## Discrete Optimization (Spring 2018)

# Assignment 11

**Problem 4** could be **submitted** until May 25 12:00 noon into the right box in front of MA C1 563. You are allowed to submit your solutions in groups of at most three students.

### Problem 1

Let  $M \in \mathbb{Z}^{n \times m}$  be totally unimodular. Prove that the following matrices are totally unimodular as well:

- 1.  $M^T$
- 2.  $(M I_n)$
- 3. (M M)
- 4.  $M \cdot (I_n 2e_i^T e_j)$  for some j.

 $I_n$  is the  $n \times n$  identity matrix and  $e_j$  is the vector having a 1 in the j-th component, and 0 in the other components.

#### Problem 2

Let G be a graph and let A be its node-edge incidence matrix. We have seen in class that if G is bipartite then A is totally unimodular. Prove the converse, *i.e.*, if A is totally unimodular then G is bipartite.

#### Problem 3

Consider a bipartite graph  $G = (A \cup B, E)$ . Assume there exist matchings  $M_A$  and  $M_B$  covering vertices  $A_1 \subseteq A$  and  $B_1 \subseteq B$ , respectively. Prove that there always exists a matching that covers  $A_1 \cup B_1$ .

Hint: The symmetric difference  $M_A \Delta M_B$  consists of only cycles and paths.

## Problem 4 $(\star)$

Given a graph G(V, E), a perfect matching of G is a matching which covers all the vertices (equivalently, a matching of cardinality |V|/2). Suppose you are given an oracle that, given a graph G, tells you whether G has a perfect matching or not. Show how to use this oracle to find a maximum cardinality matching of a graph G(V, E), using at most |V| + |E| calls to the oracle.

Hint: you should modify the graph at each call of the oracle.

#### Problem 5

A family of sets  $C \subset 2^{[n]}$  is a chain if for all  $S, T \in C$  we have either  $S \subseteq T$  or  $T \subseteq S$ . Suppose  $C_1$  and  $C_2$  are two chains. Let  $A \in \{0,1\}^{(|C_1|+|C_2|)\times n}$  be the incidence matrix of  $C_1 \cup C_2$ , i.e.  $A_{S,i} = 1$  if  $i \in S$  and 0 otherwise, for  $i = 1, \ldots, n$  and  $S \in C_1 \cup C_2$ . Prove that A is totally unimodular. Hint: use induction on the size of a square submatrix of A.