

# Discrete Optimization 2024 (EPFL): Problem set of week 8

May 2, 2024

Reminder: The min-max theorem for zero-sum games with mixed strategies says that for every  $m \times n$  matrix  $A$  we have

$$\min_y \max_x yAx = \max_x \min_y yAx,$$

where the minimum is over all  $y = (y_1, \dots, y_m) \geq 0$  such that  $\sum y_i = 1$ . The maximum is over all  $x = (x_1, \dots, x_n) \geq 0$  such that  $\sum x_i = 1$ .

1. Let  $A$  be an  $m \times n$  matrix. Assume that there is an entry in  $A$  that is the minimum in its column and the maximum in its row. Prove that this entry is the value of the zero-sum game with for two players with mixed strategies.

Solution: Let the value of that entry be equal to  $M$  and assume it is in the  $i$ 'th row and  $j$ 'th columns.

By having player  $R$  choose the  $i$ 'th row with probability 1, then no matter what player  $C$  chooses, the value of the game will be at most  $M$  because  $M$  is the largest in its row. This shows that  $\min_y \max_x yAx \leq M$ .

By having player  $C$  choose the  $j$ 'th column with probability 1, then no matter what player  $R$  chooses, the value of the game will be at least  $M$  because  $M$  is the smallest in its columns. This shows that  $\max_x \min_y yAx \geq M$ .

We are done because we know that

$$\min_y \max_x yAx = \max_x \min_y yAx.$$

2. Find  $\max_x \min_y yAx$  for the matrix

$$A = \begin{pmatrix} 5 & 1 \\ 3 & 4 \end{pmatrix}$$

What is the (mixed) strategy for the column player to guarantee the maximum possible result?

Solution: We follow the example we did in class. Denoting  $x = (a, 1-a)$  and  $y = (b, 1-b)$  we find that  $yAx = (5b-1)(a-\frac{3}{5}) + 3\frac{2}{5}$ . From here it is not hard to see that the value of the game is  $3\frac{2}{5}$  and the best strategy for the column player is to use the vector  $(\frac{3}{5}, \frac{2}{5})$  in order to guarantee this optimal value.

3. Find the min-max value for the diagonal matrix with  $\lambda_1, \dots, \lambda_n$  on the main diagonal.

Solution: If  $\lambda_i \geq 0$  and  $\lambda_j \leq 0$  (possibly  $i = j$ ), then the min-max value is equal to 0. This is because we can apply Problem 1 on the entry  $a_{ij} = 0$  that is the largest in its column and smallest in its row.

Therefore, assume  $\lambda_1, \dots, \lambda_n > 0$ .

Then  $yAx = \sum x_i y_i \lambda_i$ . Fixing  $y$  we have  $\max_x \sum x_i y_i \lambda_i$  is when  $x_i = 1$  for the  $i$  such that  $y_i \lambda_i$  is maximum.

Then  $\max_x yAx = \max_i y_i \lambda_i$ . Therefore, if we want to find  $\min_y \max_x yAx$  we better have  $y$  such that  $\max_i y_i \lambda_i$  is minimum. Similar to what we did on Problem 2, this happens when all the  $y_i \lambda_i$  are equal. Then  $y_i = \frac{1}{\lambda_i}$  times a constant that does not depend on  $i$ . Because  $\sum y_i = 1$ , then we must have

$$y_i = \frac{1}{\sum \frac{1}{\lambda_j}} \frac{1}{\lambda_i}.$$

Then the value of the min-max is  $\frac{1}{\sum \frac{1}{\lambda_j}}$ .

What is the  $\lambda_i$ 's are all negative?

In this case notice that we have (we use the fact that  $\min f = -\max(-f)$ )

$$\begin{aligned}
\frac{1}{\sum \frac{1}{\lambda_j}} &= -\min_y \max_x y(-A)x \\
&= -\min_y \max_x -yAx \\
&= -\min_y (-\min_x yAx) \\
&= \max_y \min_x yAx \\
&= \max_y \min_x xAy \\
&= \min_x \max_y xAy \\
&= \min_y \max_x yAx
\end{aligned}$$

4. Show that in a zero-sum game with a matrix  $A$  with mixed strategies the following is true: If one player knows the mixed strategy of the other player, then the best response (strategy) for him is a pure strategy. That is, the best response is choosing just one row or column.

Solution: Assume for example that the Column player knows the strategy  $y_0$  of the Row player. Then the Column player wants to maximize  $\max_x y_0 Ax$  subject to  $\sum x_i = 1$  and  $x \geq 0$ . However, this is a linear program. The simplex algorithm will find a vertex of the polyhedron  $\sum x_i = 1$  and  $x \geq 0$ . The vertices of this polyhedron are the  $n$  vectors that have all their coordinates equal to 0 except one coordinate that is equal to 1. This gives a pure strategy.

5. Find a hyperplane separating the point  $v = (4, 3, 6)$  from the ball  $\{(x, y, z) \mid (x - 1)^2 + (y - 2)^2 + (z - 3)^2 \leq 16\}$ .

Solution: The radius of the ball is  $\sqrt{16} = 4$ . The center of the ball is  $u = (1, 2, 3)$ . The closest point of the ball to  $v$  is  $w = u + 4(v - u)/|v - u|$ . Let  $z$  be the midpoint of the segment  $[v, w]$ , that is  $(v + w)/2$ . It is enough to take the hyperplane perpendicular to  $v - u$  through  $z$ .