(y)_{\mathcal{S}} - y\|_{2}^{2} \right\}$

$$= \arg \max_{\mathcal{S}: \mathcal{S} \in \bar{\mathcal{M}}_K} \| (y)_{\mathcal{S}} \|^2$$

$$\widehat{x} = \arg\min_{x \in \mathbb{R}^n} \frac{1}{2} \|y - x\|_2^2 + f_2(x) \equiv \operatorname{prox}_{f_2}(y)$$

• Analysis of the prox for *structured* sparse sets g(x) = 0

$$\operatorname{prox}_{f_2}(y) = \arg\min_{x:x\in\Sigma_{\mathcal{M}_K}} \|x-y\|$$

support of the solution <> modular approximation problem

 $\operatorname{supp}\left(\operatorname{arg\,min}_{x:\operatorname{supp}(x)\in\mathcal{M}_{K}}\|x-y\|_{2}^{2}\right) = \operatorname{arg\,max}_{\mathcal{S}:\mathcal{S}\in\bar{\mathcal{M}}_{K}}F(S;y)$

where $F(S; y) = \sum_{i \in S} |y_i|^2$.

$$\widehat{x} = \arg\min_{x \in \mathbb{R}^n} \frac{1}{2} \|y - x\|_2^2 + f_2(x) \equiv \operatorname{prox}_{f_2}(y)$$

• Analysis of the prox for *structured* sparse sets g(x) = 0

$$\operatorname{prox}_{f_2}(y) = \arg\min_{x:x\in\Sigma_{\mathcal{M}_K}} \|x-y\|$$

support of the solution <> modular approximation problem

$$\operatorname{supp}\left(\operatorname{arg\,min}_{x:\operatorname{supp}(x)\in\mathcal{M}_{K}}\|x-y\|_{2}^{2}\right) = \operatorname{arg\,max}_{\mathcal{S}:\mathcal{S}\in\bar{\mathcal{M}}_{K}}F(S;y)$$

underlying optimization problem <> integer linear program $supp \left(\arg\min_{z} \left\{ \rho^{T} z : z \in \Sigma_{\mathcal{M}_{K}} \right\} \right)$ $z_{i} \in \{0, 1\}: \text{ support indicator variables} \qquad \rho_{i} = -|y_{i}|^{2}$ [Kyrillidis and C, 2011]

$$\widehat{x} = \arg\min_{x \in \mathbb{R}^n} \frac{1}{2} \|y - x\|_2^2 + f_2(x) \equiv \operatorname{prox}_{f_2}(y)$$

• Analysis of the prox for *structured* sparse sets g(x) = 0

$$\operatorname{prox}_{f_2}(y) = \arg\min_{x:x\in\Sigma_{\mathcal{M}_K}} \|x-y\|$$

support of the solution <> modular approximation problem

$$\operatorname{supp}\left(\operatorname{arg\,min}_{x:\operatorname{supp}(x)\in\mathcal{M}_{K}}\|x-y\|_{2}^{2}\right) = \operatorname{arg\,max}_{\mathcal{S}:\mathcal{S}\in\bar{\mathcal{M}}_{K}}F(S;y)$$

underlying optimization problem <> integer linear program

Class of problems we can tractably solve: **PMAP**

• Polynomial time modular epsilon-approximation property $F(\widehat{S}_{\epsilon}; y) \ge (1 - \epsilon) \max_{S \in \overline{\mathcal{M}}_K} F(S; y)$ [Kyrillidis and C, 2011]

PMAP-0:

• Matroid structured sparse models:

$$\mathcal{M} = (\mathcal{N}, \mathcal{I} \subseteq 2^{\mathcal{N}}), \, \mathcal{N} = \{1, \dots, N\}$$

 \mathcal{N} : ground set \mathcal{I} : base set

Definition:

non-emptiness	1. $\emptyset \in \mathcal{I}$
heredity	2. $A \in \mathcal{I}$ and $B \subseteq A \Rightarrow B \in \mathcal{I}$
exchange	3. $A, B \in \mathcal{I}$ and $ A > B \Rightarrow \exists e \in A \setminus B$ such that $B \cup \{e\} \in \mathcal{I}$



PMAP-0:

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Definition:

non-emptiness1. $\emptyset \in \mathcal{I}$ heredity2. $A \in \mathcal{I}$ and $B \subseteq A \Rightarrow B \in \mathcal{I}$ exchange3. $A, B \in \mathcal{I}$ and $|A| > |B| \Rightarrow \exists e \in A \setminus B$ such that $B \cup \{e\} \in \mathcal{I}$

Let $\mathcal{N} = \{1, 2, 3, 4\}$. The smallest matroid that contains $\{1, 2\}$ and $\{3, 4\}$ is ???



