## Introduction to Combinatorics

# Homework 11 - Linear Algebra Method 

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## Questions

1. Prove the "reverse oddtown theorem": given a set $X$ of $n$ elements, and $m$ sets $S_{1}, \ldots, S_{m}$ over $X$ where each $\left|S_{i}\right|$ is even, and $\left|S_{i} \cap S_{j}\right|$ is odd, prove that $m \leq n+1$. Can you improve this to $m \leq n$ ?
Solution: Add a dummy element to each set, and apply Oddtown theorem. Therefore, $m \leq n+1$. It can also be shown that $m \leq n$. If $n$ is odd, take the complement sets, and apply Oddtown theorem. Otherwise, $n$ is even. Now note that the set system made up of $\left\{S_{i}, i=1 \ldots m\right\}$, and the complement set system made up of sets $\left\{\overline{S_{i}}, i=1 \ldots m\right\}$ satisfying same condition. And for contradiction, assume $m=n+1$. Then by making the equation for linear dependency on the characteristic vectors of both set $S_{i}$, and the complement set $\overline{S_{i}}$ (following the proof done in class), it can be deduced that $v_{i}+\overline{v_{i}}+\ldots+v_{n+1}+\overline{v_{n+1}}=0$, a contradiction as $n+1$ is odd, and so the above sum is 1 , and not 0 .
2. Prove the "bipartite oddtown theorem": given a set $X$ of $n$ elements, and sets $R_{1}, \ldots, R_{m}$ and $B_{1}, \ldots, B_{m}$, where $\left|R_{i} \cap B_{i}\right|$ is odd for every $i$, and $\left|R_{i} \cap B_{j}\right|$ is even for every $i \neq j$, prove that $m \leq n$.
Solution: Say characteristic vector of $R_{i}$ is $r_{i}$, and $B_{i}$ is $b_{i}$. Then make the linear-dependency equation for $r_{i}$. To see $\lambda_{i}=0$, multiply equation by $b_{i}$.
3. A block-design consisting of $n$ total elements $X$ and $m$ sets over $X$ where each set has $k$ elements and every $t$-sized subset of $X$ is contained in exactly $\lambda$ sets is denoted by $t-(n, k, \lambda)$. Then Fisher's inequality states that $m \geq n$.
Prove that for a block-design with $t=2, \lambda \cdot \frac{n-1}{k-1}$ has to be an integer.
Solution: Say the first element lies in $c$ sets. Then for each of those sets, it makes a pair with the other $k-1$ elements, for a total of $c(k-1)$ pairs. On the other hand, it makes a pair with $n-1$ elements, and each of those pairs appears in $\lambda$ sets. So then $c(k-1)=\lambda(n-1)$, and we're done.
4. Can there be a block-design of type 2-(16, 6, 1)? What about 2-(21, 6, 1)? $2-(25,10,3)$ ?

Solution: None are possible. First note that by double-counting, we have $m=\lambda \frac{n(n-1)}{k(k-1)}$. Now apply Fisher's inequality to get a contradiction on $m$.
5. For $\lambda=1$, prove Fisher's inequality directly.

Solution: Take any set $A$ with $k$ elements. As each element lies in $\frac{n-1}{k-1}$ sets, then for each element $i \in A$, there are $(n-1) /(k-1)-1$ other sets containing $i$. Over all $i \in A$, these sets are distinct, so the total number of sets are at least $\left(\frac{n-1}{k-1}-1\right) k+1$. It can be verified that this is at least $n$.

Bonus Problem. Suppose we have a necklace of $n$ beads. Each bead is labelled with an integer and the sum of all these labels is $n-1$. Prove that we can cut the necklace to form a string whose consecutive labels $x_{1}, x_{2}, \ldots, x_{n}$ satisfy

$$
\sum_{i=1}^{k} x_{i} \leq k-1 \quad \forall k=1,2, \ldots, n
$$

