

Two person general Sum Games

Two players: Row-Player and Column-Player

Two payoff matrices: $A \in \mathbb{R}^{m \times n}$, $B \in \mathbb{R}^{n \times m}$

x, y mixed strategies for ~~columns~~ row and column player respectively, then

$x^T \cdot A \cdot y =$ Expected payoff for row-player.

$x^T \cdot B \cdot y =$ Expected payoff for column player.

Example: Prisoners dilemma:

	Confess	Silent
Confess	4, 4	1, 5
Silent	5, 1	2, 2

(numbers are 'negative' of payoff)

Only stable pure strategy is that both confess

Nash equilibrium: Pair x, y of mixed strategies s.t.

\forall mixed strategies x' : $x'^T \cdot A \cdot y \leq x^T \cdot A \cdot y$

\forall mixed strategies y' : $x^T \cdot B \cdot y' \leq x^T \cdot B \cdot y$

Lemma: Best response to column players (mixed)

strategy y^* is pure strategy.

Best response to row players (mixed)

strategy x is pure strategy.

proof: Best response computation via LP.

Given y , compute

$$\max x^T \cdot A \cdot y$$

$$\sum_{i=1}^m x_i = 1$$

$$x_i \geq 0 \quad i=1, \dots, m.$$

Basic solutions to LP are exactly the unit vectors

Similar argument for best response of column-player. □

Lemma: Given mixed strategies x and y , there are always

pure strategies x' and y' with

$$x'^T \cdot A \cdot y \leq x^T \cdot A \cdot y$$

$$x^T \cdot B \cdot y' \leq x^T \cdot B \cdot y$$

proof: x and y are convex combinations of pure strategies

(unit vectors)



Theorem (Nash): Every Two-person general sum Game (bi-matrix Game) has Nash equilibrium.

For the proof we need:

Theorem: (Brouwer's Fixpoint Theorem)

$S \subseteq \mathbb{R}^n$ compact convex and $f: S \rightarrow S$ continuous, then there exists at least one $x \in S$ with $f(x) = x$.

Proof of Nash Theorem:

Let x, y be mixed strategies.

$$r_i(x, y) = \max \{ 0, e_i^T A \cdot y - x^T A \cdot y \}$$

$$c_j(x, y) = \max \{ 0, x^T B \cdot e_j - x^T B \cdot y \}$$

Notice: x, y Nash equilib. $\Leftrightarrow r = c = 0$

New strategies:
$$x_i' = \frac{x_i + r_i}{1 + \sum_k r_k}$$

$$y_j' = \frac{y_j + c_j}{1 + \sum_k c_k}$$

Interpretation: "Flow" x and y to improving direction, if not all r_i, c_j are zero

$T(x, y) = (x', y')$ is continuous.

$S = \{ (x, y) \in \mathbb{R}^{m+n} : x, y \text{ mixed strategies} \}$ is compact, convex.

Brouwer's Fix point Thm $\Rightarrow \exists (x, y)$ with

$$T(x, y) = (x, y)$$

Claim: (x, y) with $T(x, y) = (x, y)$ is Nash equilibrium.

proof of claim:

Suppose (x, y) is not Nash equilibrium. Then:

$$\exists i \text{ with } e_i^T \cdot A \cdot y > x^T \cdot A \cdot y \quad \text{or}$$

$$\exists j \text{ with } x^T \cdot B \cdot e_j > x^T \cdot B \cdot y.$$

Assume first cond. holds. Then $r_i > 0$ and

~~Lemma~~ Thus $\sum_{i=1}^m r_i > 0.$

Lemma $\Rightarrow \exists i'$ with $r_{i'} = 0$

$$\Rightarrow \exists i' \text{ with } x_{i'}' \neq x_{i'}.$$

implying that (x, y) is not a fixpoint.