[Nemhauser and Wolsey, 1999]

PMAP-0:

• Matroid structured sparse models:

$$\mathcal{M} = (\mathcal{N}, \mathcal{I} \subseteq 2^{\mathcal{N}}), \, \mathcal{N} = \{1, \dots, N\}$$

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Definition:

non-emptiness1. $\emptyset \in \mathcal{I}$ heredity2. $A \in \mathcal{I}$ and $B \subseteq A \Rightarrow B \in \mathcal{I}$ exchange3. $A, B \in \mathcal{I}$ and $|A| > |B| \Rightarrow \exists e \in A \setminus B$ such that $B \cup \{e\} \in \mathcal{I}$

Let $\mathcal{N} = \{1, 2, 3, 4\}$. The smallest matroid that contains $\{1, 2\}$ and $\{3, 4\}$ $\mathcal{I} = \{ \emptyset, \qquad by the non-emptiness property \\ \{1\}, \{2\}, \{3\}, \{4\}, \{1, 2\}, \{3, 4\}, by the heredity property \\ \{1, 3\}, \{1, 4\}, \{2, 3\}, \{2, 4\} \qquad by the exchange property \\ \}$



PMAP-0:

Matroid structured sparse models:

$$\mathcal{M} = (\mathcal{N}, \mathcal{I} \subseteq 2^{\mathcal{N}}), \, \mathcal{N} = \{1, \dots, N\}$$

Definition: 1.
$$\emptyset \in \mathcal{I}$$

2. $A \in \mathcal{I}$ and $B \subseteq A \Rightarrow B \in \mathcal{I}$
3. $A, B \in \mathcal{I}$ and $|A| > |B| \Rightarrow \exists e \in A \setminus B$ such that $B \cup \{e\} \in \mathcal{I}$

Greedy basis algorithm efficiently solves $\arg \max_{\mathcal{S}: \mathcal{S} \in \mathcal{M}} \sum_{i \in \mathcal{S}} w_i^2$

sort \mathcal{N} in decreasing order by weight w_i^2 start with empty set: $S_0 = \emptyset$ 1. $\mathcal{R}_i = \{r_i \in \mathcal{N} \setminus \mathcal{S}_i\}$ while keeping the order 2. $r = \arg \max_j \{ w_j^2 : (j \in \mathcal{R}_i) \land (\mathcal{S}_i \cup \{j\} \in \mathcal{I}) \}$ 3. $\mathcal{S}_{i+1} = \mathcal{S}_i \cup \{r\}$



PMAP-0:

• Matroid structured sparse models:

$$\mathcal{M} = (\mathcal{N}, \mathcal{I} \subseteq 2^{\mathcal{N}}), \, \mathcal{N} = \{1, \dots, N\}$$

Definition: 1.
$$\emptyset \in \mathcal{I}$$

2. $A \in \mathcal{I}$ and $B \subseteq A \Rightarrow B \in \mathcal{I}$
3. $A, B \in \mathcal{I}$ and $|A| > |B| \Rightarrow \exists e \in A \setminus B$ such that $B \cup \{e\} \in \mathcal{I}$

Greedy basis algorithm efficiently solves matroid constrained problems

Examples:

1. uniform matroid: $\mathcal{I} = \{\mathcal{S}: \mathcal{S} \subseteq \mathcal{N}, |\mathcal{S}| \leq K\}$ |y|

$$\operatorname{prox}_{f_2}(y) = \arg\min_{\substack{x:x \in \Sigma_K \\ \text{hard thresholding!} \\ H_K(y)}} \|x - y\|$$





PMAP-0:

• Matroid structured sparse models:

$$\mathcal{M} = (\mathcal{N}, \mathcal{I} \subseteq 2^{\mathcal{N}}), \, \mathcal{N} = \{1, \dots, N\}$$

Greedy basis algorithm efficiently solves matroid constrained problems

Ex	amples:		[Kyrillidis and C, 2011]		
1.	uniform matroid	<>	simple sparsity		
	intersection with the fo	ollowing matroid	ds (result is still a matroid!*)		
2.	partition matroid	<>	distributed sparsity		
3.	graphic matroid	<>	spanning tree sparsity		
4.	matching matroid	<>	graph matching sparsity		

*: in general, the intersection of two matroids is not a matroid.



PMAP-0:

• Linear support constraints:

Definition: $\Sigma_{\mathcal{M}_K} = \bigcup_{\forall z \in \mathfrak{Z}} \operatorname{supp}(z)$, where $\mathfrak{Z} := \{z \in \{0, 1\}^N : Az \leq b\}$

A and b<>integralfirst row of A<>all 1'sfirst entry of b<>K



PMAP-0:

• Linear support constraints:

Definition:
$$\Sigma_{\mathcal{M}_K} = \bigcup_{\forall z \in \mathfrak{Z}} \operatorname{supp}(z)$$
, where $\mathfrak{Z} := \{z \in \{0, 1\}^N : Az \leq b\}$

 \mathcal{U}

Example: neuronal spike model

 $z \in \{0, 1\}^N$: binary support variables

$$z_1 + z_2 + \ldots + z_N \le K$$

 $z_1 + z_2 + \ldots + z_\Delta \le 1$
 $z_2 + z_3 + \ldots + z_{\Delta+1} \le 1$

 $z_{N-\Delta+1} + z_{N-\Delta+2} + \ldots + z_N \le 1$



 $|\mathcal{M}_K|$

 \mathbf{R}^N

PMAP-0:

• Linear support constraints:

Definition:
$$\Sigma_{\mathcal{M}_{K}} = \bigcup_{\forall z \in \mathfrak{Z}} \operatorname{supp}(z)$$
, where $\mathfrak{Z} := \{z \in \{0, 1\}^{N} : Az \leq b\}$

We can use LP can relax the LS constrained ILPs:

$$\arg\min_{z} \left\{ \rho^{T} z : z \in [0, 1]^{N}, Az \le b \right\} \qquad \rho_{i} = -|y_{i}|^{2}$$

...but, when is the result binary?


PMAP-0:

• Linear support constraints:

Definition: $\Sigma_{\mathcal{M}_K} = \bigcup_{\forall z \in \mathfrak{Z}} \operatorname{supp}(z)$, where $\mathfrak{Z} := \{z \in \{0, 1\}^N : Az \leq b\}$

LP can exactly solve the LS constrained ILPs:

$$\arg\min_{z} \left\{ \rho^{T} z : z \in [0, 1]^{N}, Az \le b \right\} \qquad \rho_{i} = -|y_{i}|^{2}$$

 \mathbf{R}^N

 $|\mathcal{M}_K|$

...when A is totally unimodular (TU)*! [Nemhauser and Wolsey, 1999]

- the determinant of each square submatrix is {-1,0,1}

Examples: interval matrices, perfect matrices, network matrices

*: if we want LP relaxation to work for all b, TU is a necessary condition.

PMAP-0:

• Linear support constraints:

Definition:
$$\Sigma_{\mathcal{M}_{K}} = \bigcup_{\forall z \in \mathfrak{Z}} \operatorname{supp}(z)$$
, where $\mathfrak{Z} := \{z \in \{0,1\}^{N} : Az \leq b\}$
Example: neuronal spike model
 $z \in \{0,1\}^{N}$: binary support variables
 $\begin{array}{c}z_{1} + z_{2} + \ldots + z_{N} \leq K\\z_{1} + z_{2} + \ldots + z_{\Delta} \leq 1\\z_{2} + z_{3} + \ldots + z_{\Delta+1} \leq 1\\\vdots\\z_{N-\Delta+1} + z_{N-\Delta+2} + \ldots + z_{N} \leq 1\end{array}$
 \mathbf{TU}
 u
 \mathbf{U}
 \mathbf{U}

 \mathbf{R}^N

 $x \in \Sigma_{\mathcal{M}_K}$

 $|\mathcal{M}_K|$

[Hegde, Duarte, and C, 2009]

PMAP-0:

• prox-sparse models

Definition: define algorithmically!

$$\operatorname{prox}_{f_2}(y) = \arg\min_{x:x\in\Sigma_{\mathcal{M}_K}} \|x-y\| \qquad g(x) = 0$$