If second condition holds, then similar argument applies. \Rightarrow claim

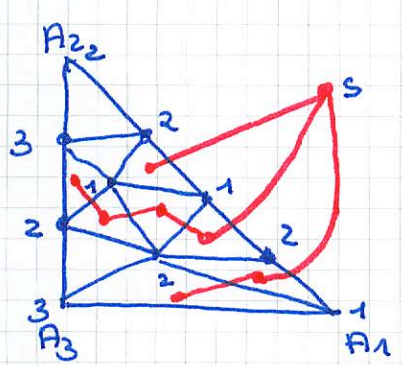
$\Rightarrow \exists$ Nash equilibrium



Why does Brouwer's Fixpoint Thm. Hold?

We provide elementary Proof for the case, where $S \subseteq \mathbb{R}^2$ is a triangle. (unit triangle)

Sperner Lemmma:



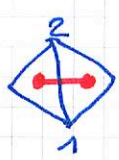
- A triangle with ^{vertices A_1, A_2, A_3} ~~vertex labels~~
- points in triangle have labels in $\{1, 2, 3\}$
- Each point on $\overline{A_i A_j}$ has labels i, j
- Labels of A_i are i respectively.

Sperner's Lemmma: Each Triangulation of

points has rainbow triangle (labels 1, 2, and 3)

proof: Consider Graph induced by Triangulation.

Draw only edges of Dual Graph that correspond to neighboring triangles with common 1,2 border

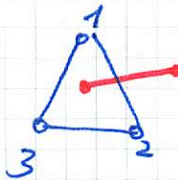


Recall that for graph $G=(V,E)$:

$$\sum_{v \in V} \deg(v) = 2E.$$

Degree of 3 is odd, since odd number of 1-2 swaps!

$\Rightarrow \exists$ other odd triangle



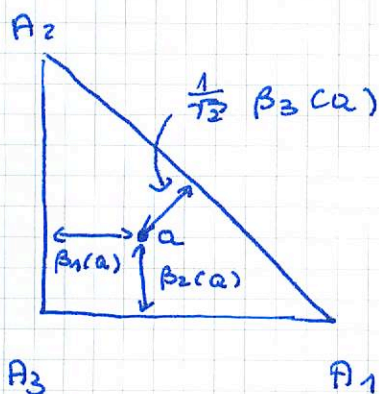
Rainbow Triangle!

□

Brouwer's Fixpoint Thm for unit Triangle

$f: \Delta \rightarrow \Delta$ continuous, then f has a fix point

proof:



$$a = (x, y)$$

$$\beta_1(a) = x$$

$$\beta_2(a) = y$$

$$\beta_3(a) = 1 - x - y$$

We have:

$$1) \beta_i(a) \geq 0 \quad \forall a \in \Delta$$

$$2.) \beta_1(a) + \beta_2(a) + \beta_3(a) = 1 \quad \forall a \in \Delta$$

Label points in triangle with $\{1, 2, 3\}$

$$\Pi_i := \{a \in \Delta : \beta_i(a) \geq \beta_i(f(a))\} \quad , \Pi_1 \cup \Pi_2 \cup \Pi_3 = \Delta$$

We want:

$$\text{label}(a) = i \quad \Rightarrow \quad a \in \Pi_i \quad (*)$$

~~most important?~~

Notice $A_i \in \Pi_i \Rightarrow \text{label}(A_i) = i$ is admissible.

Notice that line segment $\overline{A_i A_j} \in \Pi_i \cup \Pi_j \quad i, j \in \{1, 2, 3\}$

\Rightarrow points on $\overline{A_i A_j}$ have label in $\{i, j\}$

\Rightarrow Labeling (*) as in prerequisites of Sperner

Lemma is possible.

Subdivide Triangle into smaller and smaller
Triangles



Each subdivision contains rainbow triangle (Sperner)

Sequence of Rainbow triangles has converging subsequence

(Bolzano-Weierstrass). Common 3-vertices of limit
triangle is point $p \in \Delta$. ~~is~~ (Δ is closed & bounded).

We have $\beta_i(p) \geq \beta_i(f(p)) \quad i=1,2,3$

$\Rightarrow p$ is fixed point of f . $f(p) = p$

□