PMAP-0:

prox-sparse models

Definition: define algorithmically!

$$\operatorname{prox}_{f_2}(y) = \arg\min_{x:x\in\Sigma_{\mathcal{M}_K}} \|x-y\| \qquad \qquad g(x) = 0$$

Example: clustered sparsity models

- tree-sparse <> dynamic program





clustered sparse <> dynamic program

[Baraniuk, C, Wakin 2010; Baraniuk, C, Duarte, Hegde 2010]



Pop-quiz: A prox with convex and non-convex terms

Let us consider $f_2(x) = \|x\|_1 + \iota_{\{x:\|x\|_0 \le K\}}(x)$ $g(x) = \|x\|_1$

$$\operatorname{prox}_{f_2}(y) = \arg\min_{x \in \mathbb{R}^n} \frac{1}{2} \|y - x\|_2^2 + f_2(x)$$

Is it PMAP-0?

Pop-answer: A prox with convex and non-convex terms

Let us consider $f_2(x) = \|x\|_1 + \iota_{\{x:\|x\|_0 \le K\}}(x)$ $g(x) = \|x\|_1$

$$\operatorname{prox}_{f_2}(y) = \arg\min_{x \in \mathbb{R}^n} \frac{1}{2} \|y - x\|_2^2 + f_2(x)$$

$$\operatorname{supp}\left(\operatorname{prox}_{f_2}(y)\right) = \arg \max_{\mathcal{S}:|\mathcal{S}| \le K} F(\mathcal{S};y)$$

 $F(\mathcal{S}; y) = \frac{1}{2} \|y\|^2 - \min_{x: \text{supp}(x) = \mathcal{S}} \frac{1}{2} \|y - x\|_2^2 + \|x\|_1$

$$\Rightarrow F(\mathcal{S}; y) = \sum_{i \in \mathcal{S}} w_i^2$$

$$w_i^2 = y_i \times \text{soft}(y_i, 1) - \frac{1}{2} |\text{soft}(y_i, 1)|^2 - |\text{soft}(y_i, 1)|$$

Hard thresholding followed by soft thresholding!

YES: certified PMAP-0





PMAP-epsilon: $F(\widehat{\mathcal{S}}_{\epsilon}; y) \ge (1 - \epsilon) \max_{\mathcal{S} \in \overline{\mathcal{M}}_{K}} F(\mathcal{S}; y)$

Knapsack

multi-knapsack constraints weighted multi-knapsack

Ex: Nested group sparse problems

quadratically-constrained

Define algorithmically!

approximate solutions for computational reasons





[Kyrillidis and C, 2011]

PMAP-epsilon: $F(\widehat{\mathcal{S}}_{\epsilon}; y) \ge (1 - \epsilon) \max_{\mathcal{S} \in \overline{\mathcal{M}}_{K}} F(\mathcal{S}; y)$

Knapsack

multi-knapsack constraints weighted multi-knapsack

Ex: Nested group sparse problems

quadratically-constrained

• Define algorithmically!

approximate solutions for computational reasons



• **Pairwise overlapping groups** <> quadratic binary w/ cardinality cons.

 $\max_{\mathcal{S}:\mathcal{S}\in\bar{\mathcal{M}}_{K}} F(S;y) = -\min\left\{\sum_{i>j} \|(y)_{g_{i}\cap g_{j}}\|_{2}^{2} z_{i} z_{j} - \sum_{i} \|(y)_{g_{i}}\|_{2}^{2} z_{i} : \sum_{i} z_{i} \leq G\right\}.$

we can only approximate... and epsilon is large!



PMAP-epsilon: $F(\widehat{\mathcal{S}}_{\epsilon}; y) \ge (1 - \epsilon) \max_{\mathcal{S} \in \overline{\mathcal{M}}_{K}} F(\mathcal{S}; y)$

Knapsack

multi-knapsack constraints weighted multi-knapsack

Ex: Nested group sparse problems

quadratically-constrained

• Define algorithmically!

approximate solutions for computational reasons



- Pairwise overlapping groups <> quadratic binary w/ cardinality cons.
- Multi-knapsack + multi-matroids

[Lee et al., 2009]

we can only approximate... and epsilon is large!



Matrix examples!



• Rank constrained projections $\operatorname{prox}_{f_2}(Y) = \arg \min_{X:\operatorname{rank}(X) \leq R} \|X - Y\|_F$

Matrix examples!



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$$\arg\min_{X:\operatorname{rank}(X)\leq R} \|X-Y\|_F = \arg\min_{X:\operatorname{rank}(X)\leq R} \|X-U\Lambda_Y V^T\|_F \text{ singular value decomposition}$$



Matrix examples!



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$$\arg\min_{X:\operatorname{rank}(X)\leq R} \|X-Y\|_F = \arg\min_{X:\operatorname{rank}(X)\leq R} \|X-U\Lambda_Y V^T\|_F$$

$$= \arg \min_{X: \operatorname{rank}(X) \leq R} \| U^T X V - \Lambda_Y \|_F \text{ invariance to}$$
unitary transform



Matrix examples!



• Rank constrained projections $\operatorname{prox}_{f_2}(Y) = \arg \min_{X:\operatorname{rank}(X) \leq R} \|X - Y\|_F$

$$\arg\min_{X:\operatorname{rank}(X)\leq R} \|X-Y\|_F = \arg\min_{X:\operatorname{rank}(X)\leq R} \|X-U\Lambda_Y V^T\|_F$$

$$= U \left(\arg \min_{\tilde{X}: \operatorname{rank}(\tilde{X}) \le R} \| \tilde{X} - \Lambda_Y \|_F \right) V^T$$

sparse approximation problem!



Matrix examples!



• Rank constrained projections $\operatorname{prox}_{f_2}(Y) = \arg \min_{X:\operatorname{rank}(X) \leq R} \|X - Y\|_F$

$$\arg\min_{X:\operatorname{rank}(X)\leq R} \|X-Y\|_F = \arg\min_{X:\operatorname{rank}(X)\leq R} \|X-U\Lambda_Y V^T\|_F$$

$$= UH_R(\Lambda_Y)V^T$$

singular value (hard) thresholding



Matrix examples!



• Rank constrained projections $\operatorname{prox}_{f_2}(Y) = \arg \min_{X:\operatorname{rank}(X) \leq R} \|X - Y\|_F$

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 $= UH_R(\Lambda_Y)V^T$

singular value (hard) thresholding

- Non-convex spectral projections <> sets described by their eigenvalue properties
 - exact projections
 basic operations on eigenvalues

[Lewis and Malick 2008]

Matrix examples!

- Rank constrained projections $\operatorname{prox}_{f_2}(Y) = \arg \min_{X:\operatorname{rank}(X) \leq R} \|X Y\|_F$
- Non-convex spectral projections <> sets described by their eigenvalue properties
- epsilon-approximate projections (note the difference with PMAP)

 $\|\operatorname{prox}_{f_2}^{\epsilon}(Y) - Y\|_F \le (1+\epsilon) \min_{X:\operatorname{rank}(X) \le R} \|X - Y\|_F$

Two highlights:

- structure from randomness/power methods [Halko,]
- column subset selection approaches

[Halko, Martinsson, Tropp, 2010]

[Boutsidis, Mahoney, Drineas, 2010]

Recovery algorithms for low-dimensional models

Now that we have projections...



A common criteria covering a broad set of applications:

$$\min_X ||u - \Phi(X)||^2$$
 s.t. $X = S + L, ||S||_0 \le K, \operatorname{rank}(L) \le R$

- affine rank minimization, matrix completion, robust PCA...

[Candes and Recht 2009; Waters, Sankaranayanan, Baraniuk, 2011] A common algorithm: projected gradient

 $||S||_0 = \#\{S_i \neq 0\}$

Recovery algorithms for low-dimensional models

To highlight the salient differences, we will consider

	Non-convex $\binom{N}{K}$	Convex	Probabilistic
Encoding	combinatorial /	atomic norm /	compressible /
	manifolds	convex relaxation	sparse priors

compressive sensing recovery

$$\min_{x:\|x\|_0 \le K} \|u - \Phi x\|^2$$

A common algorithm:

projected gradient

 $||x||_0 = \#\{x_i \neq 0\}$

• Soft thresholding

 $\min_{x:\|x\|_1 \le \lambda} f(x)$







• Soft thresholding

 $\min_{x:\|x\|_1 \le \lambda} f(x)$







Structure in optimization:

(1)
$$f(y) - f(x) - \langle \nabla f(x), y - x \rangle = \|\Phi(y - x)\|^2 \quad \forall x, y \in \mathbb{R}^N,$$



(2)
$$f(y) - f(x) - \langle \nabla f(x), y - x \rangle \leq \frac{L}{2} \|y - x\|^2 \quad L = 2 \|\Phi\|^2, \forall x, y \in \mathbb{R}^N,$$





• Soft thresholding $\min_{x:||x||_1 \le \lambda} f(x)$

 Is x* what we are looking for?

local "unverifiable" assumptions:

- ERC/URC/RSC condition
- coherence based conditions ...

(local \rightarrow global / dual certification)

$$f(y) = \|u - \|u\|$$

$$(2)$$

$$(1)$$

$$(1)$$

$$x \in \{\|x\|_{1} \le \lambda\} \xrightarrow{x_{2}} x_{1}$$

$$x_{i+1} = \operatorname{St}_{\{\|x\|_{1} \le \lambda\}} (x_{i} - \frac{1}{L} \nabla f(x_{i}))$$

 $f(r) = \| u - \Phi r \|^2$

[Buhlmann and van de Geer 2011]









What could possibly go wrong with this naïve approach?

[C, 2011]







Restricted Isometry Property

• **Model:** *K*-sparse coefficients

Remark: implies convergence of convex relaxations also e.g., $\delta_{2K} < .465$ is sufficient for BP

• **RIP:** stable embedding





Restricted Isometry Property

• Model: *K*-sparse coefficients

Remark: implies convergence of convex relaxations also e.g., $\delta_{2K} < .465$ is sufficient for BP

RIP: stable embedding





Restricted Isometry Property for Matrices!

• Model: rank-R matrices

Remark: bi-Lipschitz embedding of low-rank matrices

• **RIP:** stable embedding



[Plan 2011]



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 x_{2} x_{1} x_{2} x_{2} x_{2} x_{2} x_{2} x_{2} x_{2} x_{2} x_{3} x_{4} x_{2} x_{2} x_{2} x_{3} x_{4} x_{2} x_{2} x_{3} x_{4} x_{4} x_{2} x_{4} x_{2} x_{3} x_{4} x_{4} x_{2} x_{4} x_{5} x_{5} x_{5} x_{6} x_{7} x_{7

Projected gradient method for non-convex sets



Model: *K*-sparse coefficients + significant coefficients lie on a rooted subtree



Sparse approx:

find best set of coefficients

- sorting
- hard thresholding

Tree-sparse approx: find best rooted subtree of coefficients

- condensing sort and select
- dynamic programming

[Baraniuk]

[Donoho]

[Baraniuk, C, Duarte, Hegde 2010]

• **Model:** *K*-sparse coefficients



RIP: stable embedding



- Model: K-sparse coefficients
 + significant coefficients lie on a rooted subtree
- Tree-RIP: stable embedding





- Number samples for correct recovery
- Piecewise cubic signals + wavelets
- Models/algorithms:
 - compressible (CoSaMP)
 - tree-compressible (tree-CoSaMP)



[Baraniuk, C, Duarte, Hegde 2010]

Recovery algorithms for low-dimensional models

The Clash Operator

Encoding combinatorial / atomic norm	m / compressible /
manifolds convex relaxa	ation sparse priors
Example $\min_{x: x _0 \le K} u - \Phi x ^2$ $\min_{x: x _1 \le \lambda} u - \Phi x ^2$	$\Phi x ^2 \qquad E\{x u\}$
Algorithm IHT, CoSaMP, SP, ALPS,	so, Variational Bayes, EP,
OMP Basis pursuit, Lass	ising Approximate message
basis pursuit denoi	passing (AMP)

 $\widehat{x}_{\text{Clash}} = \arg\min_{x:\|x\|_0 \le K, \|x\|_1 \le \lambda} \|u - \Phi x\|^2$

$$||x||_0 = \#\{x_i \neq 0\}$$

[Kyrillidis and C, 2011]

Recovery algorithms for low-dimensional models

 $\widehat{x} = \arg \min \|x\|_0 \text{ s.t. } u = \Phi x$



 $\widehat{x} = \arg \min \|x\|_1 \text{ s.t. } u = \Phi x$ $\{x' : u = \Phi x'\}$ $\|x\|_1 = c$
• A subtle issue



Which one is correct?

• A subtle issue





"Greed is good." – Joel Tropp 2004

• A subtle issue



Which one is correct?



The CLASH algorithm



$$H_{\{\|x\|_0 \le K\}}(t) = \arg \min_{\|x\|_0 \le K} \|x - t\|$$

$$St_{\{\|x\|_1 \le \lambda\}}(t) = \arg \min_{\|x\|_1 \le \lambda} \|x - t\|$$



The CLASH algorithm



The Clash Operator

	Non-convex $\binom{N}{K}$	Convex	Probabilistic
Encoding	combinatorial / manifolds	atomic norm / convex relaxation	compressible / sparse priors
Example	$\min_{x:\ x\ _0 \le K} \ u - \Phi x\ ^2$	$\min_{x:\ x\ _1 \le \lambda} \ u - \Phi x\ ^2$	$E\{x u\}$
Algorithm	IHT, CoSaMP, SP, ALPS, OMP	Basis pursuit, Lasso, basis pursuit denoising	Variational Bayes, EP, Approximate message passing (AMP)

 $\widehat{x}_{\text{Clash}} = \arg\min_{x:\|x\|_0 \le K, \|x\|_1 \le \lambda} \|u - \Phi x\|^2$

The idea is much more general

$$\widehat{x}_{\text{Normed Pursuit}} = \arg\min_{x:\|x\|_0 \le K, \|x\|_* \le \lambda} \|u - \Phi x\|^2$$

 $||x||_0 = \#\{x_i \neq 0\}$

[Kyrillidis, Puy, and C, 2012]

- Using projected gradient with exact non-convex projections with RIP/ERC/URC/RSC...
- Exact low-dimensional model
 - noise-free measurements: exact recovery
 - noisy measurements: stable recovery
- Approximately low-dimensional model
 - recovery as good as K-model-sparse approximation

$$\|x - \hat{x}\|_{\ell_2} \le C_1 \log\left(\frac{N}{K}\right) \frac{\|x - x_{\mathcal{M}_K}\|_{\ell_1}}{K^{1/2}} + C_2 \epsilon$$
recovery
signal K-term noise

recovery error signal *K*-term model approx error

[Baraniuk, C, Duarte, Hegde 2010]

- Using projected gradient with exact non-convex projections with RIP/ERC/URC/RSC...
- Exact low-dimensional model
 - noise-free measurements:
 exact recovery
 - noisy measurements: stable recovery
- Approximately low-dimensional model
 - recovery as good as K-model-sparse approximation

$$\begin{aligned} \|x - \hat{x}\|_{\ell_2} &\leq C_1 \log \left(\frac{N}{K}\right) \frac{\|x - x_{\mathcal{M}_K}\|_{\ell_1}}{K^{1/2}} + C_2 \epsilon \\ \\ \xrightarrow{\text{recovery}}_{\text{error}} & \text{signal } \textit{K-term}_{\text{model approx error}} & \text{noise} \end{aligned}$$

- the bound remains qualitatively the same for other models!!!

 Projected gradient with (non)exact non-convex projections without RIP/ERC/URC/RSC...

• Not much!

- convergence to stationary point with *Kurdyka-Lojasiewicz*

[Attouch et al., 2010]



Coded Aperture Snapshot Spectral Imager





http://www.disp.duke.edu/projects/CASSI/





Acceleration of non-convex algorithms

• Several approaches

step-size selection

$$x_{i+1} = H_{\Sigma_{\mathcal{M}_K}} (y_i - \mu_i \nabla f(y_i))$$

$$y_{i+1} = x_{i+1} + \tau_i (x_{i+1} - x_i)$$

memory based methods similar to Nesterov acceleration / double overrelaxation

non-convex splitting

(adaptive) block coordinate descent

epsilon-approximate projections



Acceleration of non-convex algorithms

• Several approaches

step-size selection

memory based methods similar to Nesterov acceleration / double overrelaxation 34.8s Original



Low rank

144 x 176 x 200









non-convex splitting

(adaptive) block coordinate descent

epsilon-approximate projections

[Zhou and Tao 2011; Kyrillidis and C, 2012]







MATRIX ALPS





Final remarks

- non-convex algorithms
 - possible performance gains
 - non-convexifiable priors
 - matching prox operator with optimal space/time bounds

complexity of structured approximation

- non-convex algorithms
 vs. convex algorithms
 - no clear winner / scenario dependent
 - decades of research in both

low-dimensional scaffold





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